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Hitachi Single-Chip Microcomputer

H8/3802 Series

 H8/3802
 HD6473802, HD6433802

 H8/3801
 HD6433801

 H8/3800
 HD6433800

Hardware Manual

HITACHI

ADE-602-203A Rev. 2.0 1/9/01 Hitachi Ltd.





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Preface

The H8/300L Series of single-chip microcomputers has the high-speed H8/300L CPU at its core, with many necessary peripheral functions on-chip. The H8/300L CPU instruction set is compatible with the H8/300 CPU.

The H8/3802 Series has a system-on-a-chip architecture that includes such peripheral functions as an LCD controller/driver, three timers, a two-channel 10-bit PWM, a serial communication interface, and an A/D converter. This allows H8/3802 Series devices to be used as embedded microcomputers in systems requiring LCD display.

This manual describes the hardware of the H8/3802 Series. For details on the H8/3802 Series instruction set, refer to the H8/300L Series Programming Manual.

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Section 1 Overview

1.1 Overview

The H8/300L Series is a series of single-chip microcomputers (MCU: microcomputer unit), built around the high-speed H8/300L CPU and equipped with peripheral system functions on-chip.

Within the H8/300L Series, the H8/3802 Series comprises single-chip microcomputers equipped with a controller/driver. Other on-chip peripheral functions include three timers, a two-channel 10-bit pulse width modulator (PWM), a serial communication interface, and an A/D converter. Together, these functions make the H8/3800 Series ideally suited for embedded applications in systems requiring low power consumption and LCD display. Models in the H8/3802 Series are the H8/3802, with on-chip 16-kbyte ROM and 1-kbyte RAM, the H8/3801, with 12-kbyte ROM and 512 byte RAM, and the H8/3800, with 8-kbyte ROM and 512 byte RAM.

The H8/3802 is also available in a ZTAT^{TM*} version with on-chip PROM which can be programmed as required by the user.

Table 1.1 summarizes the features of the H8/3802 Series.

Note: * ZTAT (Zero Turn Around Time) is a trademark of Hitachi, Ltd.

Table 1.1Features

Item	Specification							
CPU	High-speed H8/300L CPU							
	General-register architecture General registers: Sixteen 8-bit registers (can be used as eight 16-bit							
	registers)							
	Operating speed							
	 Max. operating speed: 8 MHz 							
	 — Add/subtract: 0.25 μs (operating at 8 MHz) 							
	— Multiply/divide: 1.75 μs (operating at 8 MHz)							
	 — Can run on 32.768 kHz or 38.4 kHz subclock 							
	 Instruction set compatible with H8/300 CPU 							
	 Instruction length of 2 bytes or 4 bytes 							
	 Basic arithmetic operations between registers 							
	 MOV instruction for data transfer between memory and registers 							
	Typical instructions							
	— Multiply (8 bits \times 8 bits)							
	— Divide (16 bits ÷ 8 bits)							
	— Bit accumulator							
	Register-indirect designation of bit position							
Interrupts	18 interrupt sources							
	 11 external interrupt sources (IRQ₁, IRQ₀, WKP₇ to WKP₀, IRQAEC) 							
	7 internal interrupt sources							
Clock pulse generators	Two on-chip clock pulse generators							
	System clock pulse generator: 1.0 to 16 MHz							
	Subclock pulse generator: 32.768 kHz, 38.4 kHz							
Power-down modes	Seven power-down modes							
	Sleep (high-speed) mode							
	Sleep (medium-speed) mode							
	Standby mode							
	Watch mode							
	Subsleep mode							
	Subactive mode							
	Active (medium-speed) mode							

Item	Specification							
Memory	Large on-chip memory							
	H8/3802: 16-kbyte ROM, 1-kbyte RAM							
	• H8/3801: 12-kbyte ROM, 512 byte RAM							
	• H8/3800: 8-kbyte ROM, 512 byte RAM							
I/O ports	50 pins							
	• 39 I/O pins							
	5 input pins							
	6 output pins							
Timers	Three on-chip timers							
	Timer A: 8-bit timer							
	Count-up timer with selection of eight internal clock signals divided from the system clock (\emptyset)* and four clock signals divided from the watch clock (\emptyset_w)*							
	Asynchronous event counter: 16-bit timer							
	 Count-up timer able to count asynchronous external events independently of the MCU's internal clocks 							
	Asynchronous external events can be counted (both rising and falling edge detection possible)							
	Timer F: 16-bit timer							
	 Can be used as two independent 8-bit timers 							
	 Count-up by an event input from the four internal clocks 							
	 Provision for toggle output by means of compare-match function 							
Serial communication	SCI3: 8-bit synchronous/asynchronous serial interface							
interface	Incorporates multiprocessor communication function							
10-bit PWM	Pulse-division PWM output for reduced ripple							
	 Can be used as a 10-bit D/A converter by connecting to an external low-pass filter. 							
A/D converter	Successive approximations using a resistance ladder							
	4-channel analog input pins							
	Conversion time: 31/ø or 62/ø per channel							

ltem	Specification										
LCD controller/driver	LCD controller/driver equipped with a maximum of 25 segment pins and four common pins										
	• Choice of four duty cycles (static, 1/2, 1/3, or 1/4)										
	Segment pin	s can be switched	to general-purpose port fu	nction in 4-bit units							
Product lineup	Produ	ct Code									
	Mask ROM Version	ZTAT Version	– Package	ROM/RAM Size							
	HD6433802H	HD6473802H	64-pin QFP (FP-64A)	ROM 16 kbytes							
	HD6433802FP	HD6473802FP	64-pin LQFP (FP-64E)	RAM 1 kbytes							
	HD6433802P	HD6473802P	64-pin DILP (DP-64S)	_							
	HD6433801H	·	64-pin QFP (FP-64A)	ROM 12 kbytes							
	HD6433801FP	·	64-pin LQFP (FP-64E)	RAM 512 bytes							
	HD6433801P	———	64-pin DILP (DP-64S)								
	HD6433800H	— —	64-pin QFP (FP-64A)	ROM 8 kbytes							
	HD6433800FP	— —	64-pin LQFP (FP-64E)	RAM 512 bytes							
	HD6433800P		64-pin DILP (DP-64S)								

Note: * See section 4, Clock Pulse Generator, for the definition of ø and. $ø_w$.

1.2 Internal Block Diagram

Figure 1.1 shows a block diagram of the H8/3802 Series.

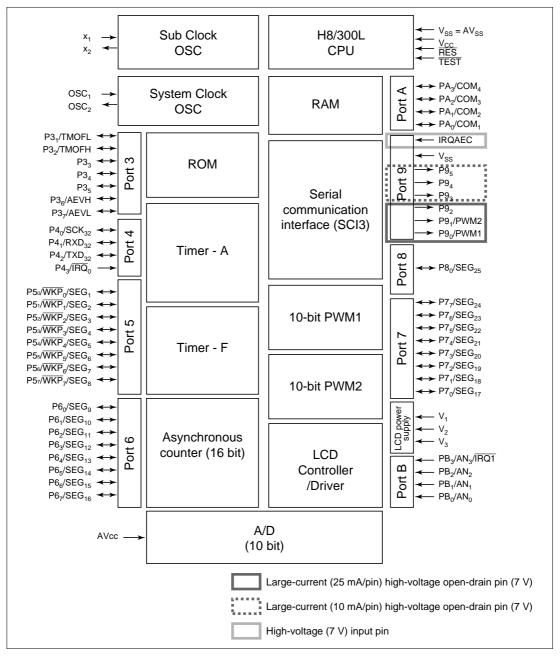


Figure 1.1 Block Diagram

1.3 Pin Arrangement and Functions

1.3.1 Pin Arrangement

The H8/3802 Series pin arrangement is shown in figures 1.2 and 1.3.

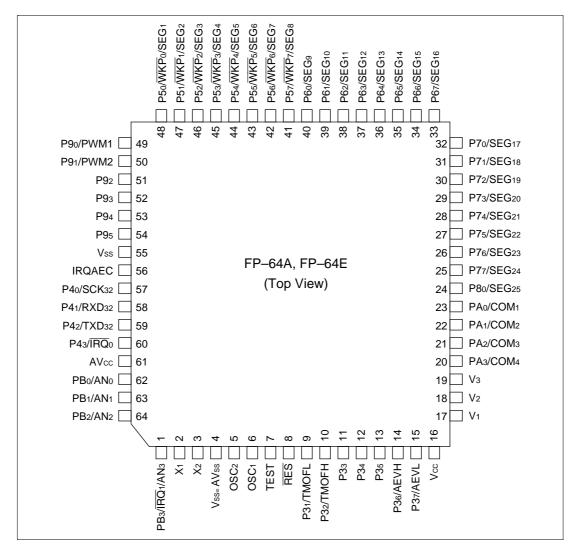


Figure 1.2 Pin Arrangement (FP-64A, FP-64E: Top View)

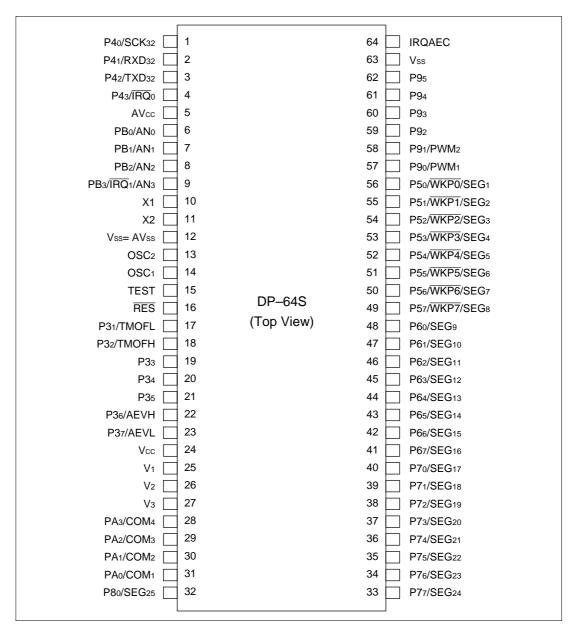


Figure 1.3 Pin Arrangement (DP-64S: Top View)

1.3.2 Pin Functions

Table 1.2 outlines the pin functions of the H8/3802 Series.

Table 1.2Pin Functions

		Pin	No.					
Туре	Symbol	FP-64A FP-64E	DP-64S	I/O	Name and Functions			
Power source pins	V _{CC}	16	24	Input	Power supply: All V _{CC} pins should be connected to the system power supply.			
	V _{SS}	4 (= AV _{ss}) 55	12 (= AV _{ss}) 63	Input	Ground: All V _{SS} pins should be connected to the system power supply (0 V).			
	AV _{CC} 61 5		5	Input	Analog power supply: This is the power supply pin for the A/D converter. When the A/D converter is not used, connect this pin to the system power supply.			
	AV _{SS}	4 (= V _{SS})	12 (= V _{SS})	Input	Analog ground: This is the A/D converter ground pin. It should be connected to the system power supply (0V).			
	V_1 V_2 V_3	17 18 19	25 26 27	Input	LCD power supply: These are the power supply pins for the LCD controller/driver.			
Clock pins	OSC ₁	6	14	Input	These pins connect to a crystal or			
	OSC ₂	5	13	Output	ceramic oscillator, or can be used to input an external clock. See section 4, Clock Pulse Generators, for a typical connection diagram.			
	X ₁	1 2 10		Input	These pins connect to a 32.768-kHz or			
	X ₂	3	11	Output	 38.4-kHz crystal oscillator. See section 4, Clock Pulse Generators, for a typical connection diagram. 			
System control	RES	8	16	Input	Reset: When this pin is driven low, the chip is reset			
	TEST	7	15	Input	Test pin: This pin is reserved and cannot be used. It should be connected to V_{SS} .			

	Pin No.				
Туре	Symbol	FP-64A FP-64E	DP-64S	 I/O	Name and Functions
Interrupt pins	IRQ ₀ 60 4 Input IRQ ₁ 1 9		IRQ interrupt request 0 and 1: These are input pins for edge-sensitive external interrupts, with a selection of rising or falling edge		
	IRQAEC	56	64	Input	Asynchronous event counter event signal: This is an interrupt input pin for enabling asynchronous event input.
	\overline{WKP}_7 to \overline{WKP}_0	41 to 48	49 to 56	Input	Wakeup interrupt request 0 to 7: These are input pins for rising or falling- edge-sensitive external interrupts.
Timer pins	AEVL AEVH	15 14	23 22	Input	Asynchronous event counter event input: This is an event input pin for input to the asynchronous event counter.
	TMOFL	9	17	Output	Timer FL output: This is an output pin for waveforms generated by the timer FL output compare function.
	TMOFH	10	18	Output	Timer FH output: This is an output pin for waveforms generated by the timer FH output compare function.
10-bit PWM pin	PWM1 PWM2	49 50	57 58	Output	10-bit PWM output: These are output pins for waveforms generated by the channel 1 and 2 10-bit PWMs.
I/O ports	P3 ₇ to P3 ₁	15 to 9	23 to 17	I/O	Port 3: This is an 7-bit I/O port. Input or output can be designated for each bit by means of port control register 3 (PCR3).
	P43	60	4	Input	Port 4 (bit 3): This is a 1-bit input port.
	P4 ₂ to P4 ₀	59 to 57	3 to 1	I/O	Port 4 (bits 2 to 0): This is a 3-bit I/O port. Input or output can be designated for each bit by means of port control register 4 (PCR4).
	P5 ₇ to P5 ₀	41 to 48	49 to 56	I/O	Port 5: This is an 8-bit I/O port. Input or output can be designated for each bit by means of port control register 5 (PCR5).

		Pin	No.					
Туре	Symbol	FP-64A FP-64E	DP-64S	I/O	Name and Functions			
I/O ports	I/O ports P6 ₇ to P6 ₀ 33 to 40 41 to 48 I/O		I/O	Port 6: This is an 8-bit I/O port. Input or output can be designated for each bit by means of port control register 6 (PCR6).				
	P7 ₇ to P7 ₀	25 to 32	33 to 40	I/O	Port 7: This is an 8-bit I/O port. Input or output can be designated for each bit by means of port control register 7 (PCR7).			
	P8 ₀	24	32	I/O	Port 8: This is an 8-bit I/O port. Input or output can be designated for each bit by means of port control register 8 (PCR8).			
	P9 ₅ to P9 ₀	54 to 49	62 to 57	Output	Port 9: This is a 6-bit output port.			
	PA_3 to PA_0	20 to 23	28 to 31	I/O	Port A: This is a 4-bit I/O port. Input or output can be designated for each bit by means of port control register A (PCRA).			
	PB ₃ to PB ₀	1, 64 to 62	9 to 6	Input	Port B: This is a 4-bit input port.			
Serial communi-	RXD ₃₂	58	2	Input	SCI3 receive data input: This is the SCI3 data input pin.			
cation interface	TXD ₃₂	59	3	Output	SCI3 transmit data output: This is the SCI3 data output pin.			
(SCI)	SCK ₃₂	57	1	I/O	SCI3 clock I/O: This is the SCI3 clock I/O pin.			
A/D converter	AN3 to An0	1 64 to 62	9 to 6	Input	Analog input channels 3 to 0: These are analog data input channels to the A/D converter			
LCD controller/	COM ₄ to COM ₁	20 to 23	28 to 31	Output	LCD common output: These are the LCD common output pins.			
driver	SEG ₂₅ to SEG ₁	24 to 48	32 to 56	Output	LCD segment output: These are the LCD segment output pins.			

Section 2 CPU

2.1 Overview

The H8/300L CPU has sixteen 8-bit general registers, which can also be paired as eight 16-bit registers. Its concise instruction set is designed for high-speed operation.

2.1.1 Features

Features of the H8/300L CPU are listed below.

- General-register architecture
 - Sixteen 8-bit general registers, also usable as eight 16-bit general registers
- Instruction set with 55 basic instructions, including:
 - Multiply and divide instructions
 - Powerful bit-manipulation instructions
- Eight addressing modes
 - Register direct
 - Register indirect
 - Register indirect with displacement
 - Register indirect with post-increment or pre-decrement
 - Absolute address
 - Immediate
 - Program-counter relative
 - Memory indirect
 - 64-kbyte address space
- High-speed operation
 - All frequently used instructions are executed in two to four states
 - High-speed arithmetic and logic operations
 - 8- or 16-bit register-register add or subtract: 0.25 μs*
 - 8×8 -bit multiply: 1.75 µs^{*}
 - $16 \div 8$ -bit divide: $1.75 \ \mu s^*$

Note: * These values are at $\phi = 8$ MHz.

• Low-power operation modes

SLEEP instruction for transfer to low-power operation

2.1.2 Address Space

The H8/300L CPU supports an address space of up to 64 kbytes for storing program code and data.

See 2.8, Memory Map, for details of the memory map.

2.1.3 Register Configuration

Figure 2.1 shows the register structure of the H8/300L CPU. There are two groups of registers: the general registers and control registers.

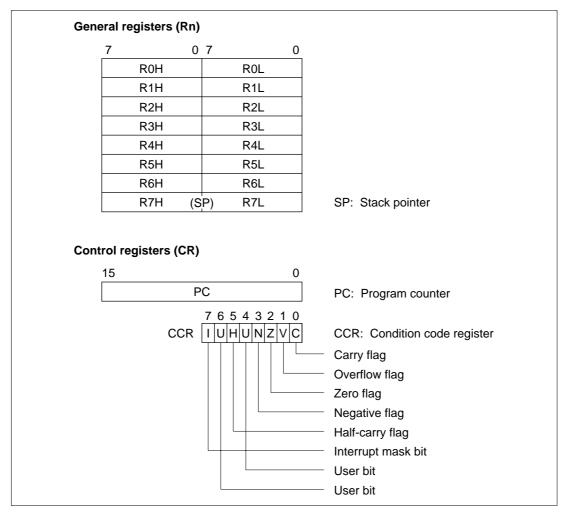


Figure 2.1 CPU Registers

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2.2 Register Descriptions

2.2.1 General Registers

All the general registers can be used as both data registers and address registers.

When used as data registers, they can be accessed as 16-bit registers (R0 to R7), or the high bytes (R0H to R7H) and low bytes (R0L to R7L) can be accessed separately as 8-bit registers.

When used as address registers, the general registers are accessed as 16-bit registers (R0 to R7).

R7 also functions as the stack pointer (SP), used implicitly by hardware in exception processing and subroutine calls. When it functions as the stack pointer, as indicated in figure 2.2, SP (R7) points to the top of the stack.

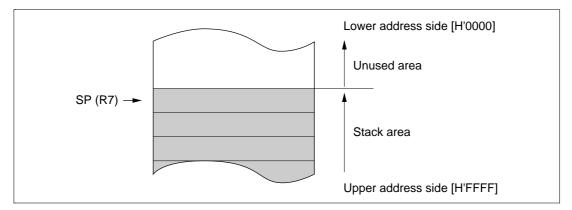


Figure 2.2 Stack Pointer

2.2.2 Control Registers

The CPU control registers include a 16-bit program counter (PC) and an 8-bit condition code register (CCR).

Program Counter (PC): This 16-bit register indicates the address of the next instruction the CPU will execute. All instructions are fetched 16 bits (1 word) at a time, so the least significant bit of the PC is ignored (always regarded as 0).

Condition Code Register (CCR): This 8-bit register contains internal status information, including the interrupt mask bit (I) and half-carry (H), negative (N), zero (Z), overflow (V), and carry (C) flags. These bits can be read and written by software (using the LDC, STC, ANDC, ORC, and XORC instructions). The N, Z, V, and C flags are used as branching conditions for conditional branching (Bcc) instructions.

Bit 7—Interrupt Mask Bit (I): When this bit is set to 1, interrupts are masked. This bit is set to 1 automatically at the start of exception handling. The interrupt mask bit may be read and written by software. For further details, see section 3.3, Interrupts.

Bit 6—User Bit (U): Can be used freely by the user.

Bit 5—Half-Carry Flag (H): When the ADD.B, ADDX.B, SUB.B, SUBX.B, CMP.B, or NEG.B instruction is executed, this flag is set to 1 if there is a carry or borrow at bit 3, and is cleared to 0 otherwise.

The H flag is used implicitly by the DAA and DAS instructions.

When the ADD.W, SUB.W, or CMP.W instruction is executed, the H flag is set to 1 if there is a carry or borrow at bit 11, and is cleared to 0 otherwise.

Bit 4—User Bit (U): Can be used freely by the user.

Bit 3—Negative Flag (N): Indicates the most significant bit (sign bit) of the result of an instruction.

Bit 2—Zero Flag (Z): Set to 1 to indicate a zero result, and cleared to 0 to indicate a non-zero result.

Bit 1—Overflow Flag (V): Set to 1 when an arithmetic overflow occurs, and cleared to 0 at other times.

Bit 0—Carry Flag (C): Set to 1 when a carry occurs, and cleared to 0 otherwise. Used by:

- Add instructions, to indicate a carry
- Subtract instructions, to indicate a borrow
- · Shift and rotate instructions, to store the value shifted out of the end bit

The carry flag is also used as a bit accumulator by bit manipulation instructions.

Some instructions leave some or all of the flag bits unchanged.

Refer to the H8/300L Series Programming Manual for the action of each instruction on the flag bits.

2.2.3 Initial Register Values

When the CPU is reset, the program counter (PC) is initialized to the value stored at address H'0000 in the vector table, and the I bit in the CCR is set to 1. The other CCR bits and the general registers are not initialized. In particular, the stack pointer (R7) is not initialized. The stack pointer should be initialized by software, by the first instruction executed after a reset.

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2.3 Data Formats

The H8/300L CPU can process 1-bit data, 4-bit (BCD) data, 8-bit (byte) data, and 16-bit (word) data.

- Bit manipulation instructions operate on 1-bit data specified as bit n in a byte operand (n = 0, 1, 2, ..., 7).
- All arithmetic and logic instructions except ADDS and SUBS can operate on byte data.
- The MOV.W, ADD.W, SUB.W, CMP.W, ADDS, SUBS, MULXU (8 bits × 8 bits), and DIVXU (16 bits ÷ 8 bits) instructions operate on word data.
- The DAA and DAS instructions perform decimal arithmetic adjustments on byte data in packed BCD form. Each nibble of the byte is treated as a decimal digit.

2.3.1 Data Formats in General Registers

Data of all the sizes above can be stored in general registers as shown in figure 2.3.

Data Type	Register	No.						Da	ata F	orm	at						
		7							0								
1-bit data	RnH	7	6	5	4	3	2	1	0				don't	care			
		;								7			1				0
1-bit data	RnL				don'i	t care				7	6	5	4	3	2	1	0
Byte data	RnH	7 MSB	1			1			0 LSB	T			don't	care			
Dyle dala		IVIOD	L	I		I		L I	LOD]			uoni				
										7							0
Byte data	RnL				don'	t care	•••••			MSB		1		1			LSB
		15	1	1	1	1				1							0
Word data	Rn	MSB	I	ı		ı				I	I	I	I	L		I	LSB
4-bit BCD data	RnH	7	Uppe	r digit	4	3	Lowe	r digit	0	I			 don't	care			
	IXIII I		- oppe	laight	1		20110	- digit]			uoni	Care			
										7			4	3			0
4-bit BCD data	RnL				don't	t care					Uppe	r digit	1		Lowe	r digit	
Notation: RnH: Upper by RnL: Lower by MSB: Most sigr LSB: Least sig	te of gene	ral reç															

Figure 2.3 Register Data Formats

2.3.2 Memory Data Formats

Figure 2.4 indicates the data formats in memory. The H8/300L CPU can access word data stored in memory (MOV.W instruction), but the word data must always begin at an even address. If word data starting at an odd address is accessed, the least significant bit of the address is regarded as 0, and the word data starting at the preceding address is accessed. The same applies to instruction codes.

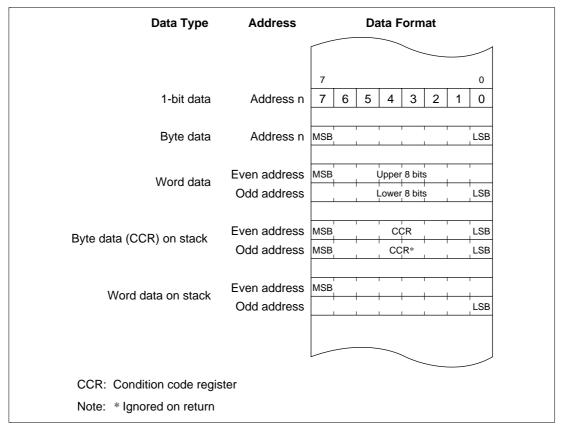


Figure 2.4 Memory Data Formats

When the stack is accessed using R7 as an address register, word access should always be performed. When the CCR is pushed on the stack, two identical copies of the CCR are pushed to make a complete word. When they are restored, the lower byte is ignored.

2.4 Addressing Modes

2.4.1 Addressing Modes

The H8/300L CPU supports the eight addressing modes listed in table 2.1. Each instruction uses a subset of these addressing modes.

No.	Address Modes	Symbol
1	Register direct	Rn
2	Register indirect	@Rn
3	Register indirect with displacement	@(d:16, Rn)
4	Register indirect with post-increment Register indirect with pre-decrement	@Rn+ @-Rn
5	Absolute address	@aa:8 or @aa:16
6	Immediate	#xx:8 or #xx:16
7	Program-counter relative	@(d:8, PC)
8	Memory indirect	@@aa:8

 Table 2.1
 Addressing Modes

1. **Register Direct—Rn:** The register field of the instruction specifies an 8- or 16-bit general register containing the operand.

Only the MOV.W, ADD.W, SUB.W, CMP.W, ADDS, SUBS, MULXU (8 bits \times 8 bits), and DIVXU (16 bits \div 8 bits) instructions have 16-bit operands.

- 2. **Register Indirect**—@**Rn:** The register field of the instruction specifies a 16-bit general register containing the address of the operand in memory.
- **3. Register Indirect with Displacement**—@(**d:16**, **Rn**): The instruction has a second word (bytes 3 and 4) containing a displacement which is added to the contents of the specified general register to obtain the operand address in memory.

This mode is used only in MOV instructions. For the MOV.W instruction, the resulting address must be even.

4. Register Indirect with Post-Increment or Pre-Decrement—@Rn+ or @-Rn:

— Register indirect with post-increment—@Rn+

The @Rn+ mode is used with MOV instructions that load registers from memory.

The register field of the instruction specifies a 16-bit general register containing the address of the operand. After the operand is accessed, the register is incremented by 1 for MOV.B or 2 for MOV.W. For MOV.W, the original contents of the 16-bit general register must be even.

Register indirect with pre-decrement—@–Rn

The @-Rn mode is used with MOV instructions that store register contents to memory.

The register field of the instruction specifies a 16-bit general register which is decremented by 1 or 2 to obtain the address of the operand in memory. The register retains the decremented value. The size of the decrement is 1 for MOV.B or 2 for MOV.W. For MOV.W, the original contents of the register must be even.

5. Absolute Address—@aa:8 or @aa:16: The instruction specifies the absolute address of the operand in memory.

The absolute address may be 8 bits long (@aa:8) or 16 bits long (@aa:16). The MOV.B and bit manipulation instructions can use 8-bit absolute addresses. The MOV.B, MOV.W, JMP, and JSR instructions can use 16-bit absolute addresses.

For an 8-bit absolute address, the upper 8 bits are assumed to be 1 (H'FF). The address range is H'FF00 to H'FFFF (65280 to 65535).

6. Immediate #xx:8 or #xx:16: The instruction contains an 8-bit operand (#xx:8) in its second byte, or a 16-bit operand (#xx:16) in its third and fourth bytes. Only MOV.W instructions can contain 16-bit immediate values.

The ADDS and SUBS instructions implicitly contain the value 1 or 2 as immediate data. Some bit manipulation instructions contain 3-bit immediate data in the second or fourth byte of the instruction, specifying a bit number.

- **7. Program-Counter Relative**—@(**d:8, PC**): This mode is used in the Bcc and BSR instructions. An 8-bit displacement in byte 2 of the instruction code is sign-extended to 16 bits and added to the program counter contents to generate a branch destination address. The possible branching range is -126 to +128 bytes (-63 to +64 words) from the current address. The displacement should be an even number.
- 8. Memory Indirect—@@aa:8: This mode can be used by the JMP and JSR instructions. The second byte of the instruction code specifies an 8-bit absolute address. The word located at this address contains the branch destination address.

The upper 8 bits of the absolute address are assumed to be 0 (H'00), so the address range is from H'0000 to H'00FF (0 to 255). Note that with the H8/300L Series, the lower end of the address area is also used as a vector area. See 3.3, Interrupts, for details on the vector area.

If an odd address is specified as a branch destination or as the operand address of a MOV.W instruction, the least significant bit is regarded as 0, causing word access to be performed at the address preceding the specified address. See 2.3.2, Memory Data Formats, for further information.

2.4.2 Effective Address Calculation

Table 2.2 shows how effective addresses are calculated in each of the addressing modes.

Arithmetic and logic instructions use register direct addressing (1). The ADD.B, ADDX, SUBX, CMP.B, AND, OR, and XOR instructions can also use immediate addressing (6).

Data transfer instructions can use all addressing modes except program-counter relative (7) and memory indirect (8).

Bit manipulation instructions can use register direct (1), register indirect (2), or 8-bit absolute addressing (5) to specify the operand. Register indirect (1) (BSET, BCLR, BNOT, and BTST instructions) or 3-bit immediate addressing (6) can be used independently to specify a bit position in the operand.

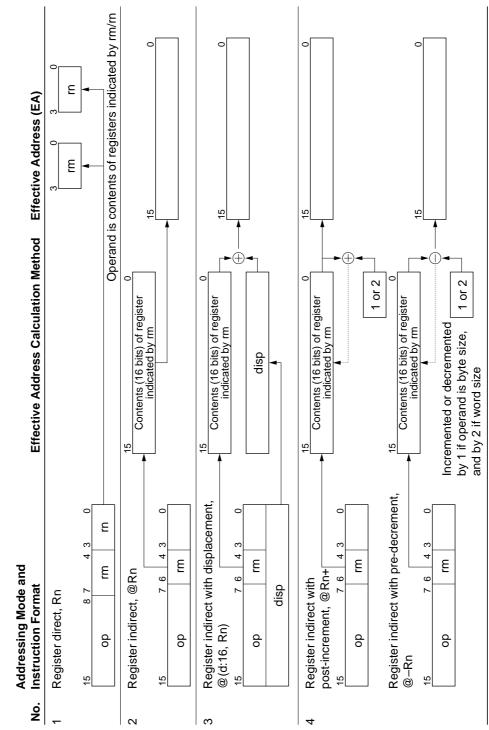
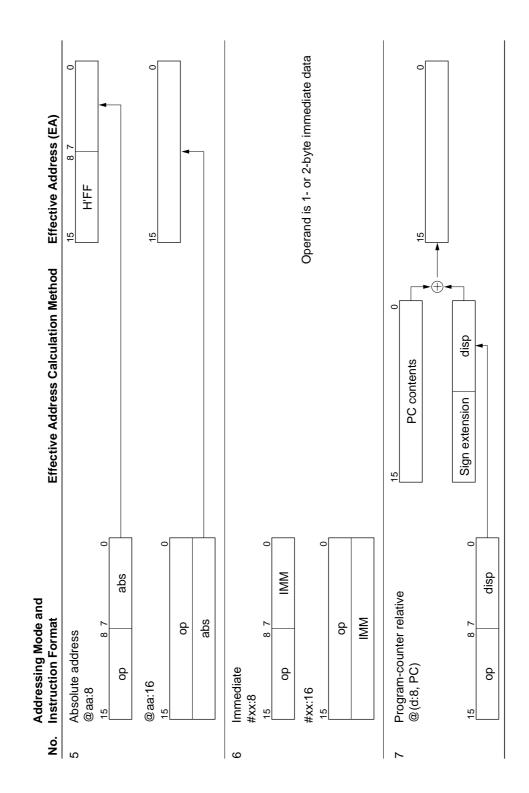


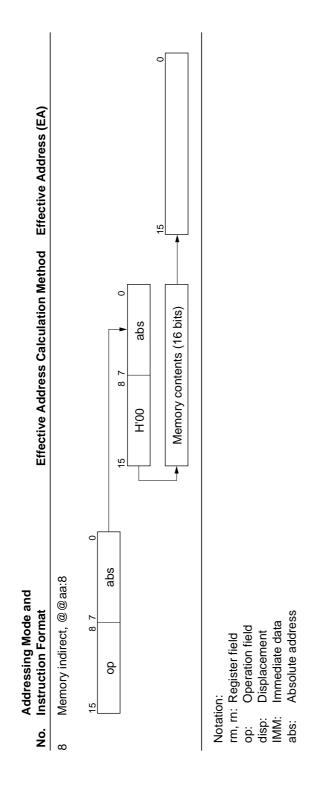
Table 2.2 Effective Address Calculation

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2.5 Instruction Set

The H8/300L Series can use a total of 55 instructions, which are grouped by function in table 2.3.

Table 2.3Instruction Set

Function	Instructions	Number
Data transfer	MOV, PUSH ^{*1} , POP ^{*1}	1
Arithmetic operations	ADD, SUB, ADDX, SUBX, INC, DEC, ADDS, SUBS, DAA, DAS, MULXU, DIVXU, CMP, NEG	14
Logic operations	AND, OR, XOR, NOT	4
Shift	SHAL, SHAR, SHLL, SHLR, ROTL, ROTR, ROTXL, ROTXR	8
Bit manipulation	BSET, BCLR, BNOT, BTST, BAND, BIAND, BOR, BIOR, BXOR, BIXOR, BLD, BILD, BST, BIST	14
Branch	Bcc ^{*2} , JMP, BSR, JSR, RTS	5
System control	RTE, SLEEP, LDC, STC, ANDC, ORC, XORC, NOP	8
Block data transfer	EEPMOV	1
		T () C

Total: 55

Notes: 1. PUSH Rn is equivalent to MOV.W Rn, @-SP.

POP Rn is equivalent to MOV.W @SP+, Rn. The same applies to the machine language.

2. Bcc is a conditional branch instruction in which cc represents a condition code.

The following sections give a concise summary of the instructions in each category, and indicate the bit patterns of their object code. The notation used is defined next.

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Notation	
Rd	General register (destination)
Rs	General register (source)
Rn	General register
(EAd), <ead></ead>	Destination operand
(EAs), <eas></eas>	Source operand
CCR	Condition code register
Ν	N (negative) flag of CCR
Z	Z (zero) flag of CCR
V	V (overflow) flag of CCR
С	C (carry) flag of CCR
PC	Program counter
SP	Stack pointer
#IMM	Immediate data
disp	Displacement
+	Addition
-	Subtraction
×	Multiplication
÷	Division
٨	AND logical
V	OR logical
\oplus	Exclusive OR logical
\rightarrow	Move
~	Logical negation (logical complement)
:3	3-bit length
:8	8-bit length
:16	16-bit length
(), < >	Contents of operand indicated by effective address

2.5.1 Data Transfer Instructions

Table 2.4 describes the data transfer instructions. Figure 2.5 shows their object code formats.

 Table 2.4
 Data Transfer Instructions

Instructio	n	Size*	Function				
MOV B/W		B/W	$(EAs) \to Rd, Rs \to (EAd)$				
			Moves data between two general registers or between a general register and memory, or moves immediate data to a general register.				
			The Rn, @Rn, @(d:16, Rn), @aa:16, #xx:16, @-Rn, and @Rn+ addressing modes are available for word data. The @aa:8 addressing mode is available for byte data only.				
			The @–R7 and @R7+ modes require word operands. Do not specify byte size for these two modes.				
POP		W	@SP+ \rightarrow Rn				
			Pops a 16-bit general register from the stack. Equivalent to MOV.W @SP+, Rn.				
PUSH		W	$Rn \rightarrow @-SP$				
			Pushes a 16-bit general register onto the stack. Equivalent to MOV.W Rn, @-SP.				
Notes: *	Size:	Operand size					
	B:	Byte					
	W:	Word					

Certain precautions are required in data access. See 2.9.1, Notes on Data Access, for details.

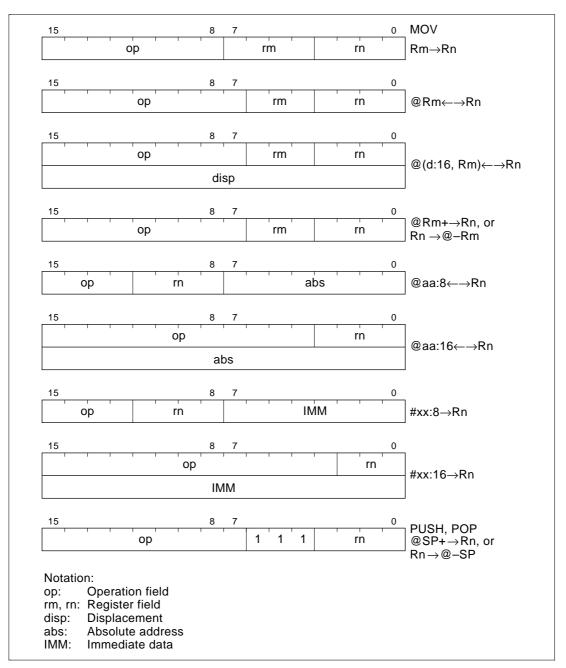


Figure 2.5 Data Transfer Instruction Codes

2.5.2 Arithmetic Operations

Table 2.5 describes the arithmetic instructions.

Table 2.5 Arithmetic Instruction
--

Instructio	n	Size*	Function			
ADD SUB B/W		B/W	$Rd \pm Rs \to Rd, Rd + \#IMM \to Rd$			
			Performs addition or subtraction on data in two general registers, or addition on immediate data and data in a general register. Immediate data cannot be subtracted from data in a general register. Word data can be added or subtracted only when both words are in general registers.			
ADDX SU	ЗX	В	$Rd \pm Rs \pm C \to Rd, Rd \pm \#IMM \pm C \to Rd$			
			Performs addition or subtraction with carry or borrow on byte data in two general registers, or addition or subtraction on immediate data and data in a general register.			
INC DEC		В	$Rd \pm 1 \rightarrow Rd$			
			Increments or decrements a general register by 1.			
ADDS SU	BS	W	$Rd \pm 1 \rightarrow Rd, Rd \pm 2 \rightarrow Rd$			
			Adds or subtracts 1 or 2 to or from a general register			
DAA DAS		В	Rd decimal adjust \rightarrow Rd			
			Decimal-adjusts (adjusts to 4-bit BCD) an addition or subtraction result in a general register by referring to the CCR			
MULXU B		В	$Rd \times Rs \rightarrow Rd$			
			Performs 8-bit \times 8-bit unsigned multiplication on data in two general registers, providing a 16-bit result			
DIVXU B		В	$Rd \div Rs \to Rd$			
			Performs 16-bit ÷ 8-bit unsigned division on data in two general registers, providing an 8-bit quotient and 8-bit remainder			
CMP		B/W	Rd – Rs, Rd – #IMM			
			Compares data in a general register with data in another general register or with immediate data, and indicates the result in the CCR. Word data can be compared only between two general registers.			
NEG		В	$0 - Rd \rightarrow Rd$			
			Obtains the two's complement (arithmetic complement) of data in a general register			
Notes: *	Size: B: W:	Operand size Byte Word				

2.5.3 **Logic Operations**

Table 2.6 describes the four instructions that perform logic operations.

Instruction	Size*	Function
AND	В	$Rd \land Rs \to Rd, Rd \land \#IMM \to Rd$
		Performs a logical AND operation on a general register and another general register or immediate data
OR	В	$Rd \lor Rs \to Rd, Rd \lor \#IMM \to Rd$
		Performs a logical OR operation on a general register and another general register or immediate data
XOR	В	$Rd \oplus Rs \to Rd, \ Rd \oplus \#IMM \to Rd$
		Performs a logical exclusive OR operation on a general register and another general register or immediate data
NOT	В	$\sim \text{Rd} \rightarrow \text{Rd}$
		Obtains the one's complement (logical complement) of general register contents
Notes: * Size	: Operand size	e

Table 2.6 **Logic Operation Instructions**

B: Byte

2.5.4 **Shift Operations**

Table 2.7 describes the eight shift instructions.

Table 2.7 **Shift Instructions**

Instruction	n	Size*	Function
SHAL		В	$Rd shift \to Rd$
SHAR			Performs an arithmetic shift operation on general register contents
SHLL		В	$Rd shift \to Rd$
SHLR			Performs a logical shift operation on general register contents
ROTL		В	$Rd rotate \rightarrow Rd$
ROTR			Rotates general register contents
ROTXL		В	Rd rotate through carry \rightarrow Rd
ROTXR			Rotates general register contents through the C (carry) bit
Notes: *	Size:	Operand size	9

B: Byte

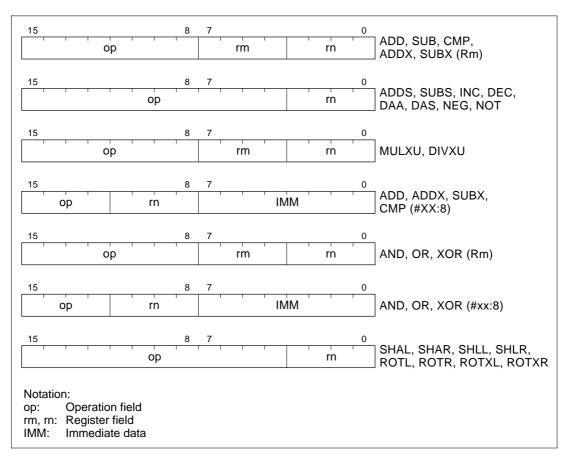


Figure 2.6 shows the instruction code format of arithmetic, logic, and shift instructions.

Figure 2.6 Arithmetic, Logic, and Shift Instruction Codes

2.5.5 Bit Manipulations

Table 2.8 describes the bit-manipulation instructions. Figure 2.7 shows their object code formats.

 Table 2.8
 Bit-Manipulation Instructions

Instructio	n	Size*	Function
BSET		В	$1 \rightarrow (\text{-bit-No.> of -EAd>})$
			Sets a specified bit in a general register or memory to 1. The bit number is specified by 3-bit immediate data or the lower three bits of a general register.
BCLR		В	$0 \rightarrow (\text{sbit-No.> of })$
			Clears a specified bit in a general register or memory to 0. The bit number is specified by 3-bit immediate data or the lower three bits of a general register.
BNOT		В	~ (<bit-no.> of <ead>) \rightarrow (<bit-no.> of <ead>)</ead></bit-no.></ead></bit-no.>
			Inverts a specified bit in a general register or memory. The bit number is specified by 3-bit immediate data or the lower three bits of a general register.
BTST		В	~ (<bit-no.> of <ead>) \rightarrow Z</ead></bit-no.>
			Tests a specified bit in a general register or memory and sets or clears the Z flag accordingly. The bit number is specified by 3-bit immediate data or the lower three bits of a general register.
BAND		В	$C \land (<\!bit-No.\!> of < EAd>) \to C$
			ANDs the C flag with a specified bit in a general register or memory, and stores the result in the C flag.
BIAND		В	$C \land [\text{~(of)}] \to C$
			ANDs the C flag with the inverse of a specified bit in a general register or memory, and stores the result in the C flag.
			The bit number is specified by 3-bit immediate data.
BOR		В	$C \lor (<\!bit-No.\!> of <\!EAd\!>) \to C$
			ORs the C flag with a specified bit in a general register or memory, and stores the result in the C flag.
BIOR		В	$C \lor [\text{~(of)}] \to C$
			ORs the C flag with the inverse of a specified bit in a general register or memory, and stores the result in the C flag.
			The bit number is specified by 3-bit immediate data.
Notes: *	Size:	Operand size	;
	D .	Dute	

B: Byte

Instructio	n	Size*	Function
BXOR		В	$C \oplus (<\!bit\!-\!No.\!> of <\!\mathsf{EAd\!\!>}) \to C$
			XORs the C flag with a specified bit in a general register or memory, and stores the result in the C flag.
BIXOR		В	$C \oplus \ [\text{-(of)}] \to C$
			XORs the C flag with the inverse of a specified bit in a general register or memory, and stores the result in the C flag.
			The bit number is specified by 3-bit immediate data.
BLD		В	(<bit-no.> of <ead>) \rightarrow C</ead></bit-no.>
			Copies a specified bit in a general register or memory to the C flag.
BILD		В	~ (<bit-no.> of <ead>) \rightarrow C</ead></bit-no.>
			Copies the inverse of a specified bit in a general register or memory to the C flag.
			The bit number is specified by 3-bit immediate data.
BST		В	$C \rightarrow (\text{sbit-No.> of })$
			Copies the C flag to a specified bit in a general register or memory.
BIST		В	~ C \rightarrow (<bit-no.> of <ead>)</ead></bit-no.>
			Copies the inverse of the C flag to a specified bit in a general register or memory.
			The bit number is specified by 3-bit immediate data.
Notes: *	Size: B:	Operand size Byte	

Certain precautions are required in bit manipulation. See 2.9.2, Notes on Bit Manipulation, for details.

15		8	7					0	BSET, BCLR, BNOT, BTST
	op	0	7	IMM		r	า	0	Operand: register direct (Rn) Bit No.: immediate (#xx:3)
15		8	7					0	
	ор	Ĩ		rm		rı	٦		Operand: register direct (Rn) Bit No.: register direct (Rm)
15		8	7					0	_
	op	1		'n	0	0	0	0	Operand: register indirect (@Rn)
	ор			IMM	0	0	0	0	Bit No.: immediate (#xx:3)
15		8	7					0	_
	ор	I		rn	0	0	0	0	Operand: register indirect (@Rn)
	ор			rm	0	0	0	0	Bit No.: register direct (Rm)
15		8	7					0	
	op	1			abs	1			Operand: absolute (@aa:8)
	ор			IMM	0	0	0	0	Bit No.: immediate (#xx:3)
15		8	7					0	
	ор	1			abs	1			Operand: absolute (@aa:8)
	ор			rm	0	0	0	0	Bit No.: register direct (Rm)
									BAND, BOR, BXOR, BLD, BST
15		8	7					0	
	ор			IMM		rı	า		Operand: register direct (Rn) Bit No.: immediate (#xx:3)
15		8	7					0	
	op	-		rn	0	0	0	0	Operand: register indirect (@Rn)
	ор			IMM	0	0	0	0	Bit No.: immediate (#xx:3)
15		8	7					0	
	ор	I			abs	I			Operand: absolute (@aa:8)
	ор			IMM	0	0	0	0	Bit No.: immediate (#xx:3)
N 1 1 1									
Notation op:	Operation field								
	Register field Absolute address								
	Immediate data								



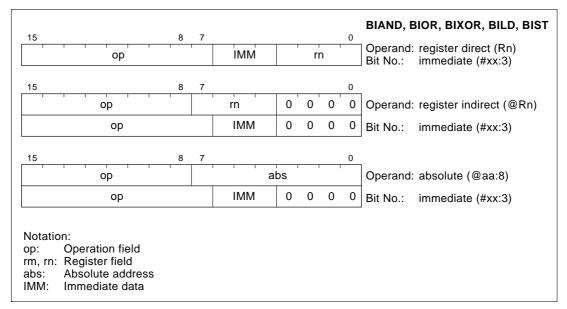


Figure 2.7 Bit Manipulation Instruction Codes (cont)

2.5.6 Branching Instructions

Table 2.9 describes the branching instructions. Figure 2.8 shows their object code formats.

Table 2.9Branching Instructions

Instruction	Size	Function	Function						
Bcc	_	Branches to the designated address if condition cc is true. The branching conditions are given below.							
			Description	Condition					
		BRA (BT)	Always (true)	Always					
		BRN (BF)	Never (false)	Never					
		BHI	High	$C \lor Z = 0$					
		BLS	Low or same	C ∨ Z = 1					
		BCC (BHS)	Carry clear (high or same)	C = 0					
		BCS (BLO)	Carry set (low)	C = 1					
		BNE	Not equal	Z = 0					
		BEQ	Equal	Z = 1					
		BVC	Overflow clear	V = 0					
		BVS	Overflow set	V = 1					
		BPL	Plus	N = 0					
		BMI	Minus	N = 1					
		BGE	Greater or equal	N ⊕ V = 0					
		BLT	Less than	N ⊕ V = 1					
		BGT	Greater than	$Z \vee (N \oplus V) = 0$					
		BLE	Less or equal	$Z \lor (N \oplus V) = 1$					
JMP		Branches unco	onditionally to a specified addres	S					
BSR	_	Branches to a	Branches to a subroutine at a specified address						
JSR		Branches to a	subroutine at a specified addres	S					
RTS		Returns from a subroutine							

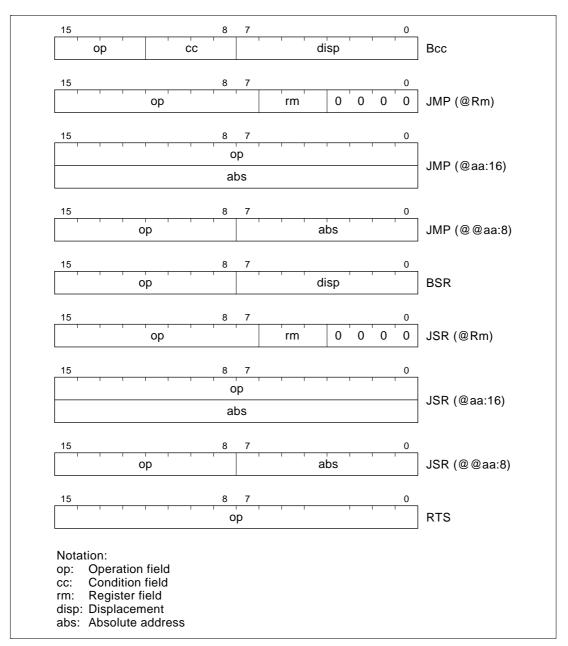


Figure 2.8 Branching Instruction Codes

2.5.7 System Control Instructions

Table 2.10 describes the system control instructions. Figure 2.9 shows their object code formats.

Table 2.10	System	Control	Instructions
-------------------	--------	---------	--------------

Instructio	n	Size*	Function
RTE		_	Returns from an exception-handling routine
SLEEP			Causes a transition from active mode to a power-down mode. See section 5, Power-Down Modes, for details.
LDC		В	$Rs \to CCR, \ \ \#IMM \to CCR$
			Moves immediate data or general register contents to the condition code register
STC		В	$CCR \rightarrow Rd$
			Copies the condition code register to a specified general register
ANDC		В	$CCR \land \#IMM \rightarrow CCR$
			Logically ANDs the condition code register with immediate data
ORC		В	$CCR \lor \#IMM \rightarrow CCR$
			Logically ORs the condition code register with immediate data
XORC		В	$CCR \oplus \#IMM \rightarrow CCR$
			Logically exclusive-ORs the condition code register with immediate data
NOP			$PC + 2 \rightarrow PC$
			Only increments the program counter
Notes: *	Size:	Operand size	9
	B٠	Buto	

B: Byte

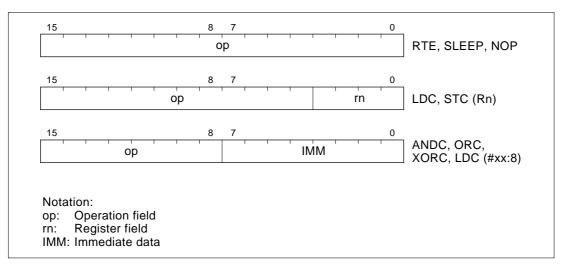


Figure 2.9 System Control Instruction Codes

2.5.8 Block Data Transfer Instruction

Table 2.11 describes the block data transfer instruction. Figure 2.10 shows its object code format.

 Table 2.11
 Block Data Transfer Instruction

Instruction	Size	Function
EEPMOV	_	If R4L ≠ 0 then
		repeat $@R5+ \rightarrow @R6+$ $R4L - 1 \rightarrow R4L$ untilR4L = 0
		else next;
		Block transfer instruction. Transfers the number of data bytes specified by R4L from locations starting at the address indicated by R5 to locations starting at the address indicated by R6. After the transfer, the next instruction is executed.

Certain precautions are required in using the EEPMOV instruction. See 2.9.3, Notes on Use of the EEPMOV Instruction, for details.

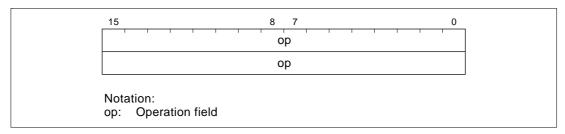


Figure 2.10 Block Data Transfer Instruction Code

2.6 Basic Operational Timing

CPU operation is synchronized by a system clock (\emptyset) or a subclock (\emptyset_{SUB}). For details on these clock signals see section 4, Clock Pulse Generators. The period from a rising edge of \emptyset or \emptyset_{SUB} to the next rising edge is called one state. A bus cycle consists of two states or three states. The cycle differs depending on whether access is to on-chip memory or to on-chip peripheral modules.

2.6.1 Access to On-Chip Memory (RAM, ROM)

Access to on-chip memory takes place in two states. The data bus width is 16 bits, allowing access in byte or word size. Figure 2.11 shows the on-chip memory access cycle.

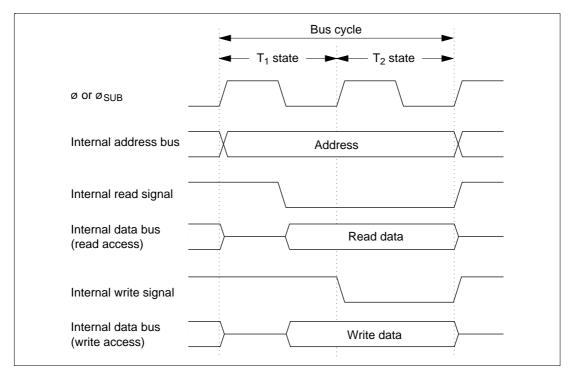


Figure 2.11 On-Chip Memory Access Cycle

2.6.2 Access to On-Chip Peripheral Modules

On-chip peripheral modules are accessed in two states or three states. The data bus width is 8 bits, so access is by byte size only. This means that for accessing word data, two instructions must be used. Figures 2.12 and 2.13 show the on-chip peripheral module access cycle.

Two-state access to on-chip peripheral modules

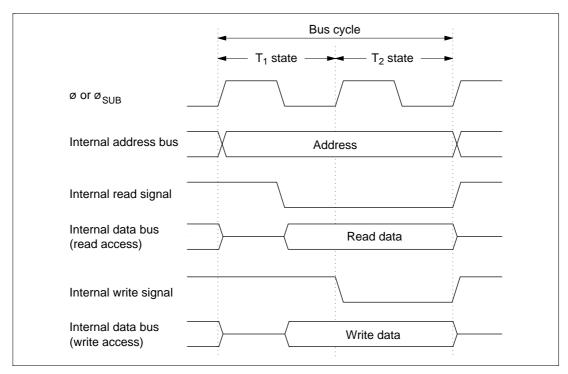
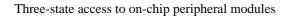


Figure 2.12 On-Chip Peripheral Module Access Cycle (2-State Access)



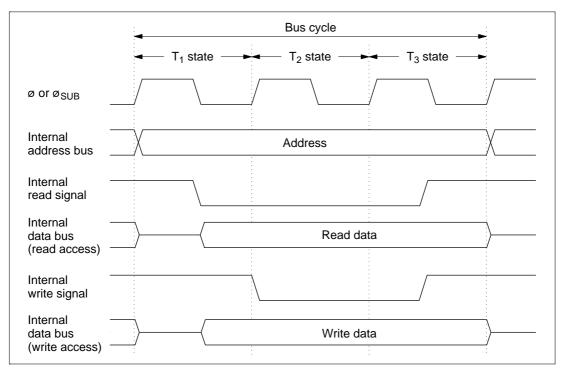


Figure 2.13 On-Chip Peripheral Module Access Cycle (3-State Access)

2.7 CPU States

2.7.1 Overview

There are four CPU states: the reset state, program execution state, program halt state, and exception-handling state. The program execution state includes active (high-speed or medium-speed) mode and subactive mode. In the program halt state there are a sleep (high-speed or medium-speed) mode, standby mode, watch mode, and sub-sleep mode. These states are shown in figure 2.14. Figure 2.15 shows the state transitions.

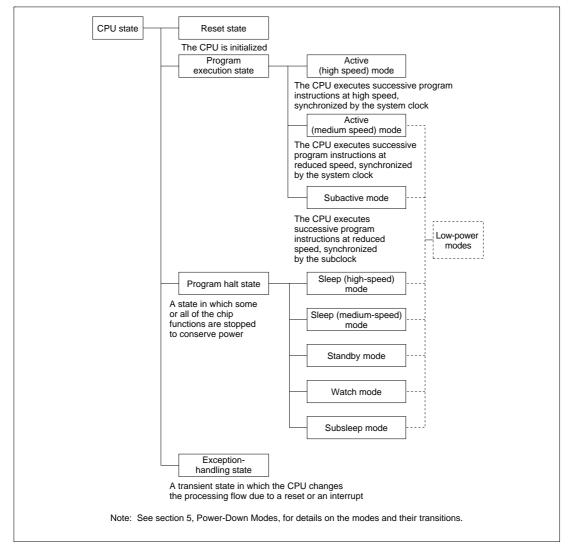


Figure 2.14 CPU Operation States

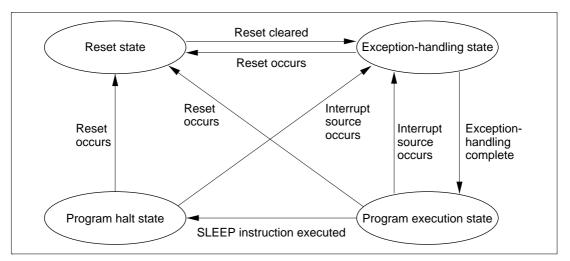


Figure 2.15 State Transitions

2.7.2 Program Execution State

In the program execution state the CPU executes program instructions in sequence.

There are three modes in this state, two active modes (high speed and medium speed) and one subactive mode. Operation is synchronized with the system clock in active mode (high speed and medium speed), and with the subclock in subactive mode. See section 5, Power-Down Modes for details on these modes.

2.7.3 Program Halt State

In the program halt state there are five modes: two sleep modes (high speed and medium speed), standby mode, watch mode, and subsleep mode. See section 5, Power-Down Modes for details on these modes.

2.7.4 Exception-Handling State

The exception-handling state is a transient state occurring when exception handling is started by a reset or interrupt and the CPU changes its normal processing flow. In exception handling caused by an interrupt, SP (R7) is referenced and the PC and CCR values are saved on the stack.

For details on interrupt handling, see section 3.3, Interrupts.

2.8 Memory Map

2.8.1 Memory Map

The memory map of the H8/3802 is shown in figure 2.16(1), that of the H8/3801 in figure 2.16(2), and that of the H8/3800 in figure 2.16(3).

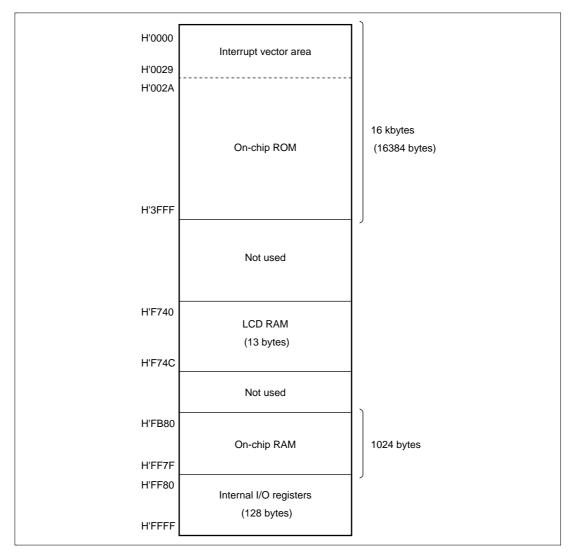


Figure 2.16 (1) H8/3802 Memory Map

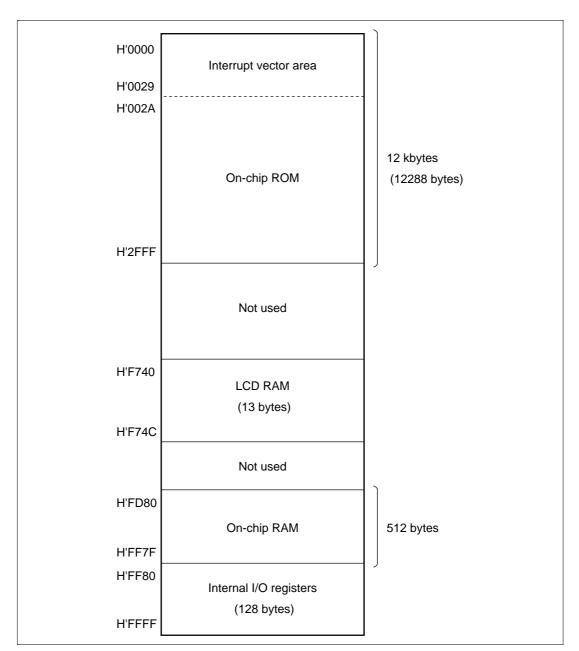


Figure 2.16 (2) H8/3801 Memory Map

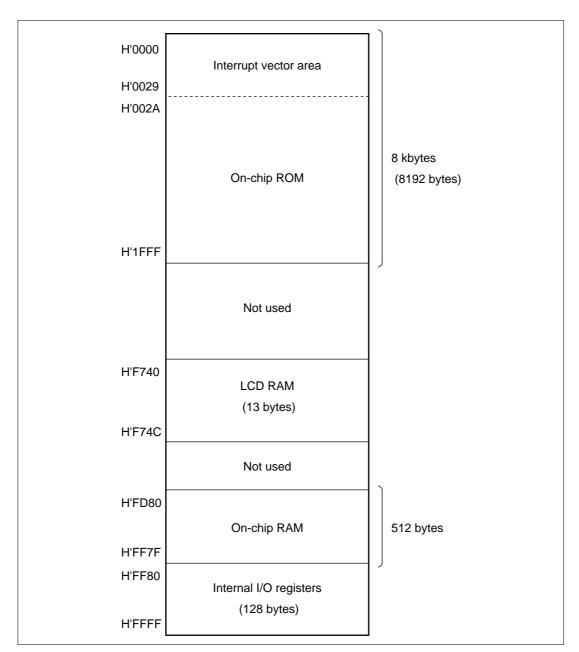


Figure 2.16 (3) H8/3800 Memory Map

2.9 Application Notes

2.9.1 Notes on Data Access

1. Access to Empty Areas:

The address space of the H8/300L CPU includes empty areas in addition to the RAM, registers, and ROM areas available to the user. If these empty areas are mistakenly accessed by an application program, the following results will occur.

Data transfer from CPU to empty area:

The transferred data will be lost. This action may also cause the CPU to misoperate.

Data transfer from empty area to CPU:

Unpredictable data is transferred.

2. Access to Internal I/O Registers:

Internal data transfer to or from on-chip modules other than the ROM and RAM areas makes use of an 8-bit data width. If word access is attempted to these areas, the following results will occur.

Word access from CPU to I/O register area:

Upper byte: Will be written to I/O register.

Lower byte: Transferred data will be lost.

Word access from I/O register to CPU:

Upper byte: Will be written to upper part of CPU register.

Lower byte: Unpredictable data will be written to lower part of CPU register.

Byte size instructions should therefore be used when transferring data to or from I/O registers other than the on-chip ROM and RAM areas. Figure 2.17 shows the data size and number of states in which on-chip peripheral modules can be accessed.

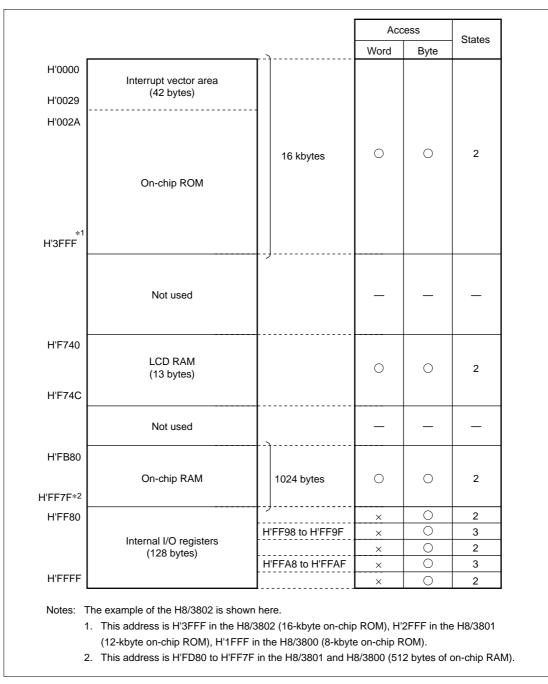


Figure 2.17 Data Size and Number of States for Access to and from On-Chip Peripheral Modules

2.9.2 Notes on Bit Manipulation

The BSET, BCLR, BNOT, BST, and BIST instructions read one byte of data, modify the data, then write the data byte again. Special care is required when using these instructions in cases where two registers are assigned to the same address, in the case of registers that include write-only bits, and when the instruction accesses an I/O port.

Orde	er of Operation	Operation
1	Read	Read byte data at the designated address
2	Modify	Modify a designated bit in the read data
3	Write	Write the altered byte data to the designated address

1. Bit manipulation in two registers assigned to the same address

Example 1: timer load register and timer counter

Figure 2.18 shows an example in which two timer registers share the same address. When a bit manipulation instruction accesses the timer load register and timer counter of a reloadable timer, since these two registers share the same address, the following operations take place.

Orde	er of Operation	Operation
1	Read	Timer counter data is read (one byte)
2	Modify	The CPU modifies (sets or resets) the bit designated in the instruction
3	Write	The altered byte data is written to the timer load register

The timer counter is counting, so the value read is not necessarily the same as the value in the timer load register. As a result, bits other than the intended bit in the timer load register may be modified to the timer counter value.

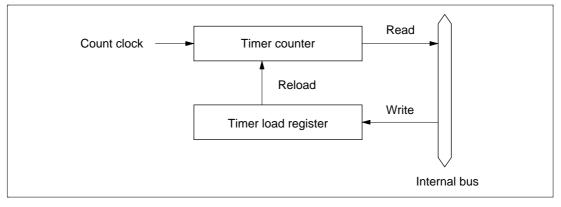


Figure 2.18 Timer Configuration Example

Example 2: BSET instruction executed designating port 3

 $P3_7$ and $P3_6$ are designated as input pins, with a low-level signal input at $P3_7$ and a high-level signal at $P3_6$. The remaining pins, $P3_5$ to $P3_1$, are output pins and output low-level signals. In this example, the BSET instruction is used to change pin $P3_1$ to high-level output.

[A: Prior to executing BSET]

	P3 ₇	P3 ₆	P3 ₅	P3 ₄	P3 ₃	P3 ₂	P3 ₁	_
Input/output	Input	Input	Output	Output	Output	Output	Output	_
Pin state	Low level	High level	Low level	Low level	Low level	Low level	Low level	—
PCR3	0	0	1	1	1	1	1	1
PDR3	1	0	0	0	0	0	0	1

[B: BSET instruction executed]

BSET	#1		@PDR3
DOGT	# 1	1	@PDR5

The BSET instruction is executed designating port 3.

[C: After executing BSET]

	P3 ₇	P3 ₆	P3 ₅	P3 ₄	P3 ₃	P3 ₂	P3 ₁	_
Input/output	Input	Input	Output	Output	Output	Output	Output	_
Pin state	Low level	High level	Low level	Low level	Low level	Low level	High level	_
PCR3	0	0	1	1	1	1	1	1
PDR3	0	1	0	0	0	0	1	1

[D: Explanation of how BSET operates]

When the BSET instruction is executed, first the CPU reads port 3.

Since $P3_7$ and $P3_6$ are input pins, the CPU reads the pin states (low-level and high-level input). $P3_5$ to $P3_1$ are output pins, so the CPU reads the value in PDR3. In this example PDR3 has a value of H'81, but the value read by the CPU is H'41.

Next, the CPU sets bit 1 of the read data to 1, changing the PDR3 data to H'43. Finally, the CPU writes this value (H'43) to PDR3, completing execution of BSET.

As a result of this operation, bit 1 in PDR3 becomes 1, and $P3_1$ outputs a high-level signal. However, bits 7 and 6 of PDR3 end up with different values.

To avoid this problem, store a copy of the PDR3 data in a work area in memory. Perform the bit manipulation on the data in the work area, then write this data to PDR3.

[A: Prior to executing BSET]

MOV.	В	#81,	ROL
MOV.	В	ROL,	@RAM0
MOV.	В	ROL,	@PDR3

The PDR3 value (H'81) is written to a work area in memory (RAM0) as well as to PDR3

	P3 ₇	P3 ₆	P3 ₅	P3 ₄	P3 ₃	P3 ₂	P3 ₁	_
Input/output	Input	Input	Output	Output	Output	Output	Output	_
Pin state	Low level	High level	Low level	Low level	Low level	Low level	Low level	_
PCR3	0	0	1	1	1	1	1	1
PDR3	1	0	0	0	0	0	0	1
RAM0	1	0	0	0	0	0	0	1

[B: BSET instruction executed]

BSET #1 , @RAM(

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The BSET instruction is executed designating the PDR3 work area (RAM0).

[C: After executing BSET]

MOV.	В	@RAM0,	ROL
MOV.	В	ROL,	@PDR3

The work area (RAM0) value is written to PDR3.

	P3 ₇	P3 ₆	P3 ₅	P3 ₄	P3 ₃	P3 ₂	P3 ₁	_
Input/output	Input	Input	Output	Output	Output	Output	Output	_
Pin state	Low level	High level	Low level	Low level	Low level	Low level	High level	—
PCR3	0	0	1	1	1	1	1	1
PDR3	1	0	0	0	0	0	1	1
RAM0	1	0	0	0	0	0	1	1

2. Bit manipulation in a register containing a write-only bit

Example 3: BCLR instruction executed designating port 3 control register PCR3

As in the examples above, P_{3_7} and P_{3_6} are input pins, with a low-level signal input at P_{3_7} and a high-level signal at P_{3_6} . The remaining pins, P_{3_5} to P_{3_1} , are output pins that output low-level signals. In this example, the BCLR instruction is used to change pin P_{3_1} to an input port. It is assumed that a high-level signal will be input to this input pin.

[A: Prior to executing BCLR]

	P3 ₇	P3 ₆	P3 ₅	P3 ₄	P3 ₃	P3 ₂	P3 ₁	-
Input/output	Input	Input	Output	Output	Output	Output	Output	_
Pin state	Low level	High level	Low level	Low level	Low level	Low level	Low level	_
PCR3	0	0	1	1	1	1	1	1
PDR3	1	0	0	0	0	0	0	1

[B: BCLR instruction executed]

BCLR #1 , @PCR3

The BCLR instruction is executed designating PCR3.

[C: After executing BCLR]

	P3 ₇	P3 ₆	P3 ₅	P3 ₄	P3 ₃	P3 ₂	P3 ₁	—
Input/output	Output	Output	Output	Output	Output	Output	Input	_
Pin state	Low level	High level	Low level	Low level	Low level	Low level	High level	_
PCR3	1	1	1	1	1	1	0	1
PDR3	1	0	0	0	0	0	0	1

[D: Explanation of how BCLR operates]

When the BCLR instruction is executed, first the CPU reads PCR3. Since PCR3 is a write-only register, the CPU reads a value of H'FF, even though the PCR3 value is actually H'3F.

Next, the CPU clears bit 1 in the read data to 0, changing the data to H'FD. Finally, this value (H'FD) is written to PCR3 and BCLR instruction execution ends.

As a result of this operation, bit 1 in PCR3 becomes 0, making $P3_1$ an input port. However, bits 7 and 6 in PCR3 change to 1, so that $P3_7$ and $P3_6$ change from input pins to output pins.

To avoid this problem, store a copy of the PCR3 data in a work area in memory. Perform the bit manipulation on the data in the work area, then write this data to PCR3.

[A: Prior to executing BCLR]

MOV.	В	#3F,	ROL
MOV.	В	ROL,	@RAM0
MOV.	В	ROL,	@PCR3

The PCR3 value (H'3F) is written to a work area in memory (RAM0) as well as to PCR3.

	P3 ₇	P3 ₆	P3 ₅	P3 ₄	P3 ₃	P3 ₂	P3 ₁	_
Input/output	Input	Input	Output	Output	Output	Output	Output	_
Pin state	Low level	High level	Low level	Low level	Low level	Low level	Low level	
PCR3	0	0	1	1	1	1	1	1
PDR3	1	0	0	0	0	0	0	1
RAM0	0	0	1	1	1	1	1	1

[B: BCLR instruction executed]

BCLR #1 , @RAMO

The BCLR instruction is executed designating the PCR3 work area (RAM0).

[C: After executing BCLR]

MOV.	В	@RAM0	, ROL
MOV.	В	ROL,	@PCR3

The work area (RAM0) value is written to PCR3.

	P3 ₇	P3 ₆	P3 ₅	P3 ₄	P3 ₃	P3 ₂	P3 ₁	_
Input/output	Input	Input	Output	Output	Output	Output	Output	_
Pin state	Low level	High level	Low level	Low level	Low level	Low level	High level	
PCR3	0	0	1	1	1	1	0	1
PDR3	1	0	0	0	0	0	0	1
RAM0	0	0	1	1	1	1	0	1

Table 2.12 lists the pairs of registers that share identical addresses. Table 2.13 lists the registers that contain write-only bits.

Table 2.12 Registers with Shared Addresses

Register Name	Abbreviation	Address
Port data register 3*	PDR3	H'FFD6
Port data register 4*	PDR4	H'FFD7
Port data register 5*	PDR5	H'FFD8
Port data register 6*	PDR6	H'FFD9
Port data register 7*	PDR7	H'FFDA
Port data register 8*	PDR8	H'FFDB
Port data register A*	PDRA	H'FFDD

Note: * Port data registers have the same addresses as input pins.

Table 2.13 Registers with Write-Only Bits

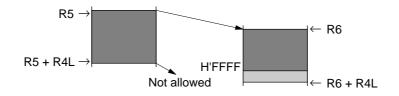
Register Name	Abbreviation	Address
Port control register 3	PCR3	H'FFE6
Port control register 4	PCR4	H'FFE7
Port control register 5	PCR5	H'FFE8
Port control register 6	PCR6	H'FFE9
Port control register 7	PCR7	H'FFEA
Port control register 8	PCR8	H'FFEB
Port control register A	PCRA	H'FFED
Timer control register F	TCRF	H'FFB6
PWM1 control register	PWCR1	H'FFD0
PWM1 data register U	PWDRU1	H'FFD1
PWM1 data register L	PWDRL1	H'FFD2
PWM2 control register	PWCR2	H'FFCD
PWM2 data register U	PWDRU2	H'FFCE
PWM2 data register L	PWDRL2	H'FFCF

2.9.3 Notes on Use of the EEPMOV Instruction

• The EEPMOV instruction is a block data transfer instruction. It moves the number of bytes specified by R4L from the address specified by R5 to the address specified by R6.



• When setting R4L and R6, make sure that the final destination address (R6 + R4L) does not exceed H'FFFF. The value in R6 must not change from H'FFFF to H'0000 during execution of the instruction.



Section 3 Exception Handling

3.1 Overview

Exception handling is performed in the H8/3802 Series when a reset or interrupt occurs. Table 3.1 shows the priorities of these two types of exception handling.

 Table 3.1
 Exception Handling Types and Priorities

Priority	Exception Source	Time of Start of Exception Handling
High	Reset	Exception handling starts as soon as the reset state is cleared
≜	Interrupt	When an interrupt is requested, exception handling starts after execution of the present instruction or the exception handling in
Low		progress is completed

3.2 Reset

3.2.1 Overview

A reset is the highest-priority exception. The internal state of the CPU and the registers of the onchip peripheral modules are initialized.

3.2.2 Reset Sequence

As soon as the $\overline{\text{RES}}$ pin goes low, all processing is stopped and the chip enters the reset state.

To make sure the chip is reset properly, observe the following precautions.

- At power on: Hold the $\overline{\text{RES}}$ pin low until the clock pulse generator output stabilizes.
- Resetting during operation: Hold the $\overline{\text{RES}}$ pin low for at least 10 system clock cycles.

Reset exception handling takes place as follows.

- The CPU internal state and the registers of on-chip peripheral modules are initialized, with the I bit of the condition code register (CCR) set to 1.
- The PC is loaded from the reset exception handling vector address (H'0000 to H'0001), after which the program starts executing from the address indicated in PC.

When system power is turned on or off, the $\overline{\text{RES}}$ pin should be held low.

Figure 3.1 shows the reset sequence starting from $\overline{\text{RES}}$ input.

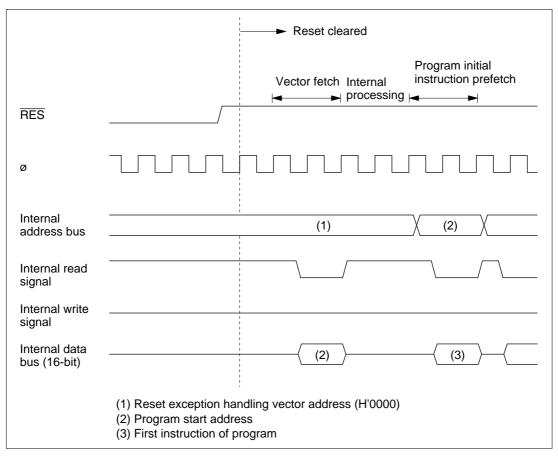


Figure 3.1 Reset Sequence

3.2.3 Interrupt Immediately after Reset

After a reset, if an interrupt were to be accepted before the stack pointer (SP: R7) was initialized, PC and CCR would not be pushed onto the stack correctly, resulting in program runaway. To prevent this, immediately after reset exception handling all interrupts are masked. For this reason, the initial program instruction is always executed immediately after a reset. This instruction should initialize the stack pointer (e.g. MOV.W #xx: 16, SP).

3.3 Interrupts

3.3.1 Overview

The interrupt sources include 11 external interrupts (WKP₇ to WKP₀, IRQ_1 to IRQ_0 , IRQAEC) and 7 internal interrupts from on-chip peripheral modules. Table 3.2 shows the interrupt sources, their priorities, and their vector addresses. When more than one interrupt is requested, the interrupt with the highest priority is processed.

The interrupts have the following features:

- Internal and external interrupts can be masked by the I bit in CCR. When the I bit is set to 1, interrupt request flags can be set but the interrupts are not accepted.
- IRQAEC, IRQ₁ to IRQ₀, and WKP₇ to WKP₀ can be set to either rising edge sensing or falling edge sensing.

Interrupt Source	Interrupt	Vector Number	Vector Address	Priority
RES	Reset	0	H'0000 to H'0001	High
ĪRQ ₀	IRQ ₀	4	H'0008 to H'0009	- ▲
ĪRQ ₁	IRQ ₁	5	H'000A to H'000B	-
IRQAEC	IRQAEC	6	H'000C to H'000D	-
	WKP ₀ WKP ₁ WKP ₂ WKP ₃ WKP ₄ WKP ₅ WKP ₆ WKP ₇	9	H'0012 to H'0013	_
Timer A	Timer A overflow	11	H'0016 to H'0017	-
Asynchronous event counter	Asynchronous event counter overflow	12	H'0018 to H'0019	-
Timer FL	Timer FL compare match Timer FL overflow	14	H'001C to H'001D	-
Timer FH	Timer FH compare match Timer FH overflow	15	H'001E to H'001F	-
SCI3	SCI3 transmit end SCI3 transmit data empty SCI3 receive data full SCI3 overrun error SCI3 framing error SCI3 parity error	18	H'0024 to H'0025	_
A/D	A/D conversion end	19	H'0026 to H'0027	- 🖌
(SLEEP instruction executed)	Direct transfer	20	H'0028 to H'0029	Low

Table 3.2 Interrupt Sources and Their Priorities

Note: Vector addresses H'0002 to H'0007, H'000E to H'0011, H'0014 to H'0015, H'001A to H'001B, and H'0020 to H'0023 are reserved and cannot be used.

3.3.2 Interrupt Control Registers

Table 3.3 lists the registers that control interrupts.

Table 3.3 Interrupt Control Registers

Name	Abbreviation	R/W	Initial Value	Address
IRQ edge select register	IEGR	R/W	—	H'FFF2
Interrupt enable register 1	IENR1	R/W	_	H'FFF3
Interrupt enable register 2	IENR2	R/W		H'FFF4
Interrupt request register 1	IRR1	R/W*	—	H'FFF6
Interrupt request register 2	IRR2	R/W*	_	H'FFF7
Wakeup interrupt request register	IWPR	R/W*	H'00	H'FFF9
Wakeup edge select register	WEGR	R/W	H'00	H'FF90

Note: * Write is enabled only for writing of 0 to clear a flag.

1. IRQ edge select register (IEGR)

Bit	7	6	5	4	3	2	1	0
	—	—	—	_	—	_	IEG1	IEG0
Initial value	1	1	1				0	0
Read/Write	_	_	_	W	W	W	R/W	R/W

IEGR is an 8-bit read/write register used to designate whether pins \overline{IRQ}_1 and \overline{IRQ}_0 are set to rising edge sensing or falling edge sensing.

Bits 7 to 5: Reserved bits

Bits 7 to 5 are reserved; they are always read as 1 and cannot be modified.

Bits 4 to 2: Reserved bits

Bits 4 to 2 are reserved; only 0 can be written to these bits.

Bit 1: IRQ₁ edge select (IEG1)

Bit 1 selects the input sensing of the \overline{IRQ}_1 pin.

Bit 1 IEG1	Description	
0	Falling edge of \overline{IRQ}_1 pin input is detected	(initial value)
1	Rising edge of \overline{IRQ}_1 pin input is detected	

Bit 0: IRQ₀ edge select (IEG0)

Bit 0 selects the input sensing of pin \overline{IRQ}_0 .

Bit 0 IEG0	Description	
0	Falling edge of \overline{IRQ}_0 pin input is detected	(initial value)
1	Rising edge of \overline{IRQ}_0 pin input is detected	

2. Interrupt enable register 1 (IENR1)

Bit	7	6	5	4	3	2	1	0
	IENTA	—	IENWP	—	_	IENEC2	IEN1	IEN0
Initial value	0	—	0	_		0	0	0
Read/Write	R/W	W	R/W	W	W	R/W	R/W	R/W

IENR1 is an 8-bit read/write register that enables or disables interrupt requests.

Bit 7: Timer A interrupt enable (IENTA)

Bit 7 enables or disables timer A overflow interrupt requests.

Bit 7 IENTA	Description	
0	Disables timer A interrupt requests	(initial value)
1	Enables timer A interrupt requests	

Bit 6: Reserved bit

Bit 6 is reserved; only 0 can be written to this bit.

Bit 5: Wakeup interrupt enable (IENWP)

Bit 5 enables or disables WKP₇ to WKP₀ interrupt requests.

Bit 5 IENWP	Description	
0	Disables \overline{WKP}_7 to \overline{WKP}_0 interrupt requests	(initial value)
1	Enables \overline{WKP}_7 to \overline{WKP}_0 interrupt requests	

Bits 4 and 3: Reserved bits

Bits 4 and 3 are reserved; only 0 can be written to these bits.

Bit 2: IRQAEC interrupt enable (IENEC2)

Bit 2 enables or disables IRQAEC interrupt requests.

Bit 2 IENEC2	Description	
0	Disables IRQAEC interrupt requests	(initial value)
1	Enables IRQAEC interrupt requests	

Bits 1 and 0: IRQ_1 and IRQ_0 interrupt enable (IEN1 and IEN0)

Bits 1 and 0 enable or disable IRQ_1 and IRQ_0 interrupt requests.

Bit n IENn	Description	
0	Disables interrupt requests from pin IRQn	(initial value)
1	Enables interrupt requests from pin IRQn	

(n = 1 or 0)

3. Interrupt enable register 2 (IENR2)

Bit	7	6	5	4	3	2	1	0
	IENDT	IENAD	—	—	IENTFH	IENTFL	—	IENEC
Initial value	0	0	_	—	0	0	_	0
Read/Write	R/W	R/W	W	W	R/W	R/W	W	R/W

IENR2 is an 8-bit read/write register that enables or disables interrupt requests.

Bit 7: Direct transfer interrupt enable (IENDT)

Bit 7 enables or disables direct transfer interrupt requests.

Bit 7 IENDT	Description	
0	Disables direct transfer interrupt requests	(initial value)
1	Enables direct transfer interrupt requests	

Bit 6: A/D converter interrupt enable (IENAD)

Bit 6 enables or disables A/D converter interrupt requests.

Bit 6 IENAD	Description	
0	Disables A/D converter interrupt requests	(initial value)
1	Enables A/D converter interrupt requests	

Bits 5 and 4: Reserved bits

Bits 5 and 4 are reserved; only 0 can be written to these bits.

Bit 3: Timer FH interrupt enable (IENTFH)

Bit 3 enables or disables timer FH compare match and overflow interrupt requests.

Bit 3 IENTFH	Description	
0	Disables timer FH interrupt requests	(initial value)
1	Enables timer FH interrupt requests	

Bit 2: Timer FL interrupt enable (IENTFL)

Bit 2 enables or disables timer FL compare match and overflow interrupt requests.

Bit 2 IENTFL	Description	
0	Disables timer FL interrupt requests	(initial value)
1	Enables timer FL interrupt requests	

Bit 1: Reserved bit

Bit 1 is reserved; only 0 can be written to this bit.

Bit 0: Asynchronous event counter interrupt enable (IENEC)

Bit 0 enables or disables asynchronous event counter interrupt requests.

Bit 0 IENEC	Description	
0	Disables asynchronous event counter interrupt requests	(initial value)
1	Enables asynchronous event counter interrupt requests	

For details of SCI3 interrupt control, see 10.2.6. Serial control register 3 (SCR3).

4. Interrupt request register 1 (IRR1)

Bit	7	6	5	4	3	2	1	0
	IRRTA	—	_	—	_	IRREC2	IRRI1	IRRI0
Initial value	0	_	1	_		0	0	0
Read/Write	R/W*	W	—	W	W	R/W*	R/W*	R/W*

Note: * Only a write of 0 for flag clearing is possible

IRR1 is an 8-bit read/write register, in which a corresponding flag is set to 1 when a timer A, IRQAEC or IRQ_1 , IRQ_0 interrupt is requested. The flags are not cleared automatically when an interrupt is accepted. It is necessary to write 0 to clear each flag.

Bit 7: Timer A interrupt request flag (IRRTA)

Bit 7 IRRTA	Description	
0	Clearing conditions: When IRRTA = 1, it is cleared by writing 0	(initial value)
1	Setting conditions: When the timer A counter value overflows from H'FF to H'00	

Bits 6, 4, and 3: Reserved bits

Bits 6, 4, and 3 are reserved; only 0 can be written to these bits.

Bit 5: Reserved bit

Bit 5 is reserved; it is always read as 1 and cannot be modified.

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Bit 2: IRQAEC interrupt request flag (IRREC2)

Bit 2 IRREC2	Description	
0	Clearing conditions: When IRREC2 = 1, it is cleared by writing 0	(initial value)
1	Setting conditions: When pin IRQAEC is designated for interrupt input and the designated s input	ignal edge is

Bits 1 and 0: IRQ_1 and IRQ_0 interrupt request flags (IRRI1 and IRRI0)

Bit n IRRIn	Description
0	Clearing conditions: (initial value) When IRRIn = 1, it is cleared by writing 0
1	Setting conditions: When pin IRQn is designated for interrupt input and the designated signal edge is input
	(n = 1 or 0

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5. Interrupt request register 2 (IRR2)

Bit	7	6	5	4	3	2	1	0
	IRRDT	IRRAD	—	—	IRRTFH	IRRTFL	—	IRREC
Initial value	0	0	—		0	0		0
Read/Write	R/W*	R/W*	W	W	R/W*	R/W *	W	R/W *

Note: * Only a write of 0 for flag clearing is possible

IRR2 is an 8-bit read/write register, in which a corresponding flag is set to 1 when a direct transfer, A/D converter, Timer FH, or Timer FL asynchronous event counter interrupt is requested. The flags are not cleared automatically when an interrupt is accepted. It is necessary to write 0 to clear each flag.

Bit 7: Direct transfer interrupt request flag (IRRDT)

Bit 7 IRRDT	Description
0	Clearing conditions: (initial value) When IRRDT = 1, it is cleared by writing 0
1	Setting conditions: When a direct transfer is made by executing a SLEEP instruction while DTON = 1 in SYSCR2

Bit 6: A/D converter interrupt request flag (IRRAD)

Bit 6 IRRAD	Description	
0	Clearing conditions: When IRRAD = 1, it is cleared by writing 0	(initial value)
1	Setting conditions: When A/D conversion is completed and ADSF is cleared to 0 in ADSR	

Bits 5 and 4: Reserved bits

Bits 5 and 4 are reserved; only 0 can be written to these bits.

Bit 3: Timer FH interrupt request flag (IRRTFH)

Bit 3 IRRTFH	Description	
0	Clearing conditions: (ini When IRRTFH = 1, it is cleared by writing 0	tial value)
1	Setting conditions: When TCFH and OCRFH match in 8-bit timer mode, or when TCF (TCFL, and OCRF (OCRFL, OCRFH) match in 16-bit timer mode	TCFH)

Bit 2: Timer FL interrupt request flag (IRRTFL)

Bit 2 IRRTFL	Description	
0	Clearing conditions: When IRRTFL= 1, it is cleared by writing 0	(initial value)
1	Setting conditions: When TCFL and OCRFL match in 8-bit timer mode	

Bit 1: Reserved bit

Bit 1 is reserved; only 0 can be written to this bit.

Bit 0: Asynchronous event counter interrupt request flag (IRREC)

Bit 0 IRREC	Description	
0	Clearing conditions: When IRREC = 1, it is cleared by writing 0	(initial value)
1	Setting conditions: When ECH overflows in 16-bit counter mode, or ECH or ECL overflows counter mode	in 8-bit

6. Wakeup Interrupt Request Register (IWPR)

Bit	7	6	5	4	3	2	1	0
	IWPF7	IWPF6	IWPF5	IWPF4	IWPF3	IWPF2	IWPF1	IWPF0
Initial value	0	0	0	0	0	0	0	0
Read/Write	R/W*	R/W*	R/W*	R/W*	R/W*	R/W *	R/W *	R/W*

Note: * Only a write of 0 for flag clearing is possible

IWPR is an 8-bit read/write register containing wakeup interrupt request flags. When one of pins \overline{WKP}_7 to \overline{WKP}_0 is designated for wakeup input and a rising or falling edge is input at that pin, the corresponding flag in IWPR is set to 1. A flag is not cleared automatically when the corresponding interrupt is accepted. Flags must be cleared by writing 0.

Bits 7 to 0: Wakeup interrupt request flags (IWPF7 to IWPF0)

Bit n IWPFn	Description
0	Clearing conditions: (initial value) When IWPFn= 1, it is cleared by writing 0
1	Setting conditions: When pin \overline{WKP}_n is designated for wakeup input and a rising or falling edge is input a that pin

(n = 7 to 0)

7. Wakeup Edge Select Register (WEGR)

Bit	7	6	5	4	3	2	1	0
	WKEGS7	WKEGS6	WKEGS5	WKEGS4	WKEGS3	WKEGS2	WKEGS1	WKEGS0
Initial value	0	0	0	0	0	0	0	0
Read/Write	R/W							

WEGR is an 8-bit read/write register that specifies rising or falling edge sensing for pins $\overline{WKP}n$.

WEGR is initialized to H'00 by a reset.

Bit n: WKPn edge select (WKEGSn)

Bit n selects $\overline{WKP}n$ pin input sensing.

Bit n WKEGSn	Description	
0	WKPn pin falling edge detected	(initial value)
1	WKPn pin rising edge detected	

(n = 7 to 0)

3.3.3 External Interrupts

There are 11 external interrupts: WKP7 to WKP0, IRQ1 to IRQ0, and IRQAEC.

1. Interrupts WKP₇ to WKP₀

Interrupts WKP₇ to WKP₀ are requested by either rising or falling edge input to pins $\overline{WKP_7}$ to $\overline{WKP_0}$. When these pins are designated as pins $\overline{WKP_7}$ to $\overline{WKP_0}$ in port mode register 5 and a rising or falling edge is input, the corresponding bit in IWPR is set to 1, requesting an interrupt. Recognition of wakeup interrupt requests can be disabled by clearing the IENWP bit to 0 in IENR1. These interrupts can all be masked by setting the I bit to 1 in CCR.

When WKP_7 to WKP_0 interrupt exception handling is initiated, the I bit is set to 1 in CCR. Vector number 9 is assigned to interrupts WKP_7 to WKP_0 . All eight interrupt sources have the same vector number, so the interrupt-handling routine must discriminate the interrupt source.

2. Interrupts IRQ_1 and IRQ_0

Interrupts IRQ_1 and IRQ_0 are requested by input signals to pins \overline{IRQ}_1 and \overline{IRQ}_0 . These interrupts are detected by either rising edge sensing or falling edge sensing, depending on the settings of bits IEG1 and IEG0 in IEGR.

When these pins are designated as pins \overline{IRQ}_1 and \overline{IRQ}_0 in port mode register B and 2 and the designated edge is input, the corresponding bit in IRR1 is set to 1, requesting an interrupt. Recognition of these interrupt requests can be disabled individually by clearing bits IEN1 and IEN0 to 0 in IENR1. These interrupts can all be masked by setting the I bit to 1 in CCR.

When IRQ_1 and IRQ_0 interrupt exception handling is initiated, the I bit is set to 1 in CCR. Vector numbers 5 and 4 are assigned to interrupts IRQ_1 and IRQ_0 . The order of priority is from IRQ_0 (high) to IRQ_1 (low). Table 3.2 gives details.

3. IRQAEC Interrupt

The IRQAEC interrupt is requested by an input signal to pin IRQAEC. This interrupt is detected by rising edge, falling edge, or both edge sensing, depending on the settings of bits AIAGS1 and AIAGS0 in AEGSR.

When bit IENEC2 in IENR1 is 1 and the designated edge is input, the corresponding bit in IRR1 is set to 1, requesting an interrupt.

When IRQAEC interrupt exception handling is initiated, the I bit is set to 1 in CCR. Vector number 6 is assigned to the IRQAEC interrupt. Table 3.2 gives details.

3.3.4 Internal Interrupts

There are 7 internal interrupts that can be requested by the on-chip peripheral modules. When a peripheral module requests an interrupt, the corresponding bit in IRR1 or IRR2 is set to 1. Recognition of individual interrupt requests can be disabled by clearing the corresponding bit in IENR1 or IENR2. All these interrupts can be masked by setting the I bit to 1 in CCR. When internal interrupt handling is initiated, the I bit is set to 1 in CCR. Vector numbers from 20 to 18, 15, 14, 12, and 11 are assigned to these interrupts. Table 3.2 shows the order of priority of interrupts from on-chip peripheral modules.

3.3.5 Interrupt Operations

Interrupts are controlled by an interrupt controller. Figure 3.2 shows a block diagram of the interrupt controller. Figure 3.3 shows the flow up to interrupt acceptance.

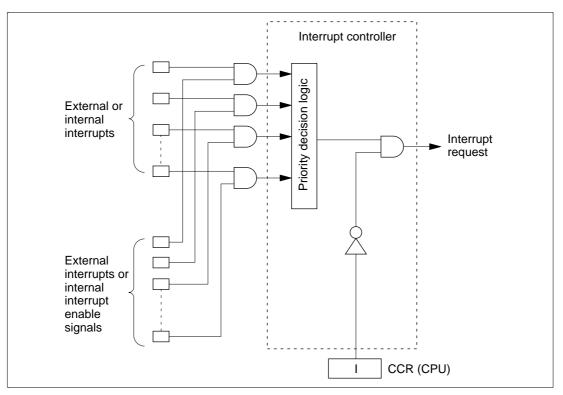


Figure 3.2 Block Diagram of Interrupt Controller

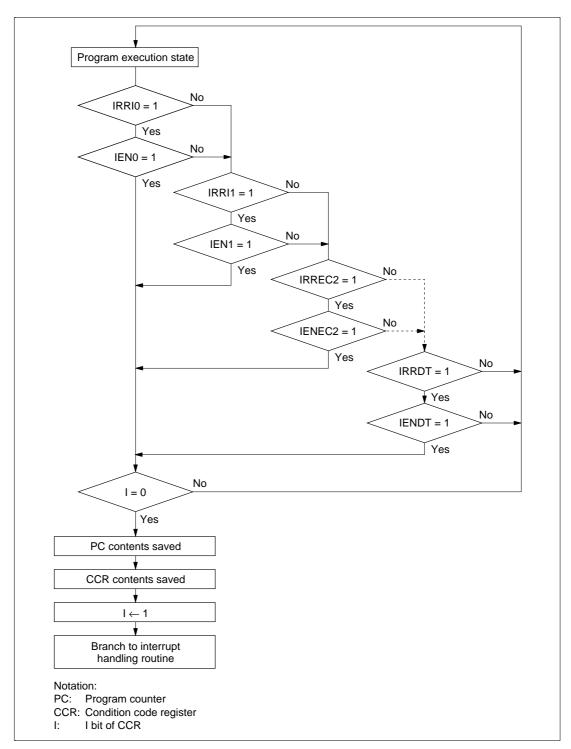
Interrupt operation is described as follows.

- When an interrupt condition is met while the interrupt enable register bit is set to 1, an interrupt request signal is sent to the interrupt controller.
- When the interrupt controller receives an interrupt request, it sets the interrupt request flag.
- From among the interrupts with interrupt request flags set to 1, the interrupt controller selects the interrupt request with the highest priority and holds the others pending. (Refer to table 3.2 for a list of interrupt priorities.)
- The interrupt controller checks the I bit of CCR. If the I bit is 0, the selected interrupt request is accepted; if the I bit is 1, the interrupt request is held pending.

- If the interrupt is accepted, after processing of the current instruction is completed, both PC and CCR are pushed onto the stack. The state of the stack at this time is shown in figure 3.4. The PC value pushed onto the stack is the address of the first instruction to be executed upon return from interrupt handling.
- The I bit of CCR is set to 1, masking further interrupts.
- The vector address corresponding to the accepted interrupt is generated, and the interrupt handling routine located at the address indicated by the contents of the vector address is executed.

Notes:

- 1. When disabling interrupts by clearing bits in an interrupt enable register, or when clearing bits in an interrupt request register, always do so while interrupts are masked (I = 1).
- 2. If the above clear operations are performed while I = 0, and as a result a conflict arises between the clear instruction and an interrupt request, exception processing for the interrupt will be executed after the clear instruction has been executed.





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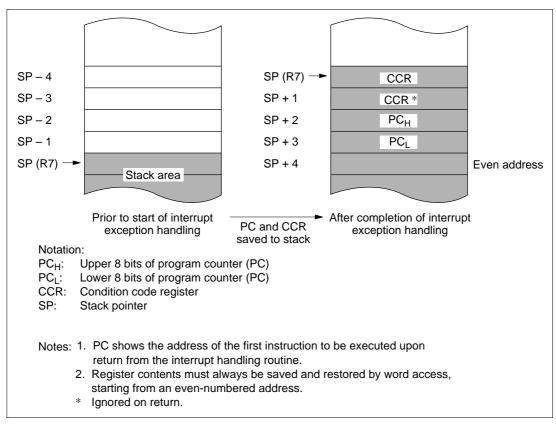


Figure 3.4 Stack State after Completion of Interrupt Exception Handling

Figure 3.5 shows a typical interrupt sequence.

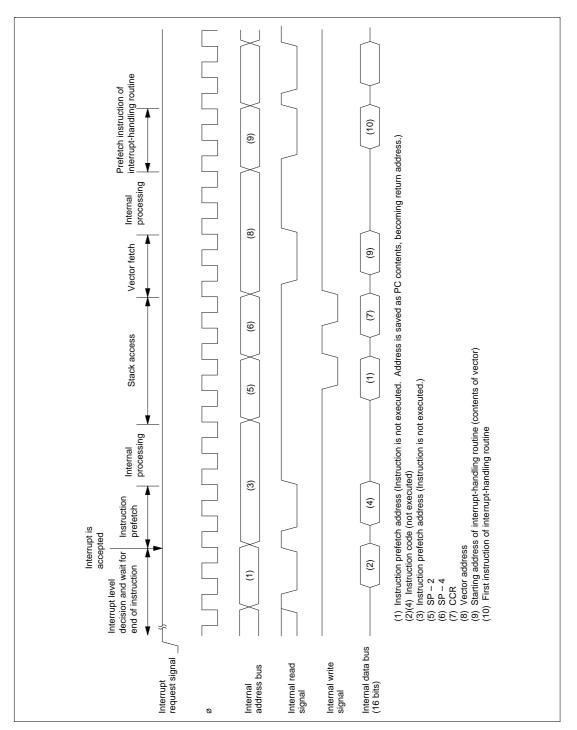


Figure 3.5 Interrupt Sequence

3.3.6 Interrupt Response Time

Table 3.4 shows the number of wait states after an interrupt request flag is set until the first instruction of the interrupt handler is executed.

Table 3.4 Interrupt Wait States

Item	States	Total
Waiting time for completion of executing instruction*	1 to 13	15 to 27
Saving of PC and CCR to stack	4	
Vector fetch	2	
Instruction fetch	4	
Internal processing	4	

Note: * Not including EEPMOV instruction.

3.4 Application Notes

3.4.1 Notes on Stack Area Use

When word data is accessed in the H8/3802 Series, the least significant bit of the address is regarded as 0. Access to the stack always takes place in word size, so the stack pointer (SP: R7) should never indicate an odd address. Use PUSH Rn (MOV.W Rn, @–SP) or POP Rn (MOV.W @SP+, Rn) to save or restore register values.

Setting an odd address in SP may cause a program to crash. An example is shown in figure 3.6.

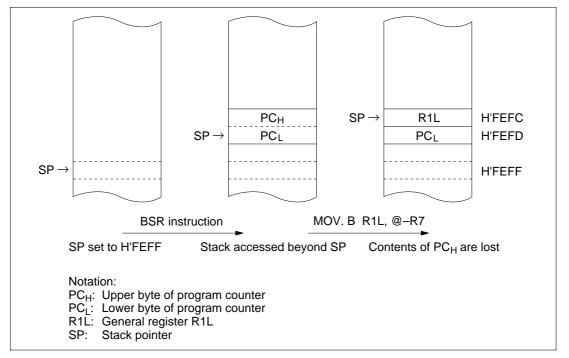


Figure 3.6 Operation when Odd Address is Set in SP

When CCR contents are saved to the stack during interrupt exception handling or restored when RTE is executed, this also takes place in word size. Both the upper and lower bytes of word data are saved to the stack; on return, the even address contents are restored to CCR while the odd address contents are ignored.

3.4.2 Notes on Rewriting Port Mode Registers

When a port mode register is rewritten to switch the functions of external interrupt pins, the following points should be observed.

When an external interrupt pin function is switched by rewriting the port mode register that controls pins IRQAEC, \overline{IRQ}_1 , \overline{IRQ}_0 , \overline{WKP}_7 to \overline{WKP}_0 , the interrupt request flag may be set to 1 at the time the pin function is switched, even if no valid interrupt is input at the pin. Be sure to clear the interrupt request flag to 0 after switching pin functions. Table 3.5 shows the conditions under which interrupt request flags are set to 1 in this way.

Interrupt Request Flags Set to 1		Conditions
IRR1 IRREC2		When the edge designated by AIEGS1 and AIEGS0 in AEGSR is input while IENEC2 in IENRI is set to 1.
	IRRI1	When PMRB bit IRQ1 is changed from 0 to 1 while pin \overline{IRQ}_1 is low and IEGR bit IEG1 = 0.
		When PMRB bit IRQ1 is changed from 1 to 0 while pin \overline{IRQ}_1 is low and IEGR bit IEG1 = 1.
	IRRI0	When PMR2 bit IRQ0 is changed from 0 to 1 while pin \overline{IRQ}_0 is low and IEGR bit IEG0 = 0.
		When PMR2 bit IRQ0 is changed from 1 to 0 while pin \overline{IRQ}_0 is low and IEGR bit IEG0 = 1.
IWPR	IWPF7	When PMR5 bit WKP7 is changed from 0 to 1 while pin $\overline{\text{WKP}}_7$ is low.
	IWPF6	When PMR5 bit WKP6 is changed from 0 to 1 while pin $\overline{\text{WKP}}_6$ is low.
	IWPF5	When PMR5 bit WKP5 is changed from 0 to 1 while pin $\overline{\text{WKP}}_5$ is low.
	IWPF4	When PMR5 bit WKP4 is changed from 0 to 1 while pin $\overline{\text{WKP}}_4$ is low.
	IWPF3	When PMR5 bit WKP3 is changed from 0 to 1 while pin $\overline{\text{WKP}}_3$ is low.
	IWPF2	When PMR5 bit WKP2 is changed from 0 to 1 while pin $\overline{\text{WKP}}_2$ is low.
	IWPF1	When PMR5 bit WKP1 is changed from 0 to 1 while pin $\overline{\text{WKP}}_1$ is low.
	IWPF0	When PMR5 bit WKP0 is changed from 0 to 1 while pin $\overline{\text{WKP}}_0$ is low.

 Table 3.5
 Conditions under which Interrupt Request Flag is Set to 1

Figure 3.7 shows the procedure for setting a bit in a port mode register and clearing the interrupt request flag.

When switching a pin function, mask the interrupt before setting the bit in the port mode register. After accessing the port mode register, execute at least one instruction (e.g., NOP), then clear the interrupt request flag from 1 to 0. If the instruction to clear the flag is executed immediately after the port mode register access without executing an intervening instruction, the flag will not be cleared.

An alternative method is to avoid the setting of interrupt request flags when pin functions are switched by keeping the pins at the high level so that the conditions in table 3.5 do not occur.

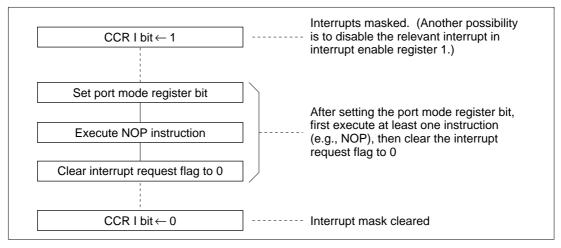


Figure 3.7 Port Mode Register Setting and Interrupt Request Flag Clearing Procedure

3.4.3 Interrupt Request Flag Clearing Method

Use the following recommended method for flag clearing in the interrupt request registers (IRR1, IRR2, and IWPR).

Recommended Method: Perform flag clearing with only one instruction. Either a bit manipulation instruction or a data transfer instruction in bytes can be used. Two examples of coding for clearing IRRI1 (bit 1 in IRR1) are shown below:

- BCLR #1,@IRR1:8
- MOV.B R1L,@IRR1:8 (Set B'11111101 into R1L in advance)

Malfunction Example: When flag clearing is performed with several instructions, a flag, other than the intended one, which was set while executing one of those instructions may be accidentally cleared, and thus cause incorrect operations to occur.

An example of coding for clearing IRR11 (bit 1 in IRR1), in which IRR10 is also cleared and the interrupt becomes invalid is shown below.

MOV.B @IRR1:8,R1L	At this point, IRRIO is O.
AND.B #B'11111101,R1L	IRRIO becomes 1 here.
MOV.B R1L,@IRR1:8	IRRIO is cleared to 0.

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In the above example, an IRQ0 interrupt occurs while the AND.B instruction is executed. Since not only the original target IRRI1, but also IRRI0 is cleared to 0, the IRQ0 interrupt becomes invalid.

Section 4 Clock Pulse Generators

4.1 Overview

Clock oscillator circuitry (CPG: clock pulse generator) is provided on-chip, including both a system clock pulse generator and a subclock pulse generator. The system clock pulse generator consists of a system clock oscillator and system clock dividers. The subclock pulse generator consists of a subclock oscillator circuit and a subclock divider.

4.1.1 Block Diagram

Figure 4.1 shows a block diagram of the clock pulse generators.

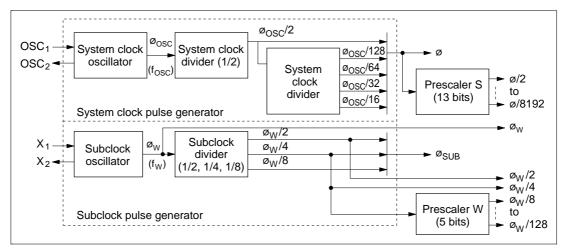


Figure 4.1 Block Diagram of Clock Pulse Generators

4.1.2 System Clock and Subclock

The basic clock signals that drive the CPU and on-chip peripheral modules are ϕ and ϕ_{SUB} . Four of the clock signals have names: ϕ is the system clock, ϕ_{SUB} is the subclock, ϕ_{OSC} is the oscillator clock, and ϕ_W is the watch clock.

The clock signals available for use by peripheral modules are $\phi/2$, $\phi/4$, $\phi/8$, $\phi/16$, $\phi/32$, $\phi/64$, $\phi/128$, $\phi/256$, $\phi/512$, $\phi/1024$, $\phi/2048$, $\phi/4096$, $\phi/8192$, ϕ_W , $\phi_W/2$, $\phi_W/4$, $\phi_W/8$, $\phi_W/16$, $\phi_W/32$, $\phi_W/64$, and $\phi_W/128$. The clock requirements differ from one module to another.

4.2 System Clock Generator

Clock pulses can be supplied to the system clock divider either by connecting a crystal or ceramic oscillator, or by providing external clock input.

1. Connecting a crystal oscillator

Figure 4.2 shows a typical method of connecting a crystal oscillator.

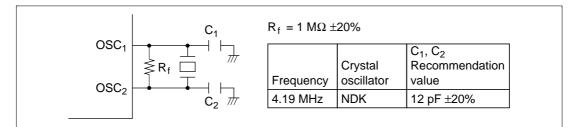


Figure 4.2 Typical Connection to Crystal Oscillator

Figure 4.3 shows the equivalent circuit of a crystal oscillator. An oscillator having the characteristics given in table 4.1 should be used.

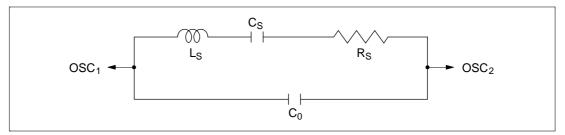


Figure 4.3 Equivalent Circuit of Crystal Oscillator

 Table 4.1
 Crystal Oscillator Parameters

Frequency (MHz)	4.193	
RS max (Ω)	100	
C ₀ max (pF)	16	

2. Connecting a ceramic oscillator

Figure 4.4 shows a typical method of connecting a ceramic oscillator.

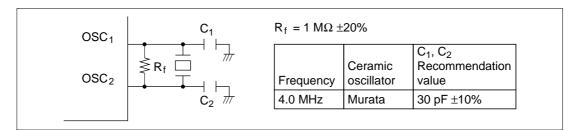


Figure 4.4 Typical Connection to Ceramic Oscillator

3. Notes on board design

When generating clock pulses by connecting a crystal or ceramic oscillator, pay careful attention to the following points.

Avoid running signal lines close to the oscillator circuit, since the oscillator may be adversely affected by induction currents. (See figure 4.5.)

The board should be designed so that the oscillator and load capacitors are located as close as possible to pins OSC_1 and OSC_2 .

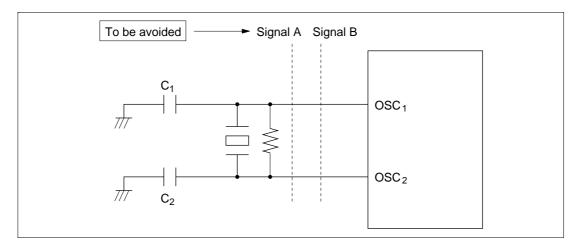


Figure 4.5 Board Design of Oscillator Circuit

4. External clock input method

Connect an external clock signal to pin OSC_1 , and leave pin OSC_2 open. Figure 4.6 shows a typical connection.

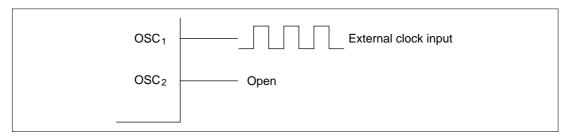


Figure 4.6 External Clock Input (Example)

Frequency	Oscillator Clock (ø _{OSC})	
Duty cycle	45% to 55%	
Note: The circuit parameters above are recommended by the crystal or ceramic oscillator		

Note: The circuit parameters above are recommended by the crystal or ceramic oscillator manufacturer.

The circuit parameters are affected by the crystal or ceramic oscillator and floating capacitance when designing the board. When using the oscillator, consult with the crystal or ceramic oscillator manufacturer to determine the circuit parameters.

4.3 Subclock Generator

1. Connecting a 32.768 kHz/38.4 kHz crystal oscillator

Clock pulses can be supplied to the subclock divider by connecting a 32.768 kHz/38.4 kHz crystal oscillator, as shown in figure 4.7. Follow the same precautions as noted under 3. notes on board design for the system clock in 4.2.

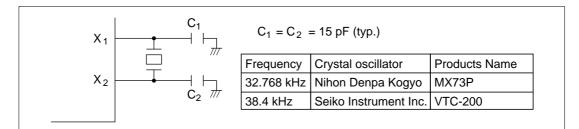


Figure 4.7 Typical Connection to 32.768 kHz/38.4 kHz Crystal Oscillator (Subclock)

Figure 4.8 shows the equivalent circuit of the 32.768 kHz/38.4 kHz crystal oscillator.

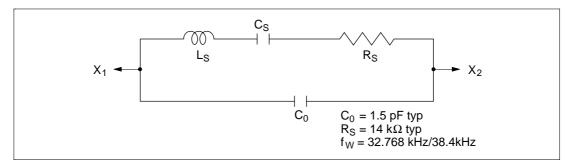


Figure 4.8 Equivalent Circuit of 32.768 kHz/38.4 kHz Crystal Oscillator

2. Pin connection when not using subclock

When the subclock is not used, connect pin X_1 to GND and leave pin X_2 open, as shown in

figure 4.9.

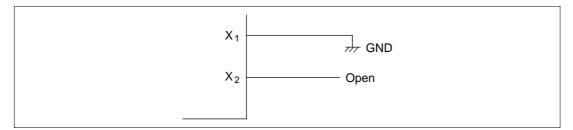


Figure 4.9 Pin Connection when not Using Subclock

3. External clock input

Connect the external clock to the X1 pin and leave the X2 pin open, as shown in figure 4.10.

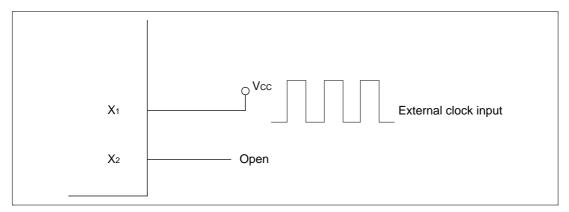


Figure 4.10 Pin Connection when Inputting External Clock

Frequency	Subclock (øw)
Duty	45% to 55%

4.4 Prescalers

The H8/3802 Series is equipped with two on-chip prescalers having different input clocks (prescaler S and prescaler W). Prescaler S is a 13-bit counter using the system clock (\emptyset) as its input clock. Its prescaled outputs provide internal clock signals for on-chip peripheral modules. Prescaler W is a 5-bit counter using a 32.768 kHz or 38.4 kHz signal divided by 4 (\emptyset_W /4) as its input clock. Its prescaled outputs are used by timer A as a time base for timekeeping.

1. Prescaler S (PSS)

Prescaler S is a 13-bit counter using the system clock (ϕ) as its input clock. It is incremented once per clock period.

Prescaler S is initialized to H'0000 by a reset, and starts counting on exit from the reset state.

In standby mode, watch mode, subactive mode, and subsleep mode, the system clock pulse generator stops. Prescaler S also stops and is initialized to H'0000.

The CPU cannot read or write prescaler S.

The output from prescaler S is shared by timer A, timer F, SCI3, the A/D converter, the LCD controller, and the 10-bit PWM. The divider ratio can be set separately for each on-chip peripheral function.

In active (medium-speed) mode the clock input to prescaler S is ϕ osc/16, ϕ osc/32, ϕ osc/64, or ϕ osc/128.

2. Prescaler W (PSW)

Prescaler W is a 5-bit counter using a 32.768 kHz/38.4 kHz signal divided by 4 ($\phi_W/4$) as its input clock.

Prescaler W is initialized to H'00 by a reset, and starts counting on exit from the reset state.

Even in standby mode, watch mode, subactive mode, or subsleep mode, prescaler W continues functioning so long as clock signals are supplied to pins X1 and X2.

Prescaler W can be reset by setting 1s in bits TMA3 and TMA2 of timer mode register A (TMA).

Output from prescaler W can be used to drive timer A, in which case timer A functions as a time base for timekeeping.

4.5 Note on Oscillators

Oscillator characteristics are closely related to board design and should be carefully evaluated by the user in mask ROM and ZTATTM versions, referring to the examples shown in this section. Oscillator circuit constants will differ depending on the oscillator element, stray capacitance in its interconnecting circuit, and other factors. Suitable constants should be determined in consultation with the oscillator element manufacturer. Design the circuit so that the oscillator element never receives voltages exceeding its maximum rating.

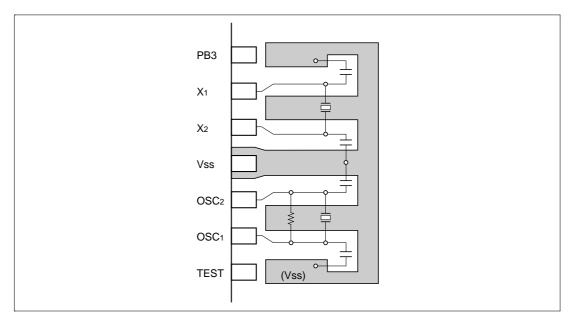


Figure 4.11 Example of Crystal and Ceramic Oscillator Element Arrangement

4.5.1 Definition of Oscillation Settling Standby Time

Figure 4.12 shows the oscillation waveform (OSC2), system clock (Ø), and microcomputer operating mode when a transition is made from standby mode, watch mode, or subactive mode, to active (high-speed/medium-speed) mode, with an oscillator element connected to the system clock oscillator.

As shown in figure 4.12, as the system clock oscillator is halted in standby mode, watch mode, and subactive mode, when a transition is made to active (high-speed/medium-speed) mode, the sum of the following two times (oscillation settling time and standby time) is required.

1. Oscillation settling time (t_{rc})

The time from the point at which the system clock oscillator oscillation waveform starts to change when an interrupt is generated, until the amplitude of the oscillation waveform increases and the oscillation frequency stabilizes.

2. Standby time

The time required for the CPU and peripheral functions to begin operating after the oscillation waveform frequency and system clock have stabilized.

The standby time setting is selected with standby timer select bits 2 to 0 (STS2 to STS0) (bits 6 to 4 in system control register 1 (SYSCR1)).

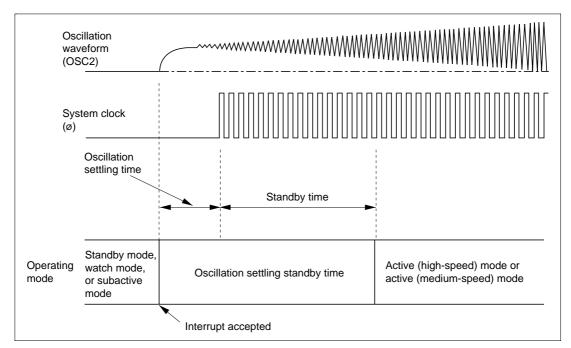


Figure 4.12 Oscillation Settling Standby Time

When standby mode, watch mode, or subactive mode is cleared by an interrupt or reset, and a transition is made to active (high-speed/medium-speed) mode, the oscillation waveform begins to change at the point at which the interrupt is accepted. Therefore, when an oscillator element is connected in standby mode, watch mode, or subactive mode, since the system clock oscillator is halted, the time from the point at which this oscillation waveform starts to change until the amplitude of the oscillation waveform increases and the oscillation frequency stabilizes—that is, the oscillation settling time—is required.

The oscillation settling time in the case of these state transitions is the same as the oscillation settling time at power-on (the time from the point at which the power supply voltage reaches the prescribed level until the oscillation stabilizes), specified by "oscillation settling time t_{rc} " in the AC characteristics.

Meanwhile, once the system clock has halted, a standby time of at least 8 states is necessary in order for the CPU and peripheral functions to operate normally.

Thus, the time required from interrupt generation until operation of the CPU and peripheral functions is the sum of the above described oscillation settling time and standby time. This total time is called the oscillation settling standby time, and is expressed by equation (1) below.

Oscillation settling standby time = oscillation settling time + standby time

 $= t_{re} + (8 \text{ to } 16,384 \text{ states})$ (1)

Therefore, when a transition is made from standby mode, watch mode, or subactive mode, to active (high-speed/medium-speed) mode, with an oscillator element connected to the system clock oscillator, careful evaluation must be carried out on the installation circuit before deciding on the oscillation settling standby time. In particular, since the oscillation settling time is affected by installation circuit constants, stray capacitance, and so forth, suitable constants should be determined in consultation with the oscillator element manufacturer.

4.5.2 Notes on Use of Crystal Oscillator Element (Excluding Ceramic Oscillator Element)

When a microcomputer operates, the internal power supply potential fluctuates slightly in synchronization with the system clock. Depending on the individual crystal oscillator element characteristics, the oscillation waveform amplitude may not be sufficiently large immediately after the oscillation settling standby time, making the oscillation waveform susceptible to influence by fluctuations in the power supply potential. In this state, the oscillation waveform may be disrupted, leading to an unstable system clock and erroneous operation of the microcomputer.

If erroneous operation occurs, change the setting of standby timer select bits 2 to 0 (STS2 to STS0) (bits 6 to 4 in system control register 1 (SYSCR1)) to give a longer standby time.

For example, if erroneous operation occurs with a standby time setting of 16 states, check the operation with a standby time setting of 1,024 states or more.

If the same kind of erroneous operation occurs after a reset as after a state transition, hold the $\overline{\text{RES}}$ pin low for a longer period.

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Section 5 Power-Down Modes

5.1 Overview

The H8/3802 Series has nine modes of operation after a reset. These include eight power-down modes, in which power dissipation is significantly reduced. Table 5.1 gives a summary of the eight operating modes.

Operating Mode	Description
Active (high-speed) mode	The CPU and all on-chip peripheral functions are operable on the system clock in high-speed operation
Active (medium-speed) mode	The CPU and all on-chip peripheral functions are operable on the system clock in low-speed operation
Subactive mode	The CPU is operable on the subclock in low-speed operation
Sleep (high-speed) mode	The CPU halts. On-chip peripheral functions are operable on the system clock
Sleep (medium-speed) mode	The CPU halts. On-chip peripheral functions operate at a frequency of 1/64, 1/32, 1/16, or 1/8 of the system clock frequency
Subsleep mode	The CPU halts. The time-base function of timer A, timer F, SCI3, AEC and LCD controller/driver are operable on the subclock
Watch mode	The CPU halts. The time-base function of timer A, timer F, AEC and LCD controller/driver are operable on the subclock
Standby mode	The CPU and all on-chip peripheral functions halt
Module standby mode	Individual on-chip peripheral functions specified by software enter standby mode and halt

Table 5.1Operating Modes

Of these nine operating modes, all but the active (high-speed) mode are power-down modes. In this section the two active modes (high-speed and medium speed) will be referred to collectively as active mode.

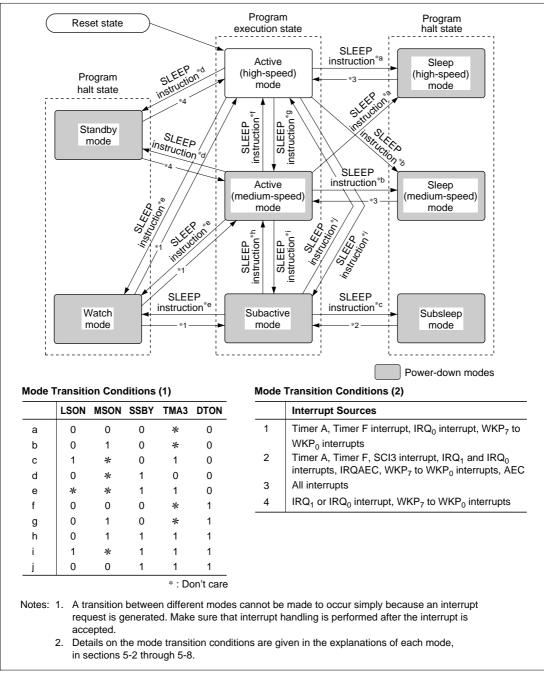


Figure 5.1 shows the transitions among these operation modes. Table 5.2 indicates the internal states in each mode.

Figure 5.1 Mode Transition Diagram

		Active	Mode	Sleep Mode					
		High-	Medium-	High-	Medium-	Watch	Subactive	Subsleep	Standby
Function		Speed	Speed	Speed	Speed	Mode	Mode	Mode	Mode
System clo	ck oscillator	Functions	Functions	Functions	Functions	Halted	Halted	Halted	Halted
Subclock of	scillator	Functions	Functions	Functions	Functions	Functions	Functions	Functions	Functions
CPU	Instructions	Functions	Functions	Halted	Halted	Halted	Functions	Halted	Halted
operations	RAM	_		Retained	Retained	Retained		Retained	Retained
	Registers	_							
	I/O ports								Retained ^{*1}
External	IRQ ₀	Functions	Functions	Functions	Functions	Functions	Functions	Functions	Functions
interrupts	IRQ ₁	_				Retained ^{*5}			
	IRQAEC								Retained ^{⁵5}
	WKP ₀	Functions	Functions	Functions	Functions	Functions	Functions	Functions	Functions
	WKP ₁	_							
	WKP ₂	_							
	WKP ₃	_							
	WKP ₄	_							
	WKP ₅	_							
	WKP ₆	_							
	WKP ₇								
Peripheral	Timer A	Functions	Functions	Functions	Functions	Functions*4	Functions*4	Functions*4	Retained
functions	Asynchronous					Functions ^{*6}	Functions	Functions	Functions*6
	counter	-					-11	11	11
	Timer F	_				Functions/ Retained*7	Functions/ Retained ^{•7}	Functions/ Retained ^{°7}	Retained
	SCI3					Reset	Functions/ Retained ^{*2}	Functions/ Retained ^{*2}	Reset
	PWM	-				Retained	Retained	Retained	Retained
	A/D converter	-				Retained	Retained	Retained	Retained
	LCD	-				Functions/ Retained ^{*3}	Functions/ Retained ^{*3}	Functions/ Retained ^{*3}	Retained

Table 5.2 Internal State in Each Operating Mode

Notes: 1. Register contents are retained, but output is high-impedance state.

2. Functions if $\phi_W/2$ is selected as the internal clock; otherwise halted and retained.

3. Functions if ϕ_W , $\phi_W/2$ or $\phi_W/4$ is selected as the operating clock; otherwise halted and retained.

4. Functions if the timekeeping time-base function is selected.

5. External interrupt requests are ignored. Interrupt request register contents are not altered.

6. Incrementing is possible, but interrupt generation is not.

7. Functions if the $\omega_W/4$ internal clock is selected; otherwise halted and retained.

5.1.1 System Control Registers

The operation mode is selected using the system control registers described in table 5.3.

Table 5.3 System Control Registers

Name	Abbreviation	R/W	Initial Value	Address
System control register 1	SYSCR1	R/W	H'07	H'FFF0
System control register 2	SYSCR2	R/W	H'F0	H'FFF1

1. System control register 1 (SYSCR1)

Bit	7	6	5	4	3	2	1	0
	SSBY	STS2	STS1	STS0	LSON	—	MA1	MA0
Initial value	0	0	0	0	0	1	1	1
Read/Write	R/W	R/W	R/W	R/W	R/W	—	R/W	R/W

SYSCR1 is an 8-bit read/write register for control of the power-down modes.

Upon reset, SYSCR1 is initialized to H'07.

Bit 7: Software standby (SSBY)

This bit designates transition to standby mode or watch mode.

Bit 7 SSBY	Description
0	• When a SLEEP instruction is executed in active mode, (initial value) a transition is made to sleep mode
	When a SLEEP instruction is executed in subactive mode, a transition is made to subsleep mode
1	When a SLEEP instruction is executed in active mode, a transition is made to standby mode or watch mode
	When a SLEEP instruction is executed in subactive mode, a transition is made to watch mode

Bits 6 to 4: Standby timer select 2 to 0 (STS2 to STS0)

These bits designate the time the CPU and peripheral modules wait for stable clock operation after exiting from standby mode or watch mode to active mode due to an interrupt. The designation should be made according to the operating frequency so that the waiting time is at least equal to the oscillation settling time.

Bit 6 STS2	Bit 5 STS1	Bit 4 STS0	Description	
0	0	0	Wait time = 8,192 states	(initial value)
0	0	1	Wait time = 16,384 states	
0	1	0	Wait time = 1,024 states	
0	1	1	Wait time = 2,048 states	
1	0	0	Wait time = 4,096 states	
1	0	1	Wait time = 2 states	(External clock mode)
1	1	0	Wait time = 8 states	
1	1	1	Wait time = 16 states	

Note: In the case that external clock is input, set up the "Standby timer select" selection to External clock mode before Mode Transition. Also, do not set up to external clock mode, in the case that it does not use external clock.

Bit 3: Low speed on flag (LSON)

This bit chooses the system clock (\emptyset) or subclock (\emptyset_{SUB}) as the CPU operating clock when watch mode is cleared. The resulting operation mode depends on the combination of other control bits and interrupt input.

Bit 3 LSON	Description	
0	The CPU operates on the system clock (ø)	(initial value)
1	The CPU operates on the subclock (øSUB)	

Bits 2: Reserved bit

Bit 2 is reserved: it is always read as 1 and cannot be modified.

Bits 1 and 0: Active (medium-speed) mode clock select (MA1, MA0)

Bits 1 and 0 choose $\phi_{osc}/128$, $\phi_{osc}/64$, $\phi_{osc}/32$, or $\phi_{osc}/16$ as the operating clock in active (medium-speed) mode and sleep (medium-speed) mode. MA1 and MA0 should be written in active (high-speed) mode or subactive mode.

Bit 1	Bit 0		
MA1	MA0	Description	
0	0	ø _{osc} /16	
0	1	ø _{osc} /32	
1	0	ø _{osc} /64	
1	1	ø _{oso} /128	(initial value)

2. System control register 2 (SYSCR2)

Bit	7	6	5	4	3	2	1	0
	_		—	NESEL	DTON	MSON	SA1	SA0
Initial value	1	1	1	1	0	0	0	0
Read/Write	_	_	_	R/W	R/W	R/W	R/W	R/W

SYSCR2 is an 8-bit read/write register for power-down mode control.

Bits 7 to 5: Reserved bits

These bits are reserved; they are always read as 1, and cannot be modified.

Bit 4: Noise elimination sampling frequency select (NESEL)

This bit selects the frequency at which the watch clock signal (ϕ_W) generated by the subclock pulse generator is sampled, in relation to the oscillator clock (ϕ_{OSC}) generated by the system clock pulse generator. When $\phi_{OSC} = 2$ to 16 MHz, clear NESEL to 0.

Bit 4 NESEL	Description	
0	Sampling rate is ø _{OSC} /16	
1	Sampling rate is ø _{OSC} /4	(initial value)

Bit 3: Direct transfer on flag (DTON)

This bit designates whether or not to make direct transitions among active (high-speed), active (medium-speed) and subactive mode when a SLEEP instruction is executed. The mode to which the transition is made after the SLEEP instruction is executed depends on a combination of this and other control bits.

Bit 3 DTON	Description
0	 When a SLEEP instruction is executed in active mode, (initial value) a transition is made to standby mode, watch mode, or sleep mode When a SLEEP instruction is executed in subactive mode, a transition is made to watch mode or subsleep mode
1	 When a SLEEP instruction is executed in active (high-speed) mode, a direct transition is made to active (medium-speed) mode if SSBY = 0, MSON = 1, and LSON = 0, or to subactive mode if SSBY = 1, TMA3 = 1, and LSON = 1 When a SLEEP instruction is executed in active (medium-speed) mode, a direct transition is made to active (high-speed) mode if SSBY = 0, MSON = 0, and LSON = 0, or to subactive mode if SSBY = 1, TMA3 = 1, and LSON = 1 When a SLEEP instruction is executed in subactive mode, a direct transition is made to active (high-speed) mode if SSBY = 1, TMA3 = 1, and LSON = 1 When a SLEEP instruction is executed in subactive mode, a direct transition is made to active (high-speed) mode if SSBY = 1, TMA3 = 1, LSON = 0, and MSON = 0, or to active (medium-speed) mode if SSBY = 1, TMA3 = 1, LSON = 0, and MSON = 1

Bit 2: Medium speed on flag (MSON)

After standby, watch, or sleep mode is cleared, this bit selects active (high-speed) or active (medium-speed) mode.

Bit 2 MSON	Description	
0	Operation in active (high-speed) mode	(initial value)
1	Operation in active (medium-speed) mode	

Bits 1 and 0: Subactive mode clock select (SA1 and SA0)

These bits select the CPU clock rate ($\phi_W/2$, $\phi_W/4$, or $\phi_W/8$) in subactive mode. SA1 and SA0 cannot be modified in subactive mode.

Bit 1 SA1	Bit 0 SA0	Description	
0	0	ø _W /8	(initial value)
0	1	ø _W /4	
1	*	ø _W /2	
			* B *

* : Don't care

5.2 Sleep Mode

5.2.1 Transition to Sleep Mode

1. Transition to sleep (high-speed) mode

The system goes from active mode to sleep (high-speed) mode when a SLEEP instruction is executed while the SSBY and LSON bits in SYSCR1 are cleared to 0, the MSON and DTON bits in SYSCR2 are cleared to 0. In sleep mode CPU operation is halted but the on-chip peripheral functions. CPU register contents are retained.

2. Transition to sleep (medium-speed) mode

The system goes from active mode to sleep (medium-speed) mode when a SLEEP instruction is executed while the SSBY and LSON bits in SYSCR1 are cleared to 0, the MSON bit in SYSCR2 is set to 1, and the DTON bit in SYSCR2 is cleared to 0. In sleep (medium-speed) mode, as in sleep (high-speed) mode, CPU operation is halted but the on-chip peripheral functions are operational. The clock frequency in sleep (medium-speed) mode is determined by the MA1 and MA0 bits in SYSCR1. CPU register contents are retained.

Furthermore, it sometimes acts with half state early timing at the time of transition to sleep (medium-speed) mode.

5.2.2 Clearing Sleep Mode

Sleep mode is cleared by any interrupt (timer A, timer F, asynchronous counter, IRQAEC, IRQ₁, IRQ₀, WKP₇ to WKP₀, SCI3, A/D converter), or by input at the $\overline{\text{RES}}$ pin.

• Clearing by interrupt

When an interrupt is requested, sleep mode is cleared and interrupt exception handling starts. A transition is made from sleep (high-speed) mode to active (high-speed) mode, or from sleep (medium-speed) mode to active (medium-speed) mode. Sleep mode is not cleared if the I bit of the condition code register (CCR) is set to 1 or the particular interrupt is disabled in the interrupt enable register.

Interrupt signal and system clock are mutually asynchronous. Synchronization error time in a maximum is $2/\phi$ (s).

• Clearing by $\overline{\text{RES}}$ input

When the $\overline{\text{RES}}$ pin goes low, the CPU goes into the reset state and sleep mode is cleared.

5.2.3 Clock Frequency in Sleep (Medium-Speed) Mode

Operation in sleep (medium-speed) mode is clocked at the frequency designated by the MA1 and MA0 bits in SYSCR1.

5.3 Standby Mode

5.3.1 Transition to Standby Mode

The system goes from active mode to standby mode when a SLEEP instruction is executed while the SSBY bit in SYSCR1 is set to 1, the LSON bit in SYSCR1 is cleared to 0, and bit TMA3 in TMA is cleared to 0. In standby mode the clock pulse generator stops, so the CPU and on-chip peripheral modules stop functioning, but as long as the rated voltage is supplied, the contents of CPU registers, on-chip RAM, and some on-chip peripheral module registers are retained. On-chip RAM contents will be further retained down to a minimum RAM data retention voltage. The I/O ports go to the high-impedance state.

5.3.2 Clearing Standby Mode

Standby mode is cleared by an interrupt (IRQ₁ or IRQ₀), WKP₇ to WKP₀ or by input at the $\overline{\text{RES}}$ pin.

• Clearing by interrupt

When an interrupt is requested, the system clock pulse generator starts. After the time set in bits STS2 to STS0 in SYSCR1 has elapsed, a stable system clock signal is supplied to the entire chip, standby mode is cleared, and interrupt exception handling starts. Operation resumes in active (high-speed) mode if MSON = 0 in SYSCR2, or active (medium-speed) mode if MSON = 1. Standby mode is not cleared if the I bit of CCR is set to 1 or the particular interrupt is disabled in the interrupt enable register.

• Clearing by $\overline{\text{RES}}$ input

When the $\overline{\text{RES}}$ pin goes low, the system clock pulse generator starts. After the pulse generator output has stabilized, if the $\overline{\text{RES}}$ pin is driven high, the CPU starts reset exception handling. Since system clock signals are supplied to the entire chip as soon as the system clock pulse generator starts functioning, the $\overline{\text{RES}}$ pin should be kept at the low level until the pulse generator output stabilizes.

5.3.3 Oscillator Settling Time after Standby Mode is Cleared

Bits STS2 to STS0 in SYSCR1 should be set as follows.

• When a crystal oscillator is used

The table below gives settings for various operating frequencies. Set bits STS2 to STS0 for a waiting time at least as long as the oscillation settling time.

STS2	STS1	STS0	Waiting Time	5 MHz	2 MHz
0	0	0	8,192 states	1.638	4.1
0	0	1	16,384 states	3.277	8.2
0	1	0	1,024 states	0.205	0.512
0	1	1	2,048 states	0.410	1.024
1	0	0	4,096 states	0.819	2.048
1	0	1	2 states (Use prohibited)	0.0004	0.001
1	1	0	8 states	0.0002	0.004
1	1	1	16 states	0.003	0.008

 Table 5.4
 Clock Frequency and Settling Time (times are in ms)

• When an external clock is used

STS2 = 1, STS1 = 0, and STS0 = 1 should be set. Other values possible use, but CPU sometimes will start operation before waiting time completion.

5.3.4 Standby Mode Transition and Pin States

When a SLEEP instruction is executed in active (high-speed) mode or active (medium-speed) mode while bit SSBY is set to 1 and bit LSON is cleared to 0 in SYSCR1, and bit TMA3 is cleared to 0 in TMA, a transition is made to standby mode. At the same time, pins go to the high-impedance state (except pins for which the pull-up MOS is designated as on). Figure 5.2 shows the timing in this case.

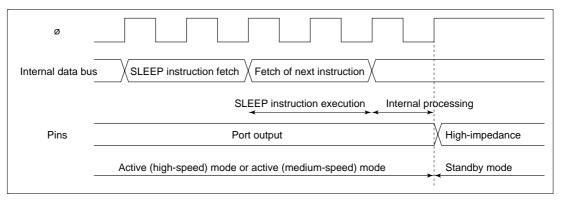


Figure 5.2 Standby Mode Transition and Pin States

5.3.5 Notes on External Input Signal Changes before/after Standby Mode

- When external input signal changes before/after standby mode or watch mode
 When an external input signal such as IRQ, WKP, or IRQAEC is input, both the high- and
 low-level widths of the signal must be at least two cycles of system clock ø or subclock ø_{SUB}
 (referred to together in this section as the internal clock). As the internal clock stops in
 standby mode and watch mode, the width of external input signals requires careful attention
 when a transition is made via these operating modes. Ensure that external input signals
 conform to the conditions stated in 3, Recommended timing of external input signals, below
- 2. When external input signals cannot be captured because internal clock stops The case of falling edge capture is illustrated in figure 5.3

As shown in the case marked "Capture not possible," when an external input signal falls immediately after a transition to active (high-speed or medium-speed) mode or subactive mode, after oscillation is started by an interrupt via a different signal, the external input signal cannot be captured if the high-level width at that point is less than 2 t_{eve} or 2 t_{subeve}.

3. Recommended timing of external input signals

To ensure dependable capture of an external input signal, high- and low-level signal widths of at least 2 t_{cyc} or 2 t_{subcyc} are necessary before a transition is made to standby mode or watch mode, as shown in "Capture possible: case 1."

External input signal capture is also possible with the timing shown in "Capture possible: case 2" and "Capture possible: case 3," in which a 2 t_{cyc} or 2 t_{subcyc} level width is secured.

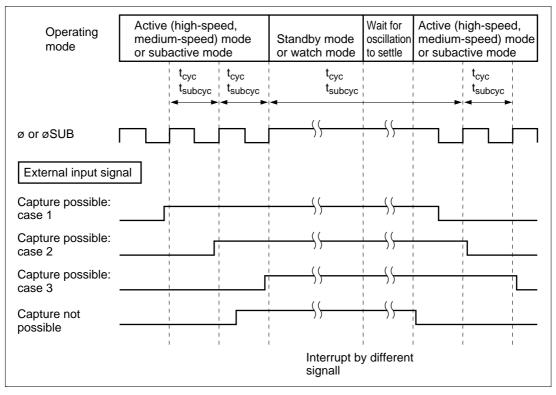


Figure 5.3 External Input Signal Capture when Signal Changes before/after Standby Mode or Watch Mode

4. Input pins to which these notes apply: \overline{IRQ}_1 to \overline{IRQ}_0 , \overline{WKP}_7 to \overline{WKP}_0 , IRQAEC

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5.4 Watch Mode

5.4.1 Transition to Watch Mode

The system goes from active or subactive mode to watch mode when a SLEEP instruction is executed while the SSBY bit in SYSCR1 is set to 1 and bit TMA3 in TMA is set to 1.

In watch mode, operation of on-chip peripheral modules is halted except for timer A, timer F, AEC and the LCD controller/driver (for which operation or halting can be set) is halted. As long as a minimum required voltage is applied, the contents of CPU registers, the on-chip RAM and some registers of the on-chip peripheral modules, are retained. I/O ports keep the same states as before the transition.

5.4.2 Clearing Watch Mode

Watch mode is cleared by an interrupt (timer A, timer F, IRQ_0 , or WKP_7 to WKP_0) or by input at the \overline{RES} pin.

• Clearing by interrupt

When watch mode is cleared by interrupt, the mode to which a transition is made depends on the settings of LSON in SYSCR1 and MSON in SYSCR2. If both LSON and MSON are cleared to 0, transition is to active (high-speed) mode; if LSON = 0 and MSON = 1, transition is to active (medium-speed) mode; if LSON = 1, transition is to subactive mode. When the transition is to active mode, after the time set in SYSCR1 bits STS2 to STS0 has elapsed, a stable clock signal is supplied to the entire chip, watch mode is cleared, and interrupt exception handling starts. Watch mode is not cleared if the I bit of CCR is set to 1 or the particular interrupt is disabled in the interrupt enable register.

• Clearing by $\overline{\text{RES}}$ input

Clearing by $\overline{\text{RES}}$ pin is the same as for standby mode; see 2. Clearing by $\overline{\text{RES}}$ pin in 5.3.2, Clearing Standby Mode.

5.4.3 Oscillator Settling Time after Watch Mode is Cleared

The waiting time is the same as for standby mode; see 5.3.3, Oscillator Settling Time after Standby Mode is Cleared.

5.4.4 Notes on External Input Signal Changes before/after Watch Mode

See 5.3.5, Notes on External Input Signal Changes before/after Standby Mode.

5.5 Subsleep Mode

5.5.1 Transition to Subsleep Mode

The system goes from subactive mode to subsleep mode when a SLEEP instruction is executed while the SSBY bit in SYSCR1 is cleared to 0, LSON bit in SYSCR1 is set to 1, and TMA3 bit in TMA is set to 1. In subsleep mode, operation of on-chip peripheral modules other than the A/D converter, and PWM is halted. As long as a minimum required voltage is applied, the contents of CPU registers, the on-chip RAM and some registers of the on-chip peripheral modules are retained. I/O ports keep the same states as before the transition.

5.5.2 Clearing Subsleep Mode

Subsleep mode is cleared by an interrupt (timer A, timer F, asynchronous counter, SCI3, IRQAEC, IRQ_1 , IRQ_0 , WKP_7 to WKP_0) or by a low input at the \overline{RES} pin.

• Clearing by interrupt

When an interrupt is requested, subsleep mode is cleared and interrupt exception handling starts. Subsleep mode is not cleared if the I bit of CCR is set to 1 or the particular interrupt is disabled in the interrupt enable register.

Interrupt signal and system clock are mutually asynchronous. Synchronization error time in a maximum is $2/\phi_{SUB}$ (s).

• Clearing by $\overline{\text{RES}}$ input

Clearing by $\overline{\text{RES}}$ pin is the same as for standby mode; see Clearing by $\overline{\text{RES}}$ pin in 5.3.2, Clearing Standby Mode.

5.6 Subactive Mode

5.6.1 Transition to Subactive Mode

Subactive mode is entered from watch mode if a timer A, timer F, IRQ_0 , or WKP_7 to WKP_0 interrupt is requested while the LSON bit in SYSCR1 is set to 1. From subsleep mode, subactive mode is entered if a timer A, timer F, asynchronous event counter, SCI3, IRQAEC, IRQ_1 , IRQ_0 , or WKP_7 to WKP_0 interrupt is requested. A transition to subactive mode does not take place if the I bit of CCR is set to 1 or the particular interrupt is disabled in the interrupt enable register.

5.6.2 Clearing Subactive Mode

Subactive mode is cleared by a SLEEP instruction or by a low input at the $\overline{\text{RES}}$ pin.

• Clearing by SLEEP instruction

If a SLEEP instruction is executed while the SSBY bit in SYSCR1 is set to 1 and TMA3 bit in TMA is set to 1, subactive mode is cleared and watch mode is entered. If a SLEEP instruction is executed while SSBY = 0 and LSON = 1 in SYSCR1 and TMA3 = 1 in TMA, subsleep mode is entered. Direct transfer to active mode is also possible; see 5.8, Direct Transfer, below.

• Clearing by $\overline{\text{RES}}$ pin

Clearing by $\overline{\text{RES}}$ pin is the same as for standby mode; see Clearing by $\overline{\text{RES}}$ pin in 5.3.2, Clearing Standby Mode.

5.6.3 Operating Frequency in Subactive Mode

The operating frequency in subactive mode is set in bits SA1 and SA0 in SYSCR2. The choices are $\phi_W/2$, $\phi_W/4$, and $\phi_W/8$.

5.7 Active (Medium-Speed) Mode

5.7.1 Transition to Active (Medium-Speed) Mode

If the $\overline{\text{RES}}$ pin is driven low, active (medium-speed) mode is entered. If the LSON bit in SYSCR2 is set to 1 while the LSON bit in SYSCR1 is cleared to 0, a transition to active (medium-speed) mode results from IRQ₀, IRQ₁ or WKP₇ to WKP₀ interrupts in standby mode, timer A, timer F, IRQ₀, or WKP₇ to WKP₀ interrupts in watch mode, or any interrupt in sleep mode. A transition to active (medium-speed) mode does not take place if the I bit of CCR is set to 1 or the particular interrupt is disabled in the interrupt enable register.

Furthermore, it sometimes acts with half state early timing at the time of transition to active (medium-speed) mode.

5.7.2 Clearing Active (Medium-Speed) Mode

Active (medium-speed) mode is cleared by a SLEEP instruction.

• Clearing by SLEEP instruction

A transition to standby mode takes place if the SLEEP instruction is executed while the SSBY bit in SYSCR1 is set to 1, the LSON bit in SYSCR1 is cleared to 0, and the TMA3 bit in TMA is cleared to 0. The system goes to watch mode if the SSBY bit in SYSCR1 is set to 1 and bit TMA3 in TMA is set to 1 when a SLEEP instruction is executed.

When both SSBY and LSON are cleared to 0 in SYSCR1 and a SLEEP instruction is executed, sleep mode is entered. Direct transfer to active (high-speed) mode or to subactive mode is also possible. See 5.8, Direct Transfer, below for details.

• Clearing by $\overline{\text{RES}}$ pin

When the $\overline{\text{RES}}$ pin is driven low, a transition is made to the reset state and active (medium-speed) mode is cleared.

5.7.3 Operating Frequency in Active (Medium-Speed) Mode

Operation in active (medium-speed) mode is clocked at the frequency designated by the MA1 and MA0 bits in SYSCR1.

5.8 Direct Transfer

5.8.1 Overview of Direct Transfer

The CPU can execute programs in three modes: active (high-speed) mode, active (medium-speed) mode, and subactive mode. A direct transfer is a transition among these three modes without the stopping of program execution. A direct transfer can be made by executing a SLEEP instruction while the DTON bit in SYSCR2 is set to 1. After the mode transition, direct transfer interrupt exception handling starts.

If the direct transfer interrupt is disabled in interrupt enable register 2, a transition is made instead to sleep mode or watch mode. Note that if a direct transition is attempted while the I bit in CCR is set to 1, sleep mode or watch mode will be entered, and it will be impossible to clear the resulting mode by means of an interrupt.

• Direct transfer from active (high-speed) mode to active (medium-speed) mode

When a SLEEP instruction is executed in active (high-speed) mode while the SSBY and LSON bits in SYSCR1 are cleared to 0, the MSON bit in SYSCR2 is set to 1, and the DTON bit in SYSCR2 is set to 1, a transition is made to active (medium-speed) mode via sleep mode.

• Direct transfer from active (medium-speed) mode to active (high-speed) mode

When a SLEEP instruction is executed in active (medium-speed) mode while the SSBY and LSON bits in SYSCR1 are cleared to 0, the MSON bit in SYSCR2 is cleared to 0, and the DTON bit in SYSCR2 is set to 1, a transition is made to active (high-speed) mode via sleep mode.

• Direct transfer from active (high-speed) mode to subactive mode

When a SLEEP instruction is executed in active (high-speed) mode while the SSBY and LSON bits in SYSCR1 are set to 1, the DTON bit in SYSCR2 is set to 1, and the TMA3 bit in TMA is set to 1, a transition is made to subactive mode via watch mode.

• Direct transfer from subactive mode to active (high-speed) mode

When a SLEEP instruction is executed in subactive mode while the SSBY bit in SYSCR1 is set to 1, the LSON bit in SYSCR1 is cleared to 0, the MSON bit in SYSCR2 is cleared to 0, the DTON bit in SYSCR2 is set to 1, and the TMA3 bit in TMA is set to 1, a transition is made directly to active (high-speed) mode via watch mode after the waiting time set in SYSCR1 bits STS2 to STS0 has elapsed.

• Direct transfer from active (medium-speed) mode to subactive mode

When a SLEEP instruction is executed in active (medium-speed) while the SSBY and LSON bits in SYSCR1 are set to 1, the DTON bit in SYSCR2 is set to 1, and the TMA3 bit in TMA is set to 1, a transition is made to subactive mode via watch mode.

• Direct transfer from subactive mode to active (medium-speed) mode

When a SLEEP instruction is executed in subactive mode while the SSBY bit in SYSCR1 is set to 1, the LSON bit in SYSCR1 is cleared to 0, the MSON bit in SYSCR2 is set to 1, the DTON bit in SYSCR2 is set to 1, and the TMA3 bit in TMA is set to 1, a transition is made directly to active (medium-speed) mode via watch mode after the waiting time set in SYSCR1 bits STS2 to STS0 has elapsed.

5.8.2 Direct Transition Times

1. Time for direct transition from active (high-speed) mode to active (medium-speed) mode

A direct transition from active (high-speed) mode to active (medium-speed) mode is performed by executing a SLEEP instruction in active (high-speed) mode while bits SSBY and LSON are both cleared to 0 in SYSCR1, and bits MSON and DTON are both set to 1 in SYSCR2. The time from execution of the SLEEP instruction to the end of interrupt exception handling (the direct transition time) is given by equation (1) below.

Direct transition time = { (Number of SLEEP instruction execution states) + (number of internal processing states) } × (tcyc before transition) + (number of interrupt exception handling execution states) × (tcyc after transition)

......(1)

Example: Direct transition time = $(2 + 1) \times 2 \text{tosc} + 14 \times 16 \text{tosc} = 230 \text{tosc}$ (when $\emptyset/8$ is selected as the CPU operating clock)

Notation:

tosc: OSC clock cycle time tcyc: System clock (ø) cycle time

2. Time for direct transition from active (medium-speed) mode to active (high-speed) mode

A direct transition from active (medium-speed) mode to active (high-speed) mode is performed by executing a SLEEP instruction in active (medium-speed) mode while bits SSBY and LSON are both cleared to 0 in SYSCR1, and bit MSON is cleared to 0 and bit DTON is set to 1 in SYSCR2. The time from execution of the SLEEP instruction to the end of interrupt exception handling (the direct transition time) is given by equation (2) below.

Direct transition time = { (Number of SLEEP instruction execution states) + (number of internal processing states) } × (tcyc before transition) + (number of interrupt exception handling execution states) × (tcyc after transition)

Example: Direct transition time = $(2 + 1) \times 16$ tosc + 14×2 tosc = 76tosc (when $\emptyset/8$ is selected as the CPU operating clock)

Notation:

tosc: OSC clock cycle time

tcyc: System clock (ø) cycle time

3. Time for direct transition from subactive mode to active (high-speed) mode

A direct transition from subactive mode to active (high-speed) mode is performed by executing a SLEEP instruction in subactive mode while bit SSBY is set to 1 and bit LSON is cleared to 0 in SYSCR1, bit MSON is cleared to 0 and bit DTON is set to 1 in SYSCR2, and bit TMA3 is set to 1 in TMA. The time from execution of the SLEEP instruction to the end of interrupt exception handling (the direct transition time) is given by equation (3) below.

Direct transition time = { (Number of SLEEP instruction execution states) + (number of internal
processing states) $\} \times ($ tsubcyc before transition $) + \{$ (wait time set in
STS2 to STS0) + (number of interrupt exception handling execution
states) $\} \times (tcyc after transition)$ (3)

Example: Direct transition time = $(2 + 1) \times 8tw + (8192 + 14) \times 2tosc = 24tw + 16412tosc$ (when øw/8 is selected as the CPU operating clock, and wait time = 8192 states)

Notation:

tosc:	OSC clock cycle time
tw:	Watch clock cycle time
tcyc:	System clock (ø) cycle time
tsubcyc:	Subclock (øSUB) cycle time

4. Time for direct transition from subactive mode to active (medium-speed) mode

A direct transition from subactive mode to active (medium-speed) mode is performed by executing a SLEEP instruction in subactive mode while bit SSBY is set to 1 and bit LSON is cleared to 0 in SYSCR1, bits MSON and DTON are both set to 1 in SYSCR2, and bit TMA3 is set to 1 in TMA. The time from execution of the SLEEP instruction to the end of interrupt exception handling (the direct transition time) is given by equation (4) below.

Direct transition time = { (Number of SLEEP instruction execution states) + (number of internal
processing states) $\} \times (tsubcyc before transition) + \{ (wait time set in the set in th$
STS2 to STS0) + (number of interrupt exception handling execution
states) $\} \times (tcyc after transition)$ (4)

Example: Direct transition time = $(2 + 1) \times 8$ tw + $(8192 + 14) \times 16$ tosc = 24tw + 131296tosc (when ϕ w/8 or ϕ 8 is selected as the CPU operating clock, and wait time = 8192 states)

Notation:

tosc:	OSC clock cycle time
tw:	Watch clock cycle time
tcyc:	System clock (ø) cycle time
tsubcyc:	Subclock (øSUB) cycle time

5.8.3 Notes on External Input Signal Changes before/after Direct Transition

- Direct transition from active (high-speed) mode to subactive mode Since the mode transition is performed via watch mode, see 5.3.5, Notes on External Input Signal Changes before/after Standby Mode.
- Direct transition from active (medium-speed) mode to subactive mode Since the mode transition is performed via watch mode, see 5.3.5, Notes on External Input Signal Changes before/after Standby Mode.
- Direct transition from subactive mode to active (high-speed) mode Since the mode transition is performed via watch mode, see 5.3.5, Notes on External Input Signal Changes before/after Standby Mode.
- Direct transition from subactive mode to active (medium-speed) mode Since the mode transition is performed via watch mode, see 5.3.5, Notes on External Input Signal Changes before/after Standby Mode.

5.9 Module Standby Mode

5.9.1 Setting Module Standby Mode

Module standby mode is set for individual peripheral functions. All the on-chip peripheral modules can be placed in module standby mode. When a module enters module standby mode, the system clock supply to the module is stopped and operation of the module halts. This state is identical to standby mode.

Module standby mode is set for a particular module by setting the corresponding bit to 0 in clock stop register 1 (CKSTPR1) or clock stop register 2 (CKSTPR2). (See table 5.5.)

5.9.2 Clearing Module Standby Mode

Module standby mode is cleared for a particular module by setting the corresponding bit to 1 in clock stop register 1 (CKSTPR1) or clock stop register 2 (CKSTPR2). (See table 5.5.)

Following a reset, clock stop register 1 (CKSTPR1) and clock stop register 2 (CKSTPR2) are both initialized to H'FF.

Table 5	5.5
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Register Name	Bit Name		Operation
CKSTPR1	TACKSTP 1		Timer A module standby mode is cleared
		0	Timer A is set to module standby mode
	TFCKSTP	1	Timer F module standby mode is cleared
		0	Timer F is set to module standby mode
	ADCKSTP	1	A/D converter module standby mode is cleared
		0	A/D converter is set to module standby mode
	S32CKSTP	1	SCI3 module standby mode is cleared
		0	SCI3 is set to module standby mode

Register Name	Bit Name		Operation
CKSTPR2	LDCKSTP	1	LCD module standby mode is cleared
		0	LCD is set to module standby mode
	PW1CKSTP	1	PWM1 module standby mode is cleared
		0	PWM1 is set to module standby mode
	AECKSTP	1	Asynchronous event counter module standby mode is cleared
		0	Asynchronous event counter is set to module standby mode
	PW2CKSTP	1	PWM2 module standby mode is cleared
		0	PWM2 is set to module standby mode

Note: For details of module operation, see the sections on the individual modules.

Section 6 ROM

6.1 Overview

The H8/3802 has 16 kbytes of on-chip mask ROM, the H8/3801 has 12 kbytes, and the H8/3800 has 8 kbytes. The ROM is connected to the CPU by a 16-bit data bus, allowing high-speed two-state access for both byte data and word data. The H8/3802 has a ZTATTM version with 16-kbyte PROM.

6.1.1 Block Diagram

Figure 6.1 shows a block diagram of the on-chip ROM.

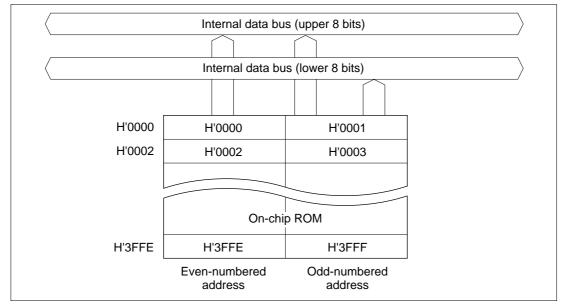


Figure 6.1 ROM Block Diagram (H8/3802)

6.2 H8/3802 PROM Mode

6.2.1 Setting to PROM Mode

If the on-chip ROM is PROM, setting the chip to PROM mode stops operation as a microcontroller and allows the PROM to be programmed in the same way as the standard HN27C101 EPROM. However, page programming is not supported. Table 6.1 shows how to set the chip to PROM mode.

Table 6.1Setting to PROM Mode

Pin Name	Setting
TEST	High level
PB ₀ /AN ₀	Low level
PB ₁ /AN ₁	—
PB ₂ /AN ₂	High level

6.2.2 Socket Adapter Pin Arrangement and Memory Map

A standard PROM programmer can be used to program the PROM. A socket adapter is required for conversion to 32 pins.

Figure 6.2 shows the pin-to-pin wiring of the socket adapter. Figure 6.3 shows a memory map.

FP-64A, FP-64E	DP-64S	Pin		Pin	HN27C101 (32-
8	16	RES			1
40	48	P60		EO0	13
39	47	P61		EO1	14
38	46	P62		EO2	15
37	45	P63		EO3	17
36	44	P64		EO4	18
35	43	P65		EO₅	19
34	42	P66		EO6	20
33	41	P67		EO7	21
57	1	P40		EA0	12
58	2	P41		EA1	11
10	18	P32		EA2	10
11	19	P33		EA3	9
12	20	P34		EA4	8
13	21	P35		EA5	7
14	22	P36		EA6	6
15	23	P37		EA7	5
32	40	P70		EA8	27
60	4	P43		EA9	26
30	38	P72		EA10	23
29	37	P73		EA11	25
28	36	P74		EA12	4
27	35	P75		EA13	28
26	34	P76		EA14	29
52	60	P93		EA15	3
53	61	P94		EA16	2
25	33	P77			22
31	39	P71			24
51	59	P92		PGM	31
16	24	Vcc		Vcc	32
61	5	AVcc			-
7	15	TEST			
2	10	X1			
64	8	PB ₂			
49	57	P90			
50	58	P91			
54	62	P95			
55	63	Vss		Vss	16
4	12	AVss			
62	6	PB ₀			
63	7	PB1			

Figure 6.2 Socket Adapter Pin Correspondence (with HN27C101)

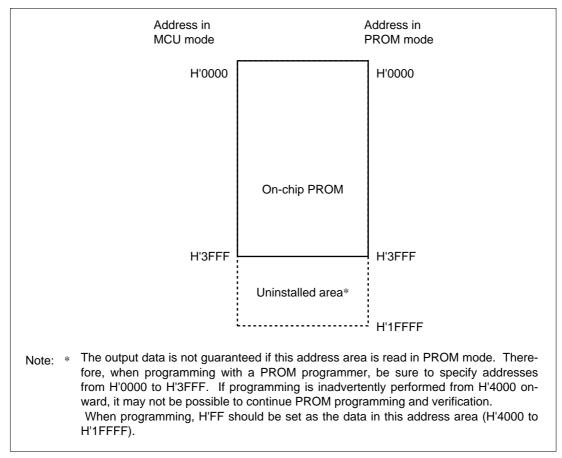


Figure 6.3 H8/3802 Memory Map in PROM Mode

6.3 H8/3802 Programming

The write, verify, and other modes are selected as shown in table 6.2 in H8/3802 PROM mode.

	Pins									
Mode	CE	ŌĒ	PGM	V _{PP}	Vcc	EO ₇ to EO ₀	EA ₁₆ to EA ₀			
Write	L	Н	L	V _{PP}	V _{CC}	Data input	Address input			
Verify	L	L	Н	V _{PP}	V _{CC}	Data output	Address input			
Programming	L	L	L	V_{PP}	V _{CC}	High impedance	Address input			
disabled	L	Н	Н							
	Н	L	L	_						
	Н	Н	Н							
Notation										
L: Low lev	el									
H: High level										
V _{PP} : V _{PP} leve	el									

Table 6.2 Mode Selection in PROM Mode (H8/3802)

V_{CC}: V_{CC} level

The specifications for writing and reading are identical to those for the standard HN27C101 EPROM. However, page programming is not supported, and so page programming mode must not be set. A PROM programmer that only supports page programming mode cannot be used. When selecting a PROM programmer, ensure that it supports high-speed, high-reliability byte-by-byte programming. Also, be sure to specify addresses from H'0000 to H'3FFF.

6.3.1 Writing and Verifying

An efficient, high-speed, high-reliability method is available for writing and verifying the PROM data. This method achieves high speed without voltage stress on the device and without lowering the reliability of written data. The basic flow of this high-speed, high-reliability programming method is shown in figure 6.4.

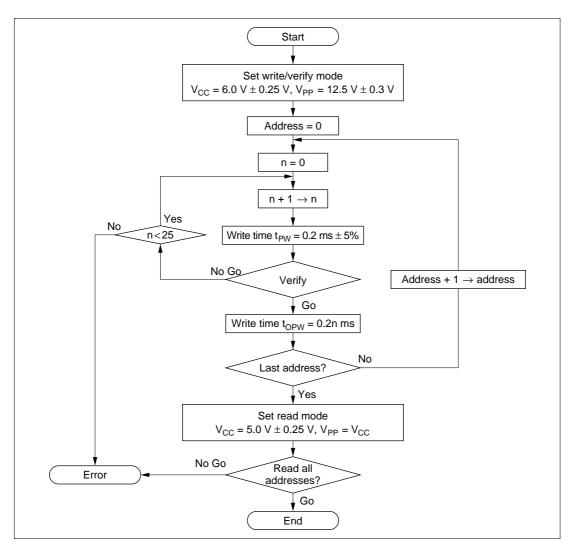


Figure 6.4 High-Speed, High-Reliability Programming Flow Chart

Table 6.3 and table 6.4 give the electrical characteristics in programming mode.

Table 6.3DC Characteristics

(Conditions: $V_{CC} = 6.0 \text{ V} \pm 0.25 \text{ V}$, $V_{PP} = 12.5 \text{ V} \pm 0.3 \text{ V}$, $V_{SS} = 0 \text{ V}$, $T_a = 25^{\circ}C \pm 5^{\circ}C$)

Item		Symbol	Min	Тур	Max	Unit	Test Condition
Input high- level voltage	$\frac{\text{EO}_7}{\text{OE}} \frac{\text{to EO}_0, \text{ EA}_{16}}{\text{OE}} \text{ to EA}_0$	V _{IH}	2.4	—	V _{CC} + 0.3	V	
Input low- level voltage	$\frac{\text{EO}_7}{\text{OE}} \frac{\text{to EO}_0, \text{EA}_{16}}{\text{OE}} \text{to EA}_0$	V _{IL}	-0.3	_	0.8	V	
Output high- level voltage	EO ₇ to EO ₀	V _{OH}	2.4		—	V	I _{OH} = -200 μA
Output low- level voltage	EO ₇ to EO ₀	V _{OL}	_		0.45	V	I _{OL} = 0.8 mA
Input leakage current	$\frac{\text{EO}_7 \text{ to } \text{EO}_0, \text{ EA}_{16} \text{ to } \text{EA}_0}{\text{OE}, \text{ CE}, \text{ PGM}}$	I _{LI}	_		2	μA	V _{in} = 5.25 V/ 0.5 V
V _{CC} current		I _{CC}	_	_	40	mA	
V _{PP} current		I _{PP}			40	mA	

Table 6.4 AC Characteristics

(Conditions: $V_{CC} = 6.0 \text{ V} \pm 0.25 \text{ V}$, $V_{PP} = 12.5 \text{ V} \pm 0.3 \text{ V}$, $T_a = 25^{\circ}C \pm 5^{\circ}C$)

Item	Symbol	Min	Тур	Мах	Unit	Test Condition
Address setup time	t _{AS}	2	—	—	μs	Figure 6.5 ^{*1}
OE setup time	t _{OES}	2	—	—	μs	
Data setup time	t _{DS}	2			μs	
Address hold time	t _{AH}	0			μs	
Data hold time	t _{DH}	2	_	_	μs	
Data output disable time	t _{DF} *2			130	ns	
V _{PP} setup time	t _{VPS}	2			μs	
Programming pulse width	t _{PW}	0.19	0.20	0.21	ms	
PGM pulse width for overwrite programming	t _{OPW} *3	0.19	_	5.25	ms	_
CE setup time	t _{CES}	2			μs	
V _{CC} setup time	t _{VCS}	2			μs	
Data output delay time	t _{OE}	0	_	200	ns	

Notes: 1. Input pulse level: 0.45 V to 2.2 V

Input rise time/fall time \leq 20 ns Timing reference levels Input: 0.8 V, 2.0 V

Output: 0.8 V, 2.0 V

2. t_{DF} is defined at the point at which the output is floating and the output level cannot be read.

3. t_{OPW} is defined by the value given in figure 6.4, High-Speed, High-Reliability Programming Flow Chart.

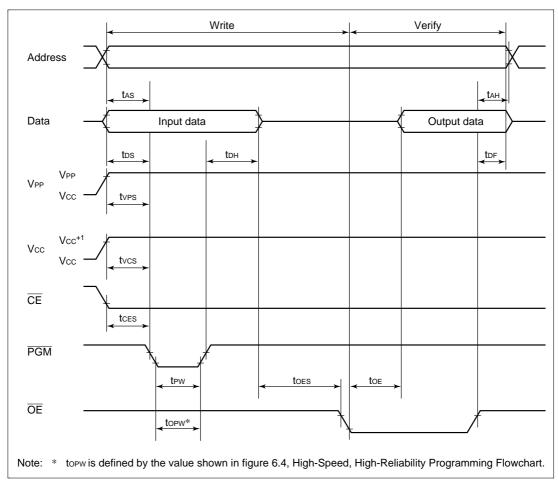


Figure 6.5 shows a PROM write/verify timing diagram.

Figure 6.5 PROM Write/Verify Timing

6.3.2 Programming Precautions

• Use the specified programming voltage and timing.

The programming voltage in PROM mode (V_{PP}) is 12.5 V. Use of a higher voltage can permanently damage the chip. Be especially careful with respect to PROM programmer overshoot.

Setting the PROM programmer to Hitachi specifications for the HN27C101 will result in correct V_{PP} of 12.5 V.

- Make sure the index marks on the PROM programmer socket, socket adapter, and chip are properly aligned. If they are not, the chip may be destroyed by excessive current flow. Before programming, be sure that the chip is properly mounted in the PROM programmer.
- Avoid touching the socket adapter or chip while programming, since this may cause contact faults and write errors.
- Take care when setting the programming mode, as page programming is not supported.
- When programming with a PROM programmer, be sure to specify addresses from H'0000 to H'3FFF. If programming is inadvertently performed from H'4000 onward, it may not be possible to continue PROM programming and verification. When programming, H'FF should be set as the data in address area H'4000 to H'1FFFF.

6.4 Reliability of Programmed Data

A highly effective way to improve data retention characteristics is to bake the programmed chips at 150°C, then screen them for data errors. This procedure quickly eliminates chips with PROM memory cells prone to early failure.

Figure 6.6 Shows the recommended screening procedure.

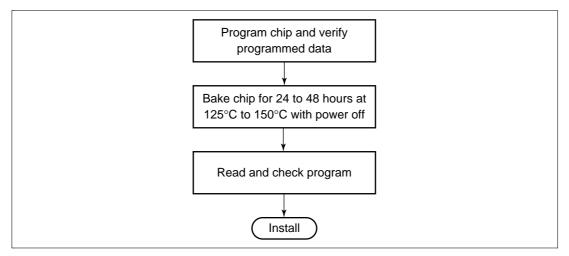


Figure 6.6 Recommended Screening Procedure

If a series of programming errors occurs while the same PROM programmer is in use, stop programming and check the PROM programmer and socket adapter for defects. Please inform Hitachi of any abnormal conditions noted during or after programming or in screening of program data after high-temperature baking.

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Section 7 RAM

7.1 Overview

The H8/3802 has 1 kbyte of high-speed static RAM on-chip, and the H8/3801 and H8/3800 have 512 bytes. The RAM is connected to the CPU by a 16-bit data bus, allowing high-speed 2-state access for both byte data and word data.

7.1.1 Block Diagram

Figure 7.1 shows a block diagram of the on-chip RAM.

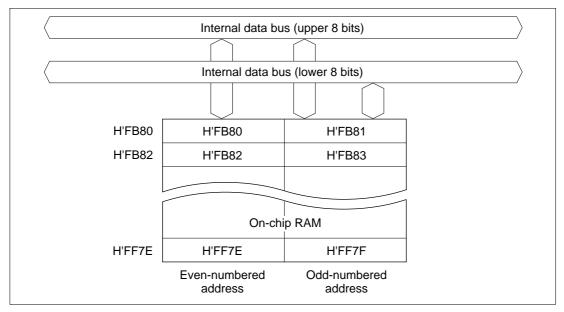


Figure 7.1 RAM Block Diagram (H8/3802)

Section 8 I/O Ports

8.1 Overview

The H8/3802 Series is provided with three 8-bit I/O ports, one 7-bit I/O port, one 4-bit I/O port, one 3-bit I/O port, one 1-bit I/O port, one 4-bit input-only port, one 1-bit input-only port, and one 6-bit output-only port. Table 8.1 indicates the functions of each port.

Each port has of a port control register (PCR) that controls input and output, and a port data register (PDR) for storing output data. Input or output can be assigned to individual bits. See 2.9.2, Notes on Bit Manipulation, for information on executing bit-manipulation instructions to write data in PCR or PDR.

Ports 5, 6, 7, 8, and A are also used as liquid crystal display segment and common pins, selectable in 4-bit units.

Block diagrams of each port are given in Appendix C, I/O Port Block Diagrams.

Port	Description	Pins	Other Functions	Function Switching Registers
Port 3	 7-bit I/O port MOS input pull-up option Large-current port 	P3 ₇ /AEVL P3 ₆ /AEVH P3 _{5,} P3 _{4,} P3 ₃	Asynchronous event counter event inputs AEVL, AEVH	PMR3
		P3 ₂ , TMOFH P3 ₁ , TMOFL	Timer F output compare output	PMR3
Port 4	1-bit input port	P4 ₃ /IRQ ₀	External interrupt 0	PMR2
	• 3-bit I/O port	P4 ₂ /TXD ₃₂ P4 ₁ /RXD ₃₂ P4 ₀ /SCK ₃₂	SCI3 data output (TXD ₃₂), data input (RXD ₃₂), clock input/output (SCK ₃₂)	SCR3 SMR
Port 5	 8-bit I/O port MOS input pull-up option	$P5_7$ to $P5_0/$ WKP ₇ to WKP ₀ / SEG ₈ to SEG ₁	Wakeup input (\overline{WKP}_7 to \overline{WKP}_0), segment output (SEG ₈ to SEG ₁)	PMR5 LPCR

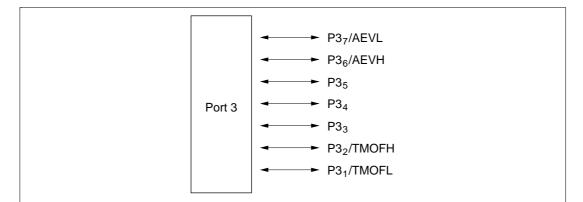
Table 8.1Port Functions

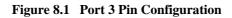
Port	Description	Pins	Other Functions	Function Switching Registers
Port 6	 8-bit I/O port MOS input pull-up option	$P6_7$ to $P6_0/SEG_{16}$ to SEG_9	Segment output (SEG ₁₆ to SEG ₉)	LPCR
Port 7	8-bit I/O port	$P7_7$ to $P7_0/$ SEG ₂₄ to SEG ₁₇	Segment output (SEG ₂₄ to SEG ₁₇)	LPCR
Port 8	1-bit I/O port	P8 ₀ /SEG ₂₅ ,	Segment output	LPCR
Port 9	6-bit output port	P9 ₅ to P9 ₂	None	
	High-voltage, large- current port	P9 ₁ , P9 ₀ / PWM2, PWM1	10-bit PWM output	PMR9
	High-voltage port	IRQAEC	None	
Port A	4-bit I/O port	PA_3 to $PA_0/$ COM ₄ to COM ₁	Common output (COM_4 to COM_1)	LPCR
Port B	4-bit input port	PB_3 to PB_0/AN_3 to AN_0	A/D converter analog input	AMR

8.2 Port 3

8.2.1 Overview

Port 3 is a 7-bit I/O port, configured as shown in figure 8.1.





8.2.2 Register Configuration and Description

Table 8.2 shows the port 3 register configuration.

Table 8.2Port 3 Registers

Name	Abbrev.	R/W	Initial Value	Address
Port data register 3	PDR3	R/W	—	H'FFD6
Port control register 3	PCR3	W	_	H'FFE6
Port pull-up control register 3	PUCR3	R/W	_	H'FFE1
Port mode register 3	PMR3	R/W	_	H'FFCA
Port mode register 2	PMR2	R/W		H'FFC9

1. Port data register 3 (PDR3)

Bit	7	6	5	4	3	2	1	0
	P37	P3 ₆	P35	P34	P3 ₃	P32	P3 ₁	—
Initial value	0	0	0	0	0	0	0	_
Read/Write	R/W	R/W	R/W	R/W	R/W	R/W	R/W	_

PDR3 is an 8-bit register that stores data for port 3 pins $P3_7$ to $P3_1$. If port 3 is read while PCR3 bits are set to 1, the values stored in PDR3 are read, regardless of the actual pin states. If port 3 is read while PCR3 bits are cleared to 0, the pin states are read.

2. Port control register 3 (PCR3)

Bit	7	6	5	4	3	2	1	0
	PCR37	PCR3 ₆	PCR3 ₅	PCR3 ₄	PCR3 ₃	PCR3 ₂	PCR3 ₁	—
Initial value	0	0	0	0	0	0	0	
Read/Write	W	W	W	W	W	W	W	W

PCR3 is an 8-bit register for controlling whether each of the port 3 pins $P3_7$ to $P3_1$ functions as an input pin or output pin. Setting a PCR3 bit to 1 makes the corresponding pin an output pin, while clearing the bit to 0 makes the pin an input pin. The settings in PCR3 and in PDR3 are valid only when the corresponding pin is designated in PMR3 as a general I/O pin.

PCR3 is a write-only register. Bits 7 to 1 are always read as 1. Bit 0 is reserved; only 0 can be written to this bit.

3. Port pull-up control register 3 (PUCR3)

Bit	7	6	5	4	3	2	1	0
	PUCR37	PUCR36	PUCR35	PUCR3 ₄	PUCR33	PUCR3 ₂	PUCR31	—
Initial value	0	0	0	0	0	0	0	_
Read/Write	R/W	R/W	R/W	R/W	R/W	R/W	R/W	W

PUCR3 controls whether the MOS pull-up of each of the port 3 pins $P3_7$ to $P3_1$ is on or off. When a PCR3 bit is cleared to 0, setting the corresponding PUCR3 bit to 1 turns on the MOS pull-up for the corresponding pin, while clearing the bit to 0 turns off the MOS pull-up.

Bit 0 is reserved; only 0 can be written to this bit.

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4. Port mode register 3 (PMR3)

Bit	7	6	5	4	3	2	1	0
	AEVL	AEVH	—	_	—	TMOFH	TMOFL	—
Initial value	0	0	_	—		0	0	_
Read/Write	R/W	R/W	W	W	W	R/W	R/W	W

PMR3 is an 8-bit read/write register, controlling the selection of pin functions for port 3 pins.

Bit 7: P3₇/AEVL pin function switch (AEVL)

This bit selects whether pin P37/AEVL is used as P37 or as AEVL.

Bit 7 AEVL	Description	
0	Functions as P3 ₇ I/O pin	(initial value)
1	Functions as AEVL input pin	

Bit 6: P3₆/AEVH pin function switch (AEVH)

This bit selects whether pin $P3_6$ /AEVH is used as $P3_6$ or as AEVH.

Bit 6 AEVH	Description	
0	Functions as P3 ₆ I/O pin	(initial value)
1	Functions as AEVH input pin	

Bits 5 to 3: Reserved bits

Bits 5 to 3 are reserved; only 0 can be written to these bits.

This bit selects whether pin $P3_2/TMOFH$ is used as $P3_2$ or as TMOFH.

Bit 2		
TMOFH	Description	
0	Functions as P3 ₂ I/O pin	
1	Functions as TMOFH output pin	(initial value)

Bit 1: P3₁/TMOFL pin function switch (TMOFL)

This bit selects whether pin $P3_1/TMOFL$ is used as $P3_1$ or as TMOFL.

Bit 1 TMOFL	Description	
0	Functions as P3 ₁ I/O pin	(initial value)
1	Functions as TMOFL output pin	

Bit 0: Reserved bit

Bit 0 is reserved; only 0 can be written to this bit.

5. Port mode register 2 (PMR2)

Bit	7	6	5	4	3	2	1	0
	—	—	POF1	—	—	—	—	IRQ ₀
Initial value	1	1	0	1	1	—	_	0
Read/Write	—	—	R/W	—	_	W	W	R/W

PMR2 is an 8-bit read/write register controlling the PMOS on/off state for the $P3_5$ pin.

Bit 5: P3₅ pin PMOS control (POF1)

This bit controls the on/off state of the $P3_5$ pin output buffer PMOS.

Bit 5		
POF1	Description	
0	CMOS output	(initial value)
1	NMOS open-drain output	

8.2.3 Pin Functions

Table 8.3 shows the port 3 pin functions.

Table 8.3Port 3 Pin Functions

Pin	Pin Functions and Selection Method							
P3 ₇ /AEVL	The pin function dep	3 ₂ in PCR3.						
	AEVL		0		1			
	PCR37	0	1 P3 ₇ output pin		*			
	Pin function	P37 input pin			AEVL input pin			
P3 ₆ /AEVH	The pin function dep	pends on bit AEVH ir	n PMR3 ar	nd bit PCR	3 ₆ in PCR3.			
	AEVH		0		1			
	PCR3 ₆	0		1	*			
	Pin function	P3 ₆ input pin	P3 ₆ ou	tput pin	AEVH input pin			
P3 ₅ to P3 ₃	to P3 ₃ The pin function depends on the corresponding bit in PCR3.							
	PCR3n	0			1			
	Pin function	P3 _n input p	in	P	P3 _n output pin			
	ι				(n = 5 to 3)			
P3 ₂ /TMOFH	The pin function dep	pends on bit TMOFH	l in PMR3	and bit PC	CR3 ₂ in PCR3.			
	TMOFH		0		1			
	PCR3 ₂	0		1	*			
	Pin function	P3 ₂ input pin	P3 ₂ ou	tput pin	TMOFH output pin			
P3 ₁ /TMOFL	The pin function dep	pends on bit TMOFL	in PMR3 a	and bit PC	R3 ₁ in PCR3.			
	TMOFL		0	1				
	PCR3 ₁	0		1	*			
	Pin function	P3 ₁ input pin	P3 ₁ ou	tput pin	THOFL output pin			
					*: Don't cara			

*: Don't care

8.2.4 Pin States

Table 8.4 shows the port 3 pin states in each operating mode.

Table 8.4Port 3 Pin States

Pins	Reset	Sleep	Subsleep	Standby	Watch	Subactive	Active
$\begin{array}{c} P3_{7/AEVL}\\ P3_{6/AEVH}\\ P3_{5}\\ P3_{4}\\ P3_{3}\\ P3_{2/TMOFH}\\ P3_{1/TMOFL} \end{array}$	High- impedance	Retains previous state	Retains previous state	High- impedance*		Functional	Functional

Note: * A high-level signal is output when the MOS pull-up is in the on state.

8.2.5 MOS Input Pull-Up

Port 3 has a built-in MOS input pull-up function that can be controlled by software. When a PCR3 bit is cleared to 0, setting the corresponding PUCR3 bit to 1 turns on the MOS pull-up for that pin. The MOS pull-up function is in the off state after a reset.

PCR3 _n	0	0	1
PUCR3 _n	0	1	*
MOS input pull-up	Off	On	Off

(n = 7 to 1) *: Don't care

8.3 Port 4

8.3.1 Overview

Port 4 is a 3-bit I/O port and 1-bit input port, configured as shown in figure 8.2.

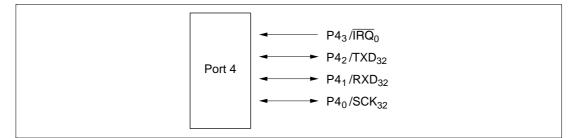


Figure 8.2 Port 4 Pin Configuration

8.3.2 Register Configuration and Description

Table 8.5 shows the port 4 register configuration.

Table 8.5Port 4 Registers

Name	Abbrev.	R/W	Initial Value	Address
Port data register 4	PDR4	R/W	H'F8	H'FFD7
Port control register 4	PCR4	W	H'F8	H'FFE7
Port mode register 2	PMR2	R/W	_	H'FFC9

1. Port data register 4 (PDR4)

Bit	7	6	5	4	3	2	1	0
			_	_	P43	P42	P4 1	P4 0
Initial value	1	1	1	1	1	0	0	0
Read/Write	—	—	—	—	R	R/W	R/W	R/W

PDR4 is an 8-bit register that stores data for port 4 pins $P4_2$ to $P4_0$. If port 4 is read while PCR4 bits are set to 1, the values stored in PDR4 are read, regardless of the actual pin states. If port 4 is read while PCR4 bits are cleared to 0, the pin states are read.

Upon reset, PDR4 is initialized to H'F8.

2. Port control register 4 (PCR4)

Bit	7	6	5	4	3	2	1	0
	_	—	—	—	_	PCR4 ₂	PCR41	PCR40
Initial value	1	1	1	1	1	0	0	0
Read/Write		_	_	_	_	W	W	W

PCR4 is an 8-bit register for controlling whether each of port 4 pins $P4_2$ to $P4_0$ functions as an input pin or output pin. Setting a PCR4 bit to 1 makes the corresponding pin an output pin, while clearing the bit to 0 makes the pin an input pin. PCR4 and PDR4 settings are valid when the corresponding pins are designated for general-purpose input/output by SCR3-2.

Upon reset, PCR4 is initialized to H'F8.

PCR4 is a write-only register, which is always read as all 1s.

3. Port mode register 2 (PMR2)

Bit	7	6	5	4	3	2	1	0
	_		POF1	_	—		—	IRQ ₀
Initial value	1	1	0	1	1	_		0
Read/Write	_	_	R/W	_	_	W	W	R/W

PMR2 is an 8-bit read/write register controlling the selection of the $P4_3/IRQ_0$ pin function and the PMOS on/off state for the $P3_5$ pin. Upon reset, PMR2 is initialized to H'DE.

Bits 7, 6, 4, and 3: Reserved bits

Bits 7, 6, and 4 to 1 are reserved; they are always read as 1 and cannot be modified.

Bit 5: P3₅ pin PMOS control (POF1)

This bit controls the on/off state of the P3₅ pin output buffer PMOS.

Bit 5 POF1	Description	
0	CMOS output	(initial value)
1	NMOS open-drain output	

Bits 2 and 1: Reserved bits

Bits 2 and 1 are reserved; only 0 can be written to these bits.

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Bit 0: P4₃/IRQ₀ pin function switch (IRQ₀)

This bit selects whether pin P4_3/ $\overline{\text{IRQ}}_0$ is used as P4_3 or as $\overline{\text{IRQ}}_0.$

Bit 0 IRQ₀	Description	
0	Functions as $P4_3$ input pin	(initial value)
1	Functions as $\overline{\mathrm{IRQ}}_{0}$ input pin	

8.3.3 **Pin Functions**

Table 8.6 shows the port 4 pin functions.

Table 8.6Port 4 Pin Functions

he pin function dep IRQ0 Pin function	ends on bit IRQ0 in 0 P4 ₃ input p		ĪF	1			
		in	ĪĒ				
Pin function	P4 ₃ input p	in	ĪF				
			11	RQ ₀ input pin			
he pin function dep PCR4.	ends on bit TE in S	CR3, bit SP	C32 in S	PCR, and bit PCR4 ₂			
SPC32		0		1			
TE		1					
PCR4 ₂	0	1		*			
Pin function	P4 ₂ input pin	P4 ₂ out	out pin	TXD ₃₂ output pin			
he pin function dep	ends on bit RE in S	CR3 and bi	t PCR4 ₁	in PCR4.			
RE			1				
PCR4 ₁	0	*					
Pin function	P4 ₁ input pin	•					
	PCR4. SPC32 TE PCR4 ₂ Pin function he pin function dep RE PCR4 ₁	PCR4. SPC32 TE PCR42 0 Pin function he pin function depends on bit RE in S RE PCR41 0	PCR4.SPC320TE0PCR420Pin functionP42 input pinP42 outputPerformed by the pin function depends on bit RE in SCR3 and bitRE0PCR410PCR410	PCR4.SPC320TE0PCR420Pin functionP42 input pinP42 output pinP10 output pinP1			

Pin	Pin Functions and Selection Method								
P40/SCK32		The pin function depends on bit CKE1 and CKE0 in SCR3, bit COM in SMR, and bit PCR4 ₀ in PCR4.							
	CKE1		0						
	CKE0		0	1	*				
	СОМ	0 1			*	*			
	PCR4 ₀	0	1	*		*			
	Pin function	P4 ₀ input pin P4 ₀ output pin SCK ₃₂ out pin			SCK ₃₂ input pin				
			h.			*: Don't caro			

*: Don't care

8.3.4 Pin States

Table 8.7 shows the port 4 pin states in each operating mode.

Table 8.7Port 4 Pin States

Pins	Reset	Sleep	Subsleep	Standby	Watch	Subactive	Active
P4 ₃ /IRQ ₀ P4 ₂ /TXD ₃₂ P4 ₁ /RXD ₃₂ P4 ₀ /SCK ₃₂	High- impedance	Retains previous state		High- impedance		Functional	Functional

8.4 Port 5

8.4.1 Overview

Port 5 is an 8-bit I/O port, configured as shown in figure 8.3.

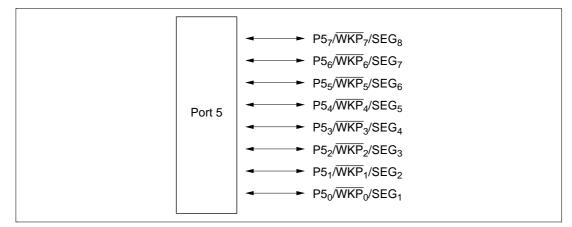


Figure 8.3 Port 5 Pin Configuration

8.4.2 Register Configuration and Description

Table 8.8 shows the port 5 register configuration.

Table 8.8Port 5 Registers

Name	Abbrev.	R/W	Initial Value	Address
Port data register 5	PDR5	R/W	H'00	H'FFD8
Port control register 5	PCR5	W	H'00	H'FFE8
Port pull-up control register 5	PUCR5	R/W	H'00	H'FFE2
Port mode register 5	PMR5	R/W	H'00	H'FFCC

1. Port data register 5 (PDR5)

Bit	7	6	5	4	3	2	1	0
	P57	P5 ₆	P5 ₅	P54	P53	P5 ₂	P5 ₁	P50
Initial value	0	0	0	0	0	0	0	0
Read/Write	R/W	R/W	R/W	R/W	R/W	R/W	R/W	R/W

PDR5 is an 8-bit register that stores data for port 5 pins $P5_7$ to $P5_0$. If port 5 is read while PCR5 bits are set to 1, the values stored in PDR5 are read, regardless of the actual pin states. If port 5 is read while PCR5 bits are cleared to 0, the pin states are read.

Upon reset, PDR5 is initialized to H'00.

2. Port control register 5 (PCR5)

Bit	7	6	5	4	3	2	1	0
	PCR57	PCR5 ₆	PCR55	PCR5 ₄	PCR53	PCR5 ₂	PCR51	PCR50
Initial value	0	0	0	0	0	0	0	0
Read/Write	W	W	W	W	W	W	W	W

PCR5 is an 8-bit register for controlling whether each of the port 5 pins $P5_7$ to $P5_0$ functions as an input pin or output pin. Setting a PCR5 bit to 1 makes the corresponding pin an output pin, while clearing the bit to 0 makes the pin an input pin. PCR5 and PDR5 settings are valid when the corresponding pins are designated for general-purpose input/output by PMR5 and bits SGS3 to SGS0 in LPCR.

Upon reset, PCR5 is initialized to H'00.

PCR5 is a write-only register, which is always read as all 1s.

3. Port pull-up control register 5 (PUCR5)

Bit	7	6	5	4	3	2	1	0
	PUCR57	PUCR5 ₆	PUCR55	PUCR54	PUCR53	PUCR52	PUCR51	PUCR50
Initial value	0	0	0	0	0	0	0	0
Read/Write	R/W	R/W	R/W	R/W	R/W	R/W	R/W	R/W

PUCR5 controls whether the MOS pull-up of each of port 5 pins $P5_7$ to $P5_0$ is on or off. When a PCR5 bit is cleared to 0, setting the corresponding PUCR5 bit to 1 turns on the MOS pull-up for the corresponding pin, while clearing the bit to 0 turns off the MOS pull-up.

Upon reset, PUCR5 is initialized to H'00.

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4. Port mode register 5 (PMR5)

Bit	7	6	5	4	3	2	1	0
	WKP ₇	WKP ₆	WKP ₅	WKP ₄	WKP ₃	WKP ₂	WKP ₁	WKP ₀
Initial value	0	0	0	0	0	0	0	0
Read/Write	R/W							

PMR5 is an 8-bit read/write register, controlling the selection of pin functions for port 5 pins.

Upon reset, PMR5 is initialized to H'00.

Bit n: $P5_n/\overline{WKP}_n/SEG_{n+1}$ pin function switch (WKPn)

When pin P5n/ $\overline{WKP}n/SEGn+1$ is not used as SEG_{n+1} , these bits select whether the pin is used as P5n or \overline{WKP}_n .

Bit n WKPn	Description	
0	Functions as P5n I/O pin	(initial value)
1	Functions as \overline{WKP}_{n} input pin	

Note: For use as SEG_{n+1} , see 13.2.1, LCD Port Control Register (LPCR).

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(n = 7 to 0)

8.4.3 Pin Functions

Table 8.9 shows the port 5 pin functions.

Table 8.9Port 5 Pin Functions

Pin	Pin Functions and Selection Method							
P5 ₇ /WKP ₇ / SEG ₈ to P5 ₀ /WKP ₀ /	The pin function depends on bit WKP_n in PMR5, bit $PCR5_n$ in PCR5, and bits SGS3 to SGS0 in LPCR.							
SEG1	$P5_7$ to $P5_4$	(n = 7 to 4)						
	SGS3 to SGS0		010, 0011, 0100 0111, 1000, 100		0010, 0011, 0100, 0101, 0110, 0111, 1000, 1001			
WKP _n 0				1	*			
	PCR5 _n	0	1 *		*			
	Pin function	P5 _n input pin	P5 _n output pin	WKPn input pin	SEGn+1 output pin			
	$P5_3 \text{ to } P5_0$ (m= 3 to 0)							
	SGS3 to SGS0	Other than 00	0001, 0010, 0011, 0100, 0101, 0110, 0111, 1000					
	WKP _m		0	1	*			
	PCR5 _m	0	1	*	*			
	Pin function	P5 _m input pin	P5 _m output pin	WKPm output pin	SEGm+1 output pin			

*: Don't care

8.4.4 Pin States

Table 8.10 shows the port 5 pin states in each operating mode.

Table 8.10Port 5 Pin States

Pins	Reset	Sleep	Subsleep	Standby	Watch	Subactive	Active
$P5_7/\overline{WKP}_7/$ SEG ₈ to P5 ₀ / WKP ₀ /SEG ₁	High- impedance	Retains previous state		High- impedance*		Functional	Functional

Note: * A high-level signal is output when the MOS pull-up is in the on state.

8.4.5 MOS Input Pull-Up

Port 5 has a built-in MOS input pull-up function that can be controlled by software. When a PCR5 bit is cleared to 0, setting the corresponding PUCR5 bit to 1 turns on the MOS pull-up for that pin. The MOS pull-up function is in the off state after a reset.

PCR5 _n	0	0	1
PUCR5 _n	0	1	*
MOS input pull-up	Off	On	Off

(n = 7 to 0) *: Don't care

8.5 Port 6

8.5.1 Overview

Port 6 is an 8-bit I/O port. The port 6 pin configuration is shown in figure 8.4.

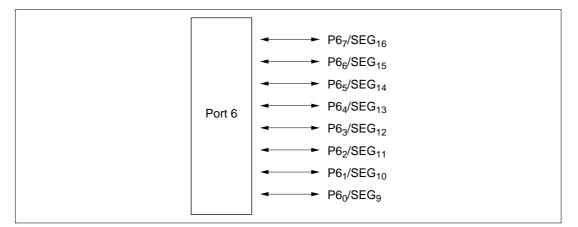


Figure 8.4 Port 6 Pin Configuration

8.5.2 Register Configuration and Description

Table 8.11 shows the port 6 register configuration.

Table 8.11 Port 6 Registers

Name	Abbrev.	R/W	Initial Value	Address
Port data register 6	PDR6	R/W	H'00	H'FFD9
Port control register 6	PCR6	W	H'00	H'FFE9
Port pull-up control register 6	PUCR6	R/W	H'00	H'FFE3

1. Port data register 6 (PDR6)

Bit	7	6	5	4	3	2	1	0
	P67	P6 ₆	P6 ₅	P64	P63	P62	P61	P6 ₀
Initial value	0	0	0	0	0	0	0	0
Read/Write	R/W	R/W	R/W	R/W	R/W	R/W	R/W	R/W

PDR6 is an 8-bit register that stores data for port 6 pins P67 to P60.

If port 6 is read while PCR6 bits are set to 1, the values stored in PDR6 are read, regardless of the actual pin states. If port 6 is read while PCR6 bits are cleared to 0, the pin states are read.

Upon reset, PDR6 is initialized to H'00.

2. Port control register 6 (PCR6)

Bit	7	6	5	4	3	2	1	0
	PCR67	PCR6 ₆	PCR6 ₅	PCR6 ₄	PCR6 ₃	PCR6 ₂	PCR6 ₁	PCR6 ₀
Initial value	0	0	0	0	0	0	0	0
Read/Write	W	W	W	W	W	W	W	W

PCR6 is an 8-bit register for controlling whether each of the port 6 pins $P6_7$ to $P6_0$ functions as an input pin or output pin.

Setting a PCR6 bit to 1 makes the corresponding pin (P6₇ to P6₀) an output pin, while clearing the bit to 0 makes the pin an input pin. PCR6 and PDR6 settings are valid when the corresponding pins are designated for general-purpose input/output by bits SGS3 to SGS0 in LPCR.

Upon reset, PCR6 is initialized to H'00.

PCR6 is a write-only register, which is always read as all 1s.

3. Port pull-up control register 6 (PUCR6)

Bit	7	6	5	4	3	2	1	0
	PUCR67	PUCR66	PUCR65	PUCR64	PUCR63	PUCR6 ₂	PUCR61	PUCR60
Initial value	0	0	0	0	0	0	0	0
Read/Write	R/W	R/W	R/W	R/W	R/W	R/W	R/W	R/W

PUCR6 controls whether the MOS pull-up of each of the port 6 pins $P6_7$ to $P6_0$ is on or off. When a PCR6 bit is cleared to 0, setting the corresponding PUCR6 bit to 1 turns on the MOS pull-up for the corresponding pin, while clearing the bit to 0 turns off the MOS pull-up.

Upon reset, PUCR6 is initialized to H'00.

8.5.3 Pin Functions

Table 8.12 shows the port 6 pin functions.

Table 8.12Port 6 Pin Functions

Pin	Pin Functions and Selection Method									
P6 ₇ /SEG ₁₆ to P6 ₀ /SEG ₉	The pin function dep LPCR.	pends on bit PCR6n	in PCR6 and bits SC	SS3 to SGS0 in						
	P6 ₇ to P6 ₄			(n = 7 to 4)						
	SEG3 to SEGS0	1000, 1001, 1010, 1011 0111, 10		0100, 0101, 0110, 0111, 1000, 1001, 1010, 1011						
	PCR6 _n	0 1		*						
	Pin function	P6 _n input pin P6 _n output pin		SEG _{n+9} output pin						
	P6 ₃ to P6 ₀	$P6_3 \text{ to } P6_0$ (m = 3 to 0)								
	SEG3 to SEGS0		Other than 0011, 0100, 0101, 0110, 0111, 1000, 1001, 1010							
	PCR6 _m	0 1		*						
	Pin function	P6 _m input pin	SEG _{n+9} output pin							

*: Don't care

8.5.4 Pin States

Table 8.13 shows the port 6 pin states in each operating mode.

Table 8.13 Port 6 Pin States

Pin	Reset	Sleep	Subsleep	Standby	Watch	Subactive	Active
P6 ₇ /SEG ₁₆ to P6 ₀ /SEG ₉	High- impedance			High- impedance*		Functional	Functional

Note: * A high-level signal is output when the MOS pull-up is in the on state.

8.5.5 MOS Input Pull-Up

Port 6 has a built-in MOS pull-up function that can be controlled by software. When a PCR6 bit is cleared to 0, setting the corresponding PUCR6 bit to 1 turns on the MOS pull-up for that pin. The MOS pull-up function is in the off state after a reset.

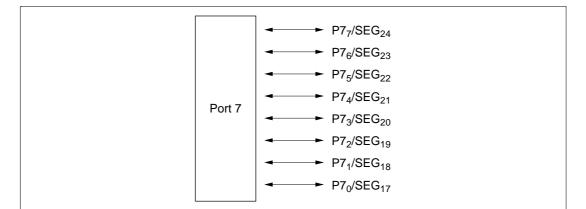
PCR6 _n	0	0	1
PUCR6 _n	0	1	*
MOS input pull-up	Off	On	Off

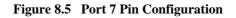
(n = 7 to 0) *: Don't care

8.6 Port 7

8.6.1 Overview

Port 7 is an 8-bit I/O port, configured as shown in figure 8.5.





8.6.2 Register Configuration and Description

Table 8.14 shows the port 7 register configuration.

Table 8.14 Port 7 Registers

Name	Abbrev.	R/W	Initial Value	Address
Port data register 7	PDR7	R/W	H'00	H'FFDA
Port control register 7	PCR7	W	H'00	H'FFEA

1. Port data register 7 (PDR7)

Bit	7	6	5	4	3	2	1	0
	P77	P76	P75	P74	P73	P72	P7 ₁	P70
Initial value	0	0	0	0	0	0	0	0
Read/Write	R/W	R/W						

PDR7 is an 8-bit register that stores data for port 7 pins $P7_7$ to $P7_0$. If port 7 is read while PCR7 bits are set to 1, the values stored in PDR7 are read, regardless of the actual pin states. If port 7 is read while PCR7 bits are cleared to 0, the pin states are read.

Upon reset, PDR7 is initialized to H'00.

2. Port control register 7 (PCR7)

Bit	7	6	5	4	3	2	1	0
	PCR77	PCR7 ₆	PCR75	PCR7 ₄	PCR73	PCR72	PCR71	PCR70
Initial value	0	0	0	0	0	0	0	0
Read/Write	W	W	W	W	W	W	W	W

PCR7 is an 8-bit register for controlling whether each of the port 7 pins $P7_7$ to $P7_0$ functions as an input pin or output pin. Setting a PCR7 bit to 1 makes the corresponding pin an output pin, while clearing the bit to 0 makes the pin an input pin. PCR7 and PDR7 settings are valid when the corresponding pins are designated for general-purpose input/output by bits SGS3 to SGS0 in LPCR.

Upon reset, PCR7 is initialized to H'00.

PCR7 is a write-only register, which is always read as all 1s.

8.6.3 Pin Functions

Table 8.15 shows the port 7 pin functions.

Table 8.15Port 7 Pin Functions

Pin	Pin Functions and Selection Method									
P7 ₇ /SEG ₂₄ to P7 ₀ /SEG ₁₇	The pin function dep LPCR.	ends on bit PCR7 _n i	n PCR7 and bits SG	S3 to SGS0 in						
	$P7_7$ to $P7_4$			(n = 7 to 4)						
	SEGS3 to SEGS0	1010, 1011, 1100, 1101 100		0110, 0111, 1000, 1001, 1010, 1011, 1100, 1101						
	PCR7 _n	0	1	*						
	Pin function	P7 _n input pin P7 _n output pin		SEG _{n+17} output pin						
	$P7_3$ to $P7_0$			(m = 3 to 0)						
	SEGS3 to SEGS0 Other than 0101, 0110, 0111, 1000, 1001, 1010, 1011, 1100			0101, 0110, 0111, 1000, 1001, 1010, 1011, 1100						
	PCR7 _m	0	1	*						
Pin function		P7 _m input pin	P7 _m output pin	SEG _{m+17} output pin						

*: Don't care

8.6.4 Pin States

Table 8.16 shows the port 7 pin states in each operating mode.

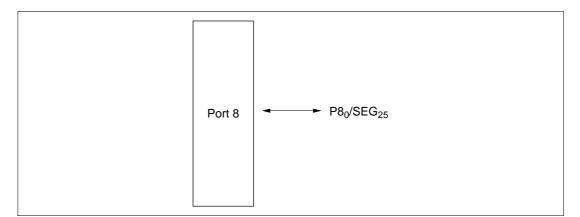
Table 8.16Port 7 Pin States

Pins	Reset	Sleep	Subsleep	Standby	Watch	Subactive	Active
$P7_7/SEG_{24}$ to $P7_0/SEG_{17}$	High- impedance	Retains previous state		High- impedance		Functional	Functional

8.7 Port 8

8.7.1 Overview

Port 8 is an 1-bit I/O port configured as shown in figure 8.6.





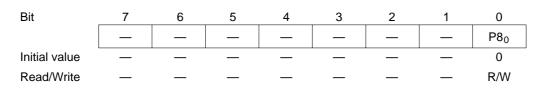
8.7.2 Register Configuration and Description

Table 8.17 shows the port 8 register configuration.

Table 8.17 Port 8 Registers

Name	Abbrev.	R/W	Initial Value	Address
Port data register 8	PDR8	R/W	—	H'FFDB
Port control register 8	PCR8	W	_	H'FFEB

1. Port data register 8 (PDR8)



PDR8 is an 8-bit register that stores data for port 8 pin $P8_0$. If port 8 is read while PCR8 bits are set to 1, the values stored in PDR8 are read, regardless of the actual pin states. If port 8 is read while PCR8 bits are cleared to 0, the pin states are read.

2. Port control register 8 (PCR8)

Bit	7	6	5	4	3	2	1	0
	—	—	_	—	—	_	_	PCR8
Initial value	_				_	_		0
Read/Write	W	W	W	W	W	W	W	W

PCR8 is an 8-bit register for controlling whether the port 8 pin $P8_0$ functions as an input or output pin. Setting a PCR8 bit to 1 makes the corresponding pin an output pin, while clearing the bit to 0 makes the pin an input pin. PCR8 and PDR8 settings are valid when the corresponding pins are designated for general-purpose input/output by bits SGS3 to SGS0 in LPCR.

PCR8 is a write-only register, which is always read as all 1s.

Bits 7 to 1: Reserved bits

Bits 7 to 1 are reserved; only 0 can be written to these bits.

8.7.3 Pin Functions

Table 8.18 shows the port 8 pin functions.

Table 8.18Port 8 Pin Functions

Pin Functions and Selection Method Pin The pin function depends on bit PCR8_n in PCR8 and bits SGS3 to SGS0 in LPCR. P80/SEG25 SEGS3 to SEGS0 Other than 0111, 1000, 1001, 1010, 1011, 0111, 1000, 1001, 1010, 1011, 1100, 1100, 1101, 1110 1101, 1110 PCR80 * 0 1 Pin function P80 input pin P8₀ output pin SEG₂₅ output pin

*: Don't care

8.7.4 Pin States

Table 8.19 shows the port 8 pin states in each operating mode.

Table 8.19Port 8 Pin States

Pins	Reset	Sleep	Subsleep	Standby	Watch	Subactive	Active
P8 ₀ /SEG ₂₅	High- impedance	Retains previous state		High- impedance		Functional	Functional

8.8 Port 9

8.8.1 Overview

Port 9 is a 6-bit output-only port, configured as shown in figure 8.7.

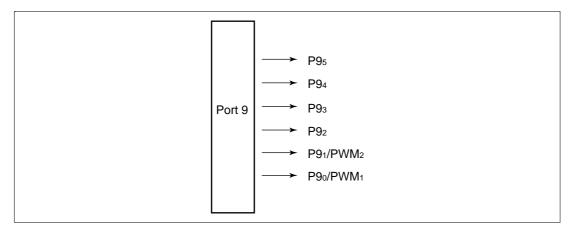


Figure 8.7 Port 5 Pin Configuration

8.8.2 Register Configuration and Description

Table 8.20 shows the port 9 register configuration.

Table 8.20 Port 9 Registers

Name	Abbrev.	R/W	Initial Value	Address
Port data register 9	PDR9	R/W	H'FF	H'FFDC
Port mode register 9	PMR9	R/W		H'FFEC

^{1.} Port data register 9 (PDR9)

Bit	7	6	5	4	3	2	1	0
	—	—	P9 ₅	P94	P93	P9 ₂	P91	P9 ₀
Initial value	1	1	1	1	1	1	1	1
Read/Write	—	—	R/W	R/W	R/W	R/W	R/W	R/W

PDR9 is an 8-bit register that stores data for port 9 pins P95 to P90.

Upon reset, PDR9 is initialized to H'FF.

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2. Port mode register 9 (PMR9)

Bit	7	6	5	4	3	2	1	0
	—	—	—	_	PIOFF		PWM ₂	PWM ₁
Initial value	1	1	1	1	0	_	0	0
Read/Write	—	—	—	—	R/W	W	R/W	R/W

PMR9 is an 8-bit read/write register controlling the selection of the $P9_0$ and $P9_1$ pin functions.

Bit 3: P9₂ to P9₀ step-up circuit control (PIOFF)

Bit 3 turns the $P9_2$ to $P9_0$ step-up circuit on and off.

Bit 3	

PIOFF	Description	
0	Large-current port step-up circuit is turned on	(initial value)
1	Large-current port step-up circuit is turned off	
Note:	When turning the step-up circuit on or off, the register must be rewritten on buffer NMOS is off (port data is 1). When turning the step-up circuit on, first clear PIOFF to 0, then wait for the system clock before turning the buffer NMOS on (clearing port data to 0). Without the elapse of the 30 system clock interval the step-up circuit will no will not be possible for a large current to flow, making operation unstable.	elapse of 30

Port 9 Pin Output Low Level Permitted Currents

Pin	Symbol	Test Conditions	Min	Тур	Max	PIOFF Bit Value
$P9_2$ to $P9_0$	I _{ol}	V_{cc} = 1.8 V to 5.5 V*		_	25 mA*	0
					10 mA	1
$\mathrm{P9}_{\scriptscriptstyle 3}$ to $\mathrm{P9}_{\scriptscriptstyle 5}$			_		10 mA	_

Note: * For details, see section 14.2.2, DC Characteristics.

Bit 2: Reserved bit

Bit 2 is reserved; only 0 can be written to this bit.

Bits 1 and 0: P9_n/PWM pin function switches

These pins select whether pin P9n/PWMn+1 is used as P9n or as PWMn+1.

Bit n		
WKPn+1	Description	
0	Functions as P9 _n output pin	(initial value)
1	Functions as PWM _{n+1} output pin	

(n = 0 or 1)

8.8.3 Pin Functions

Table 8.21 shows the port 9 pin functions.

Table 8.21Port 9 Pin Functions

Pin	Pin Functions and Selection Method
P9 ₁ /PWM _{n+1} to P9 ₀ /PWM _{n+1}	The pin function depends on bit WKP _n in PMR5, bit PCR5 _n in PCR5, and bits SGS3 to SGS0 in LPCR.
	(n = 1 or 0)

PMR9 _n	0	1
Pin function	P9 _n output pin	PWM _{n+1} output pin

*: Don't care

8.8.4 Pin States

Table 8.22 shows the port 5 pin states in each operating mode.

Table 8.22Port 5 Pin States

Pins	Reset	Sleep	Subsleep	Standby	Watch	Subactive	Active
$\begin{array}{l} P9_5 \text{ to } P9_2 \\ P9_n/PWM_{n+1} \text{ to} \\ P9_n/PWM_{n+1} \end{array}$	High- impedance		Retains previous state	High- impedance*		Functional	Functional
							(()

(n = 1 or 0)

8.9 Port A

8.9.1 Overview

Port A is a 4-bit I/O port, configured as shown in figure 8.8.

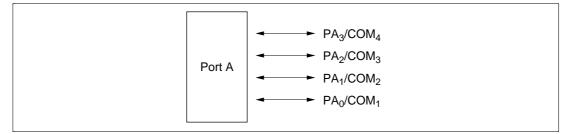


Figure 8.8 Port A Pin Configuration

8.9.2 Register Configuration and Description

Table 8.23 shows the port A register configuration.

Table 8.23Port A Registers

Name	Abbrev.	R/W	Initial Value	Address
Port data register A	PDRA	R/W	H'F0	H'FFDD
Port control register A	PCRA	W	H'F0	H'FFED

1. Port data register A (PDRA)

Bit	7	6	5	4	3	2	1	0
	—	—	—	—	PA ₃	PA ₂	PA ₁	PA ₀
Initial value	1	1	1	1	0	0	0	0
Read/Write	—	—	—	—	R/W	R/W	R/W	R/W

PDRA is an 8-bit register that stores data for port A pins PA_3 to PA_0 . If port A is read while PCRA bits are set to 1, the values stored in PDRA are read, regardless of the actual pin states. If port A is read while PCRA bits are cleared to 0, the pin states are read.

Upon reset, PDRA is initialized to H'F0.

2. Port control register A (PCRA)

Bit	7	6	5	4	3	2	1	0
		—	—		PCRA ₃	$PCRA_2$	PCRA ₁	PCRA ₀
Initial value	1	1	1	1	0	0	0	0
Read/Write	_	_	_	_	R/W	R/W	R/W	R/W

PCRA controls whether each of port A pins PA_3 to PA_0 functions as an input pin or output pin. Setting a PCRA bit to 1 makes the corresponding pin an output pin, while clearing the bit to 0 makes the pin an input pin. PCRA and PDRA settings are valid when the corresponding pins are designated for general-purpose input/output by LPCR.

Upon reset, PCRA is initialized to H'F0.

PCRA is a write-only register, which is always read as all 1s.

8.9.3 Pin Functions

Table 8.24 shows the port A pin functions.

Table 8.24Port A Pin Functions

Pin	Pin Functions and Selection Method					
PA ₃ /COM ₄	The pin function depends on bit $PCRA_3$ in PCRA and bits SGS3 to SGS0.					
	SEGS3 to SEGS0	0000	0000 0000			
	PCRA ₃	0	1	*		
	Pin function	PA3 input pin	PA ₃ output pin	COM ₄ output pin		
PA ₂ /COM ₃ The pin function depends on bit PCRA ₂ in PCRA and bits SGS3 to SGS						
	SEGS3 to SEGS0	0000	0000	Not 0000		
	PCRA ₂	0	1	*		
	Pin function	PA ₂ input pin	PA ₂ output pin	COM ₃ output pin		
PA ₁ /COM ₂	The pin function depends on bit PCRA ₁ in PCRA and bits SGS3 to SGS0.					
	SEGS3 to SEGS0	0000	0000	Not 0000		
	PCRA ₁	0	1	*		
	Pin function	PA ₁ input pin	PA ₁ output pin	COM ₂ output pin		
PA ₀ /COM ₁	The pin function depends on bit PCRA ₀ in PCRA and bits SGS3 to SGS0.					
	SEGS3 to SEGS0	00	000	Not 0000		
	PCRA ₀	0	1	*		
	Pin function	PA ₀ input pin	PA ₀ output pin	COM ₁ output pin		
	ι		1.	1]		

*: Don't care

8.9.4 Pin States

Table 8.25 shows the port A pin states in each operating mode.

Table 8.25Port A Pin States

Pins	Reset	Sleep	Subsleep	Standby	Watch	Subactive	Active
PA ₃ /COM ₄ PA ₂ /COM ₃ PA ₁ /COM ₂ PA ₀ /COM ₁	High- impedance	Retains previous state		High- impedance	Retains previous state	Functional	Functional

8.10 Port B

8.10.1 Overview

Port B is a 4-bit input-only port, configured as shown in figure 8.9.

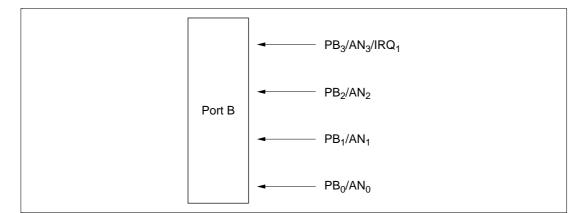


Figure 8.9 Port B Pin Configuration

8.10.2 Register Configuration and Description

Table 8.26 shows the port B register configuration.

Table 8.26Port B Register

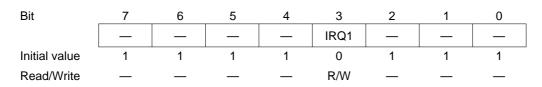
Name	Abbrev.	R/W	Initial Value	Address
Port data register B	PDRB	R	—	H'FFDE
Port mode register B	PMRB	R/W	H'F7	H'FFEE

1. Port Data Register B (PDRB)

Bit	7	6	5	4	3	2	1	0
	_	—	—	_	PB ₃	PB ₂	PB ₁	PB ₀
Read/Write	—	—	—	—	R	R	R	R

Reading PDRB always gives the pin states. However, if a port B pin is selected as an analog input channel for the A/D converter by AMR bits CH3 to CH0, that pin reads 0 regardless of the input voltage.

2. Port mode register B (PMRB)



PMRB is an 8-bit read/write register controlling the selection of the PB_3 pin function. Upon reset, PMRB is initialized to H'F7.

Bits 7 to 4 and 2 to 0: Reserved bits

Bits 7 to 4 and 2 to 0 are reserved; they are always read as 1 and cannot be modified.

Bit 3: $PB_3/AN_3/\overline{IRQ}_1$ pin function switch (IRQ1)

These bits select whether pin PB₃/AN₃/ \overline{IRQ}_1 is used as PB₃/AN₃ or as \overline{IRQ}_1 .

Bit 3 IRQ1	Description	
0	Functions as PB_3/AN_3 input pin	(initial value)
1	Functions as \overline{IRQ}_1 input pin	

Note: Rising or falling edge sensing can be selected for the \overline{IRQ}_1 pin.

8.10.3 Pin Functions

Table 8.27 shows the port B pin functions.

Table 8.27Port B Pin Functions

Pin	Pin Functions and Selection Method						
$PB_3/AN_3/\overline{IRQ}_1$	The pin function depends on bits CH3 to CH0 in AMR and bit IRQ ₁ in PMRB.						
	IRQ ₁	0		1			
	CH3 to CH0	Not 0111	1	*			
	Pin function	PB_3 input pin	AN_3 input pin	$\overline{IRQ}_{\overline{1}}$ input pin			
PB ₂ /AN ₂	The pin function dep	ends on bits CH3 to	CH0 in AMR.				
	CH3 to CH0	Not	0110				
	Pin function	PB ₂ input pin		AN ₂ input pin			
PB ₁ /AN ₁	The pin function depends on bits CH3 to CH0 in AMR.						
	CH3 to CH0	Not 0101		Not 0000			
	Pin function	PB₁ in	put pin	AN₁ input pin			
PB ₀ /AN ₀	The pin function depends on bits CH3 to CH0 in AMR.						
	CH3 to CH0	Not	0100	0100			
	Pin function	PB₀ input pin		AN ₀ input pin			
	L. L.						

*: Don't care

8.11 Input/Output Data Inversion Function

8.11.1 Overview

With input pin RXD_{32} and output pin TXD_{32} , the data can be handled in inverted form.

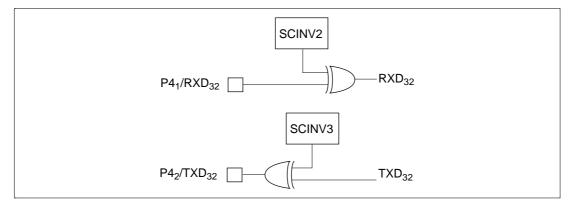


Figure 8.10 Input/Output Data Inversion Function

8.11.2 Register Configuration and Descriptions

Table 8.28 shows the registers used by the input/output data inversion function.

Table 8.28 Register Configuration

Name	Abbreviation	R/W	Address
Serial port control register	SPCR	R/W	H'FF91

Serial Port Control Register (SPCR)

Bit	7	6	5	4	3	2	1	0
	—		SPC32	—	SCINV3	SCINV2	—	—
Initial value	1	1	0		0	0	—	_
Read/Write	—	—	R/W	W	R/W	R/W	W	W

SPCR is an 8-bit readable/writable register that performs RXD_{32} and TXD_{32} pin input/output data inversion switching.

Bits 7 and 6: Reserved bits

Bits 7 and 6 are reserved; they are always read as 1 and cannot be modified.

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Bit 5: P4₂/TXD₃₂ pin function switch (SPC32)

This bit selects whether pin $P4_2/TXD_{32}$ is used as $P4_2$ or as TXD_{32} .

Bit 5 SPC32	Description	
0	Functions as P4 ₂ I/O pin	(initial value)
1	Functions as TXD ₃₂ output pin*	

Note: * Set the TE bit in SCR3 after setting this bit to 1.

Bit 4: Reserved bit

Bit 4 is reserved; only 0 can be written to this bit.

Bit 3: TXD₃₂ pin output data inversion switch

Bit 3 specifies whether or not TXD_{32} pin output data is to be inverted.

Bit 3 SCINV3	Description	
0	TXD ₃₂ output data is not inverted	(initial value)
1	TXD ₃₂ output data is inverted	

Bit 2: RXD₃₂ pin input data inversion switch

Bit 2 specifies whether or not RXD₃₂ pin input data is to be inverted.

Bit 2 SCINV2	Description	
0	RXD ₃₂ input data is not inverted	(initial value)
1	RXD ₃₂ input data is inverted	

Bits 1 and 0: Reserved bits

Bits 1 and 0 are reserved; only 0 can be written to these bits.

8.11.3 Note on Modification of Serial Port Control Register

When a serial port control register is modified, the data being input or output up to that point is inverted immediately after the modification, and an invalid data change is input or output. When modifying a serial port control register, do so in a state in which data changes are invalidated.

8.12 Application Note

8.12.1 How to Handle an Unused Pin

If an I/O pin not used by the user system is floating, pull it up or down.

- If an unused pin is an input pin, handle it in one of the following ways:
 - Pull it up to V_{cc} with an on-chip pull-up MOS.
 - Pull it up to V_{CC} with an external resistor of approximately 100 kΩ.
 - Pull it down to V_{SS} with an external resistor of approximately 100 kΩ.
 - For a pin also used by the A/D converter, pull it up to AV_{CC} .
- If an unused pin is an output pin, handle it in one of the following ways:
 - Set the output of the unused pin to high and pull it up to V_{CC} with an on-chip pull-up MOS.
 - Set the output of the unused pin to high and pull it up to V_{CC} with an external resistor of approximately 100 k Ω .
 - Set the output of the unused pin to low and pull it down to GND with an external resistor of approximately 100 k $\Omega.$

Section 9 Timers

9.1 Overview

The H8/3802 Series provides three timers: timers A, F, and an asynchronous event counter. The functions of these timers are outlined in table 9.1.

Name	Functions	Internal Clock	Event Input Pin	Waveform Output Pin Remarks
Timer A	• 8-bit interval timer	ø/8 to ø/8192	_	_
	Interval function	(8 choices)		
	Time base	ø _w /128 (choice of 4 overflow periods)	-	
Timer F	• 16-bit timer	ø/4 to ø/32, ø _w /4 (4 choices)	_	TMOFL TMOFH
	 Also usable as two independent8-bit timers. 			
	Output compare output function			
Asynchro- nous event counter	 16-bit counter Also usable as two independent 8-bit counters Counts events asynchronous to ø and øw Can count asynchronous events (rising/falling/both edges) independ- ently of the MCU's internal clock 	ø/2 to ø/8 (3 choices)	AEVL AEVH IRQAEC	_

Table 9.1Timer Functions

9.2 Timer A

9.2.1 Overview

Timer A is an 8-bit timer with interval timing and real-time clock time-base functions. The clock time-base function is available when a 32.768 kHz crystal oscillator is connected.

1. Features

Features of timer A are given below.

- Choice of eight internal clock sources (ø/8192, ø/4096, ø/2048, ø/512, ø/256, ø/128, ø/32, ø/8).
- Choice of four overflow periods (1 s, 0.5 s, 0.25 s, 31.25 ms) when timer A is used as a clock time base (using a 32.768 kHz crystal oscillator).
- An interrupt is requested when the counter overflows.
- Use of module standby mode enables this module to be placed in standby mode independently when not used.

2. Block diagram

Figure 9.1 shows a block diagram of timer A.

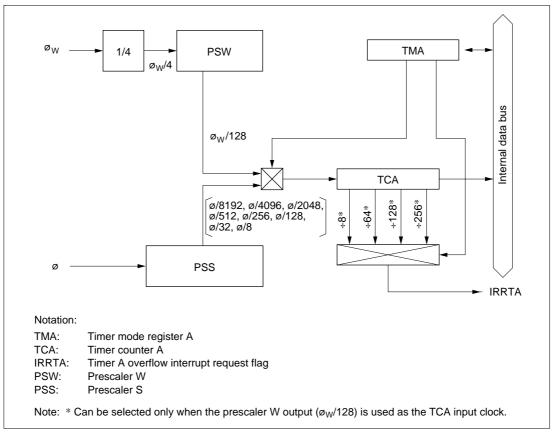


Figure 9.1 Block Diagram of Timer A

3. Register configuration

Table 9.2 shows the register configuration of timer A.

Table 9.2Timer A Registers

Name	Abbrev.	R/W	Initial Value	Address
Timer mode register A	ТМА	R/W	—	H'FFB0
Timer counter A	TCA	R	H'00	H'FFB1
Clock stop register 1	CKSTPR1	R/W	H'FF	H'FFFA

9.2.2 Register Descriptions

1. Timer mode register A (TMA)

Bit	7	6	5	4	3	2	1	0
	—	—	—		ТМАЗ	TMA2	TMA1	TMA0
Initial value	_	—	—	1	0	0	0	0
Read/Write	W	W	W	—	R/W	R/W	R/W	R/W

TMA is an 8-bit read/write register for selecting the prescaler, and input clock.

Bits 7 to 5: Reserved bits

Bits 7 to 5 are reserved; only 0 can be written to these bits.

Bit 4: Reserved bit

Bit 4 is reserved; it is always read as 1, and cannot be modified.

Bits 3 to 0: Internal clock select (TMA3 to TMA0)

Bits 3 to 0 select the clock input to TCA. The selection is made as follows.

				Description	
Bit 3 TMA3	Bit 2 TMA2	Bit 1 TMA1	Bit 0 TMA0	Prescaler and Divider Ratio or Overflow Period	Function
0	0	0	0	PSS, ø/8192 (initial value)	Interval timer
			1	PSS, ø/4096	_
		1	0	PSS, ø/2048	_
			1	PSS, ø/512	_
	1	0	0	PSS, ø/256	_
			1	PSS, ø/128	_
		1	0	PSS, ø/32	_
			1	PSS, ø/8	_
1	0	0	0	PSW, 1 s	Clock time
			1	PSW, 0.5 s	base
		1	0	PSW, 0.25 s	(when using
			1	PSW, 0.03125 s	32.768 kHz)
	1	0	0	PSW and TCA are reset	-
			1	_	
		1	0	_	
			1	_	

2. Timer counter A (TCA)

Bit	7	6	5	4	3	2	1	0
	TCA7	TCA6	TCA5	TCA4	TCA3	TCA2	TCA1	TCA0
Initial value	0	0	0	0	0	0	0	0
Read/Write	R	R	R	R	R	R	R	R

TCA is an 8-bit read-only up-counter, which is incremented by internal clock input. The clock source for input to this counter is selected by bits TMA3 to TMA0 in timer mode register A (TMA). TCA values can be read by the CPU in active mode, but cannot be read in subactive mode. When TCA overflows, the IRRTA bit in interrupt request register 1 (IRR1) is set to 1.

TCA is cleared by setting bits TMA3 and TMA2 of TMA to 11.

Upon reset, TCA is initialized to H'00.

3. Clock stop register 1 (CKSTPR1)

Bit:	7	6	5	4	3	2	1	0
	—	_	S32CKSTP	ADCKSTP	—	TFCKSTP	—	TACKSTP
Initial value:	1	1	1	1	1	1	1	1
Read/Write:	—	_	R/W	R/W	_	R/W	_	R/W

CKSTPR1 is an 8-bit read/write register that performs module standby mode control for peripheral modules. Only the bit relating to timer A is described here. For details of the other bits, see the sections on the relevant modules.

Bit 0: Timer A module standby mode control (TACKSTP)

Bit 0 controls setting and clearing of module standby mode for timer A.

TACKSTP Description 0 Timer A is set to module standby mode 1 Timer A module standby mode is cleared

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9.2.3 Timer Operation

1. Interval timer operation

When bit TMA3 in timer mode register A (TMA) is cleared to 0, timer A functions as an 8-bit interval timer.

Upon reset, TCA is cleared to H'00 and bit TMA3 is cleared to 0, so up-counting and interval timing resume immediately. The clock input to timer A is selected by bits TMA2 to TMA0 in TMA; any of eight internal clock signals output by prescaler S can be selected.

After the count value in TCA reaches H'FF, the next clock signal input causes timer A to overflow, setting bit IRRTA to 1 in interrupt request register 1 (IRR1). If IENTA = 1 in interrupt enable register 1 (IENR1), a CPU interrupt is requested.*

At overflow, TCA returns to H'00 and starts counting up again. In this mode timer A functions as an interval timer that generates an overflow output at intervals of 256 input clock pulses.

Note: * For details on interrupts, see 3.3, Interrupts.

2. Real-time clock time base operation

When bit TMA3 in TMA is set to 1, timer A functions as a real-time clock time base by counting clock signals output by prescaler W. The overflow period of timer A is set by bits TMA1 and TMA0 in TMA. A choice of four periods is available. In time base operation (TMA3 = 1), setting bit TMA2 to 1 clears both TCA and prescaler W to their initial values of H'00.

9.2.4 Timer A Operation States

Table 9.3 summarizes the timer A operation states.

Table 9.3Timer A Operation States

Oper	ation Mode	Reset	Active	Sleep	Watch	Sub- active	Sub- sleep	Standby	Module Standby
TCA	Interval	Reset	Functions	Functions	Halted	Halted	Halted	Halted	Halted
	Clock time base	Reset	Functions	Functions	Functions	Functions	Functions	Halted	Halted
TMA		Reset	Functions	Retained	Retained	Functions	Retained	Retained	Retained

Note: When the real-time clock time base function is selected as the internal clock of TCA in active mode or sleep mode, the internal clock is not synchronous with the system clock, so it is synchronized by a synchronizing circuit. This may result in a maximum error of 1/ø (s) in the count cycle.

9.2.5 Application Note

When bit 0 (TACKSTP) of the clock stop register 1 (CKSTPR1) is cleared to 0, bit 3 (TMA3) of the timer mode register A (TMA) cannot be rewritten.

Set bit 0 (TACKSTP) of the clock stop register 1 (CKSTPR1) to 1 before rewriting bit 3 (TMA3) of the timer mode register A (TMA).

9.3 Timer F

9.3.1 Overview

Timer F is a 16-bit timer with a built-in output compare function. Timer F also provides for counter resetting, interrupt request generation, toggle output, etc., using compare match signals. Timer F can also be used as two independent 8-bit timers (timer FH and timer FL).

1. Features

Features of timer F are given below.

- Choice of four internal clock sources ($\phi/32$, $\phi/16$, $\phi/4$, ϕ w/4)
- TMOFH pin toggle output provided using a single compare match signal (toggle output initial value can be set)
- Counter resetting by a compare match signal
- Two interrupt sources: one compare match, one overflow
- Can operate as two independent 8-bit timers (timer FH and timer FL) (in 8-bit mode).

		Timer FL
	Timer FH 8-Bit Timer*	8-Bit Timer/Event Counter
Internal clock	Choice of 4 (ø/32, ø/16, ø/4, øw/4)	
Toggle output	One compare match signal, output to TMOFH pin(initial value settable)	One compare match signal, output to TMOFL pin (initial value settable)
Counter reset	Counter can be reset by compare mate	ch signal
Interrupt sources	One compare match One overflow	

Note: * When timer F operates as a 16-bit timer, it operates on the timer FL overflow signal.

- Operation in watch mode, subactive mode, and subsleep mode
 When øw/4 is selected as the internal clock, timer F can operate in watch mode, subactive mode, and subsleep mode.
- Use of module standby mode enables this module to be placed in standby mode independently when not used.

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2. Block diagram

Figure 9.2 shows a block diagram of timer F.

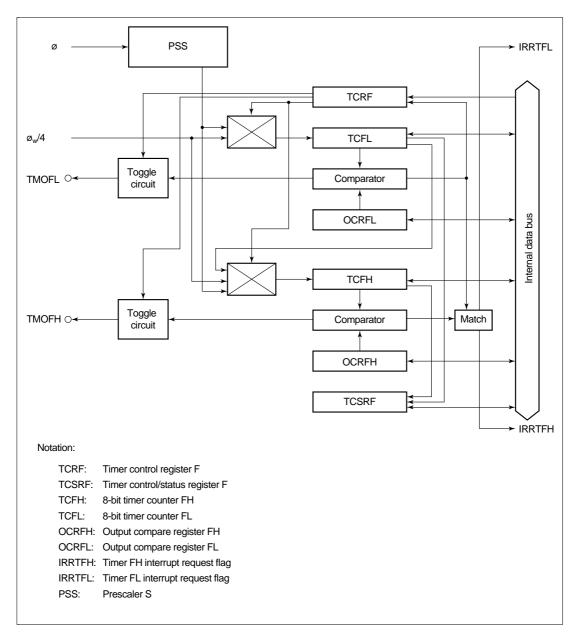


Figure 9.2 Block Diagram of Timer F

3. Pin configuration

Table 9.4 shows the timer F pin configuration.

Table 9.4Pin Configuration

Name	Abbrev.	I/O	Function
Timer FH output	TMOFH	Output	Timer FH toggle output pin
Timer FL output	TMOFL	Output	Timer FL toggle output pin

4. Register configuration

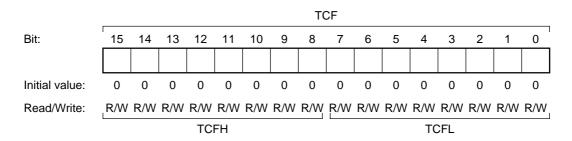
Table 9.5 shows the register configuration of timer F.

Table 9.5Timer F Registers

Name	Abbrev.	R/W	Initial Value	Address
Timer control register F	TCRF	W	H'00	H'FFB6
Timer control/status register F	TCSRF	R/W	H'00	H'FFB7
8-bit timer counter FH	TCFH	R/W	H'00	H'FFB8
8-bit timer counter FL	TCFL	R/W	H'00	H'FFB9
Output compare register FH	OCRFH	R/W	H'FF	H'FFBA
Output compare register FL	OCRFL	R/W	H'FF	H'FFBB
Clock stop register 1	CKSTPR1	R/W	H'FF	H'FFFA

9.3.2 Register Descriptions

16-bit timer counter (TCF)
 8-bit timer counter (TCFH)
 8-bit timer counter (TCFL)



TCF is a 16-bit read/write up-counter configured by cascaded connection of 8-bit timer counters TCFH and TCFL. In addition to the use of TCF as a 16-bit counter with TCFH as the upper 8 bits and TCFL as the lower 8 bits, TCFH and TCFL can also be used as independent 8-bit counters.

TCFH and TCFL can be read and written by the CPU, but when they are used in 16-bit mode, data transfer to and from the CPU is performed via a temporary register (TEMP). For details of TEMP, see 9.3.3, CPU Interface.

TCFH and TCFL are each initialized to H'00 upon reset.

a. 16-bit mode (TCF)

When CKSH2 is cleared to 0 in TCRF, TCF operates as a 16-bit counter. The TCF input clock is selected by bits CKSL2 to CKSL0 in TCRF.

TCF can be cleared in the event of a compare match by means of CCLRH in TCSRF.

When TCF overflows from H'FFFF to H'0000, OVFH is set to 1 in TCSRF. If OVIEH in TCSRF is 1 at this time, IRRTFH is set to 1 in IRR2, and if IENTFH in IENR2 is 1, an interrupt request is sent to the CPU.

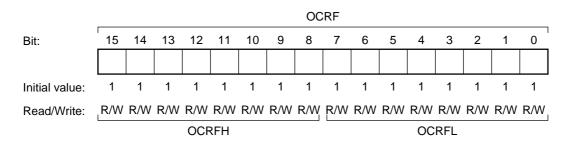
b. 8-bit mode (TCFL/TCFH)

When CKSH2 is set to 1 in TCRF, TCFH and TCFL operate as two independent 8-bit counters. The TCFH (TCFL) input clock is selected by bits CKSH2 to CKSH0 (CKSL2 to CKSL0) in TCRF.

TCFH (TCFL) can be cleared in the event of a compare match by means of CCLRH (CCLRL) in TCSRF.

When TCFH (TCFL) overflows from H'FF to H'00, OVFH (OVFL) is set to 1 in TCSRF. If OVIEH (OVIEL) in TCSRF is 1 at this time, IRRTFH (IRRTFL) is set to 1 in IRR2, and if IENTFH (IENTFL) in IENR2 is 1, an interrupt request is sent to the CPU.

16-bit output compare register (OCRF)
 8-bit output compare register (OCRFH)
 8-bit output compare register (OCRFL)



OCRF is a 16-bit read/write register composed of the two registers OCRFH and OCRFL. In addition to the use of OCRF as a 16-bit register with OCRFH as the upper 8 bits and OCRFL as the lower 8 bits, OCRFH and OCRFL can also be used as independent 8-bit registers.

OCRFH and OCRFL can be read and written by the CPU, but when they are used in 16-bit mode, data transfer to and from the CPU is performed via a temporary register (TEMP). For details of TEMP, see 9.3.3, CPU Interface.

OCRFH and OCRFL are each initialized to H'FF upon reset.

a. 16-bit mode (OCRF)

When CKSH2 is cleared to 0 in TCRF, OCRF operates as a 16-bit register. OCRF contents are constantly compared with TCF, and when both values match, CMFH is set to 1 in TCSRF. At the same time, IRRTFH is set to 1 in IRR2. If IENTFH in IENR2 is 1 at this time, an interrupt request is sent to the CPU.

Toggle output can be provided from the TMOFH pin by means of compare matches, and the output level can be set (high or low) by means of TOLH in TCRF.

b. 8-bit mode (OCRFH/OCRFL)

When CKSH2 is set to 1 in TCRF, OCRFH and OCRFL operate as two independent 8-bit registers. OCRFH contents are compared with TCFH, and OCRFL contents are with TCFL. When the OCRFH (OCRFL) and TCFH (TCFL) values match, CMFH (CMFL) is set to 1 in TCSRF. At the same time, IRRTFH (IRRTFL) is set to 1 in IRR2. If IENTFH (IENTFL) in IENR2 is 1 at this time, an interrupt request is sent to the CPU.

Toggle output can be provided from the TMOFH pin (TMOFL pin) by means of compare matches, and the output level can be set (high or low) by means of TOLH (TOLL) in TCRF.

3. Timer control register F (TCRF)

Bit:	7	6	5	4	3	2	1	0
	TOLH	CKSH2	CKSH1	CKSH0	TOLL	CKSL2	CKSL1	CKSL0
Initial value:	0	0	0	0	0	0	0	0
Read/Write:	W	W	W	W	W	W	W	W

TCRF is an 8-bit write-only register that switches between 16-bit mode and 8-bit mode, selects the input clock from among four internal clock sources and sets the output level of the TMOFH and TMOFL pins.

TCRF is initialized to H'00 upon reset.

Bit 7: Toggle output level H (TOLH)

Bit 7 sets the TMOFH pin output level. The output level is effective immediately after this bit is written.

Bit 7 TOLH	Description	
0	Low level	(initial value)
1	High level	

Bits 6 to 4: Clock select H (CKSH2 to CKSH0)

Bits 6 to 4 select the clock input to TCFH from among four internal clock sources or TCFL overflow.

Bit 6 CKSH2	Bit 5 CKSH1	Bit 4 CKSH0	Description	
0	0	0	16-bit mode, counting on TCFL overflow signal	(initial value)
0	0	1	_	
0	1	0	_	
0	1	1	Use prohibited	
1	0	0	Internal clock: counting on ø/32	
1	0	1	Internal clock: counting on ø/16	
1	1	0	Internal clock: counting on ø/4	
1	1	1	Internal clock: counting on øw/4	

Bit 3: Toggle output level L (TOLL)

Bit 3 sets the TMOFL pin output level. The output level is effective immediately after this bit is written.

Bit 3 TOLL	Description	
0	Low level	(initial value)
1	High level	

Bits 2 to 0: Clock select L (CKSL2 to CKSL0)

Bits 2 to 0 select the clock input to TCFL from among four internal clock sources or external event input.

Bit 2 CKSL2	Bit 1 CKSL1	Bit 0 CKSL0	Description	
0	0	0	Non-operational	(initial value)
0	0	1	Use prohibited	
0	1	0		
0	1	1	_	
1	0	0	Internal clock: counting on ø/32	
1	0	1	Internal clock: counting on ø/16	
1	1	0	Internal clock: counting on ø/4	
1	1	1	Internal clock: counting on øw/4	

4. Timer control/status register F (TCSRF)

Bit:	7	6	5	4	3	2	1	0
	OVFH	CMFH	OVIEH	CCLRH	OVFL	CMFL	OVIEL	CCLRL
Initial value:	0	0	0	0	0	0	0	0
Read/Write:	R/(W)*	R/(W)*	R/W	R/W	R/(W)*	R/(W)*	R/W	R/W

Note: * Bits 7, 6, 3, and 2 can only be written with 0, for flag clearing.

TCSRF is an 16-bit read/write register that performs counter clear selection, overflow flag setting, and compare match flag setting, and controls enabling of overflow interrupt requests.

TCSRF is initialized to H'00 upon reset.

Bit 7: Timer overflow flag H (OVFH)

Bit 7 is a status flag indicating that TCFH has overflowed from H'FF to H'00. This flag is set by hardware and cleared by software. It cannot be set by software.

Bit 7 OVFH	Description	
0	Clearing conditions: After reading OVFH = 1, cleared by writing 0 to OVFH	(initial value)
1	Setting conditions: Set when TCFH overflows from H'FF to H'00	

Bit 6: Compare match flag H (CMFH)

Bit 6 is a status flag indicating that TCFH has matched OCRFH. This flag is set by hardware and cleared by software. It cannot be set by software.

Bit 6 CMFH	Description	
0	Clearing conditions: After reading CMFH = 1, cleared by writing 0 to CMFH	(initial value)
1	Setting conditions: Set when the TCFH value matches the OCRFH value	

Bit 5: Timer overflow interrupt enable H (OVIEH)

Bit 5 selects enabling or disabling of interrupt generation when TCFH overflows.

Bit 5 OVIEH	Description	
0	TCFH overflow interrupt request is disabled	(initial value)
1	TCFH overflow interrupt request is enabled	

Bit 4: Counter clear H (CCLRH)

In 16-bit mode, bit 4 selects whether TCF is cleared when TCF and OCRF match.

In 8-bit mode, bit 4 selects whether TCFH is cleared when TCFH and OCRFH match.

Bit 4 CCLRH	Description	
0	16-bit mode: TCF clearing by compare match is disabled 8-bit mode: TCFH clearing by compare match is disabled	(initial value)
1	16-bit mode: TCF clearing by compare match is enabled 8-bit mode: TCFH clearing by compare match is enabled	

Bit 3: Timer overflow flag L (OVFL)

Bit 3 is a status flag indicating that TCFL has overflowed from H'FF to H'00. This flag is set by hardware and cleared by software. It cannot be set by software.

Bit 3 OVFL	Description	
0	Clearing conditions: After reading OVFL = 1, cleared by writing 0 to OVFL	(initial value)
1	Setting conditions: Set when TCFL overflows from H'FF to H'00	

Bit 2: Compare match flag L (CMFL)

Bit 2 is a status flag indicating that TCFL has matched OCRFL. This flag is set by hardware and cleared by software. It cannot be set by software.

Bit 2 CMFL	Description	
0	Clearing conditions: After reading CMFL = 1, cleared by writing 0 to CMFL	(initial value)
1	Setting conditions: Set when the TCFL value matches the OCRFL value	

Bit 1: Timer overflow interrupt enable L (OVIEL)

Bit 1 selects enabling or disabling of interrupt generation when TCFL overflows.

Bit 1 OVIEL	Description	
0	TCFL overflow interrupt request is disabled	(initial value)
1	TCFL overflow interrupt request is enabled	

Bit 0: Counter clear L (CCLRL)

Bit 0 selects whether TCFL is cleared when TCFL and OCRFL match.

Bit 0 CCLRL	Description	
0	TCFL clearing by compare match is disabled	(initial value)
1	TCFL clearing by compare match is enabled	

5. Clock stop register 1 (CKSTPR1)

Bit:	7	6	5	4	3	2	1	0
	—	—	S32CKSTP	ADCKSTP	—	TFCKSTP	—	TACKSTP
Initial value:	1	1	1	1	1	1	1	1
Read/Write:	—		R/W	R/W	_	R/W	_	R/W

CKSTPR1 is an 8-bit read/write register that performs module standby mode control for peripheral modules. Only the bit relating to timer F is described here. For details of the other bits, see the sections on the relevant modules.

Bit 2: Timer F module standby mode control (TFCKSTP)

Bit 2 controls setting and clearing of module standby mode for timer F.

TFCKSTP	Description	
0	Timer F is set to module standby mode	
1	Timer F module standby mode is cleared	(initial value)

9.3.3 CPU Interface

TCF and OCRF are 16-bit read/write registers, but the CPU is connected to the on-chip peripheral modules by an 8-bit data bus. When the CPU accesses these registers, it therefore uses an 8-bit temporary register (TEMP).

In 16-bit mode, TCF read/write access and OCRF write access must be performed 16 bits at a time (using two consecutive byte-size MOV instructions), and the upper byte must be accessed before the lower byte. Data will not be transferred correctly if only the upper byte or only the lower byte is accessed.

In 8-bit mode, there are no restrictions on the order of access.

1. Write access

Write access to the upper byte results in transfer of the upper-byte write data to TEMP. Next, write access to the lower byte results in transfer of the data in TEMP to the upper register byte, and direct transfer of the lower-byte write data to the lower register byte.

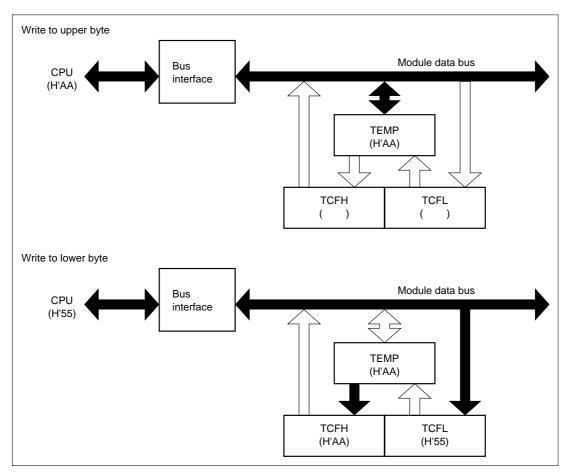


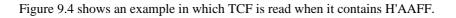
Figure 9.3 shows an example in which H'AA55 is written to TCF.

Figure 9.3 Write Access to TCR (CPU \rightarrow TCF)

2. Read access

In access to TCF, when the upper byte is read the upper-byte data is transferred directly to the CPU and the lower-byte data is transferred to TEMP. Next, when the lower byte is read, the lower-byte data in TEMP is transferred to the CPU.

In access to OCRF, when the upper byte is read the upper-byte data is transferred directly to the CPU. When the lower byte is read, the lower-byte data is transferred directly to the CPU.



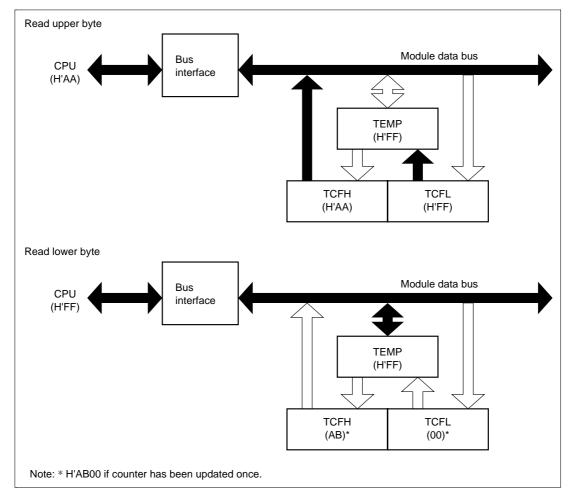


Figure 9.4 Read Access to TCF (TCF \rightarrow CPU)

9.3.4 Operation

Timer F is a 16-bit counter that increments on each input clock pulse. The timer F value is constantly compared with the value set in output compare register F, and the counter can be cleared, an interrupt requested, or port output toggled, when the two values match. Timer F can also function as two independent 8-bit timers.

1. Timer F operation

Timer F has two operating modes, 16-bit timer mode and 8-bit timer mode. The operation in each of these modes is described below.

a. Operation in 16-bit timer mode

When CKSH2 is cleared to 0 in timer control register F (TCRF), timer F operates as a 16-bit timer.

Following a reset, timer counter F (TCF) is initialized to H'0000, output compare register F (OCRF) to H'FFFF, and timer control register F (TCRF) and timer control/status register F (TCSRF) to H'00.

The timer F operating clock can be selected from three internal clocks output by prescaler S by means of bits CKSL2 to CKSL0 in TCRF.

OCRF contents are constantly compared with TCF, and when both values match, CMFH is set to 1 in TCSRF. If IENTFH in IENR2 is 1 at this time, an interrupt request is sent to the CPU, and at the same time, TMOFH pin output is toggled. If CCLRH in TCSRF is 1, TCF is cleared. TMOFH pin output can also be set by TOLH in TCRF.

When TCF overflows from H'FFFF to H'0000, OVFH is set to 1 in TCSRF. If OVIEH in TCSRF and IENTFH in IENR2 are both 1, an interrupt request is sent to the CPU.

b. Operation in 8-bit timer mode

When CKSH2 is set to 1 in TCRF, TCF operates as two independent 8-bit timers, TCFH and TCFL. The TCFH/TCFL input clock is selected by CKSH2 to CKSH0/CKSL2 to CKSL0 in TCRF.

When the OCRFH/OCRFL and TCFH/TCFL values match, CMFH/CMFL is set to 1 in TCSRF. If IENTFH/IENTFL in IENR2 is 1, an interrupt request is sent to the CPU, and at the same time, TMOFH pin/TMOFL pin output is toggled. If CCLRH/CCLRL in TCSRF is 1, TCFH/TCFL is cleared. TMOFH pin/TMOFL pin output can also be set by TOLH/TOLL in TCRF.

When TCFH/TCFL overflows from H'FF to H'00, OVFH/OVFL is set to 1 in TCSRF. If OVIEH/OVIEL in TCSRF and IENTFH/IENTFL in IENR2 are both 1, an interrupt request is sent to the CPU.

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2. TCF increment timing

TCF is incremented by clock input (internal clock input). Bits CKSH2 to CKSH0 or CKSL2 to CKSL0 in TCRF select one of four internal clock sources ($\phi/32$, $\phi/16$, $\phi/4$, or $\phi w/4$) created by dividing the system clock (ϕ or ϕw).

3. TMOFH/TMOFL output timing

In TMOFH/TMOFL output, the value set in TOLH/TOLL in TCRF is output. The output is toggled by the occurrence of a compare match. Figure 9.5 shows the output timing.

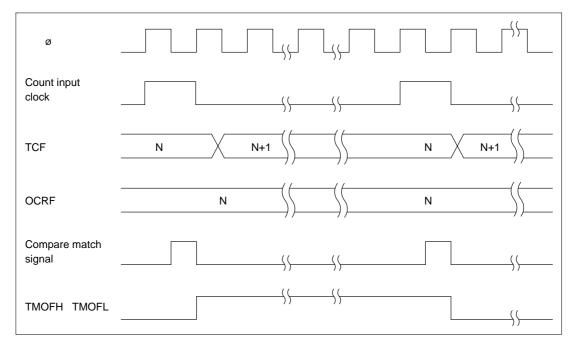


Figure 9.5 TMOFH/TMOFL Output Timing

4. TCF clear timing

TCF can be cleared by a compare match with OCRF.

5. Timer overflow flag (OVF) set timing

OVF is set to 1 when TCF overflows from H'FFFF to H'0000.

6. Compare match flag set timing

The compare match flag (CMFH or CMFL) is set to 1 when the TCF and OCRF values match. The compare match signal is generated in the last state during which the values match (when TCF is updated from the matching value to a new value). When TCF matches OCRF, the compare match signal is not generated until the next counter clock.

7. Timer F operation modes

Timer F operation modes are shown in table 9.6.

Table 9.6Timer F Operation Modes

					Sub-	Sub-		Module
Operation Mode	Reset	Active	Sleep	Watch	active	sleep	Standby	Standby
TCF	Reset	Functions	Functions	Functions/ Halted*	Functions/ Halted*	Functions/ Halted*	Halted	Halted
OCRF	Reset	Functions	Held	Held	Functions	Held	Held	Held
TCRF	Reset	Functions	Held	Held	Functions	Held	Held	Held
TCSRF	Reset	Functions	Held	Held	Functions	Held	Held	Held

Note: * When $ø_w/4$ is selected as the TCF internal clock in active mode or sleep mode, since the system clock and internal clock are mutually asynchronous, synchronization is maintained by a synchronization circuit. This results in a maximum count cycle error of 1/ø (s). When the counter is operated in subactive mode, watch mode, or subsleep mode, $ø_w/4$ must be selected as the internal clock. The counter will not operate if any other internal clock is selected.

9.3.5 Application Notes

The following types of contention and operation can occur when timer F is used.

1. 16-bit timer mode

In toggle output, TMOFH pin output is toggled when all 16 bits match and a compare match signal is generated. If a TCRF write by a MOV instruction and generation of the compare match signal occur simultaneously, TOLH data is output to the TMOFH pin as a result of the TCRF write. TMOFL pin output is unstable in 16-bit mode, and should not be used; the TMOFL pin should be used as a port pin.

If an OCRFL write and compare match signal generation occur simultaneously, the compare match signal is invalid. However, if the written data and the counter value match, a compare match signal will be generated at that point. As the compare match signal is output in synchronization with the TCFL clock, a compare match will not result in compare match signal generation if the clock is stopped.

Compare match flag CMFH is set when all 16 bits match and a compare match signal is generated. Compare match flag CMFL is set if the setting conditions for the lower 8 bits are satisfied.

When TCF overflows, OVFH is set. OVFL is set if the setting conditions are satisfied when the lower 8 bits overflow. If a TCFL write and overflow signal output occur simultaneously, the overflow signal is not output.

- 2. 8-bit timer mode
- a. TCFH, OCRFH

In toggle output, TMOFH pin output is toggled when a compare match occurs. If a TCRF write by a MOV instruction and generation of the compare match signal occur simultaneously, TOLH data is output to the TMOFH pin as a result of the TCRF write.

If an OCRFH write and compare match signal generation occur simultaneously, the compare match signal is invalid. However, if the written data and the counter value match, a compare match signal will be generated at that point. The compare match signal is output in synchronization with the TCFH clock.

If a TCFH write and overflow signal output occur simultaneously, the overflow signal is not output.

b. TCFL, OCRFL

In toggle output, TMOFL pin output is toggled when a compare match occurs. If a TCRF write by a MOV instruction and generation of the compare match signal occur simultaneously, TOLL data is output to the TMOFL pin as a result of the TCRF write.

If an OCRFL write and compare match signal generation occur simultaneously, the compare match signal is invalid. However, if the written data and the counter value match, a compare match signal will be generated at that point. As the compare match signal is output in synchronization with the TCFL clock, a compare match will not result in compare match signal generation if the clock is stopped.

If a TCFL write and overflow signal output occur simultaneously, the overflow signal is not output.

3. Clear timer FH, timer FL interrupt request flags (IRRTFH, IRRTFL), timer overflow flags H, L (OVFH, OVFL) and compare match flags H, L (CMFH, CMFL)

When ϕ w/4 is selected as the internal clock, "Interrupt factor generation signal" will be operated with ϕ w and the signal will be outputted with ϕ w width. And, "Overflow signal" and "Compare match signal" are controlled with 2 cycles of ϕ w signals. Those signals are outputted with 2 cycles width of ϕ w (figure 9.6)

In active (high-speed, medium-speed) mode, even if you cleared interrupt request flag during the term of validity of "Interrupt factor generation signal", same interrupt request flag is set. (figure 9.6.1) And, you cannot be cleared timer overflow flag and compare match flag during the term of validity of "Overflow signal" and "Compare match signal".

For interrupt request flag is set right after interrupt request is cleared, interrupt process to one time timer FH, timer FL interrupt might be repeated. (figure 9.6 2) Therefore, to definitely clear interrupt request flag in active (high-speed, medium-speed) mode, clear should be processed after the time that calculated with below (1) formula. And, to definitely clear timer overflow flag and compare match flag, clear should be processed after read timer control status register F (TCSRF) after the time that calculated with below (1) formula. For ST of (1) formula, please substitute the longest number of execution states in used instruction. (10 states of RTE instruction when MULXU, DIVXU instruction is not used, 14 states when MULXU, DIVXU instruction is used) In subactive mode, there are not limitation for interrupt request flag, timer overflow flag, and compare match flag clear.

The term of validity of "Interrupt factor generation signal"

= 1 cycle of ϕ w + waiting time for completion of executing instruction + interrupt time synchronized with $\phi = 1/\phi$ w + ST × (1/ ϕ) + (2/ ϕ) (second).....(1)

ST: Executing number of execution states

Method 1 is recommended to operate for time efficiency.

Method 1

- 1. Prohibit interrupt in interrupt handling routine (set IENFH, IENFL to 0).
- 2. After program process returned normal handling, clear interrupt request flags (IRRTFH, IRRTFL) after more than that calculated with (1) formula.

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- 3. After read timer control status register F (TCSRF), clear timer overflow flags (OVFH, OVFL) and compare match flags (CMFH, CMFL).
- 4. Operate interrupt permission (set IENFH, IENFL to 1).

Method 2

- 1. Set interrupt handling routine time to more than time that calculated with (1) formula.
- 2. Clear interrupt request flags (IRRTFH, IRRTFL) at the end of interrupt handling routine.
- 3. After read timer control status register F (TCSRF), clear timer overflow flags (OVFH, OVFL) and compare match flags (CMFH, CMFL).

All above attentions are also applied in 16-bit mode and 8-bit mode.

	Inte flag I	rrupt request clear f 2	nterrupt request lag clear l	
Program process	 Interrupt	Interrupt	Normal	
øw _				
Interrupt factor – generation signal (Internal signal, nega-active)			1 	
Overflow signal, Compare match signal (Internal signal, nega-active)			- 	
Interrupt request flag (IRRTFH, IRRTFL) [—]	 1		1	

Figure 9.6 Clear Interrupt Request Flag when Interrupt Factor Generation Signal is Valid

4. Timer counter (TCF) read/write

When ϕ w/4 is selected as the internal clock in active (high-speed, medium-speed) mode, write on TCF is impossible. And, when read TCF, as the system clock and internal clock are mutually asynchronous, TCF synchronizes with synchronization circuit. This results in a maximum TCF read value error of ±1.

When read/write TCF in active (high-speed, medium-speed) mode is needed, please select internal clock except for øw/4 before read/write.

In subactive mode, even øw/4 is selected as the internal clock, normal read/write TCF is possible.

9.4 Asynchronous Event Counter (AEC)

9.4.1 Overview

The asynchronous event counter is incremented by external event clock or internal clock input.

1. Features

Features of the asynchronous event counter are given below.

• Can count asynchronous events

Can count external events input asynchronously without regard to the operation of base clocks ϕ and ϕ_{SUB} .

The counter has a 16-bit configuration, enabling it to count up to $65536 (2^{16})$ events.

- Can also be used as two independent 8-bit event counter channels.
- Can be used as single-channel independent 16-bit event counter.
- Event/clock input is enabled only when IRQAEC is high or event counter PWM output (IECPWM) is high.
- Both edge sensing can be used for IRQAEC or event counter PWM output (IECPWM) interrupts. When the asynchronous counter is not used, independent interrupt function use is possible.
- When an event counter PWM is used, event clock input enabling/disabling can be performed automatically in a fixed cycle.
- External event input or a prescaler output clock can be selected by software for the ECH and ECL clock sources. Ø/2, Ø/4, or Ø/8 can be selected as the prescaler output clock.
- Both edge counting is possible for AEVL and AEVH.
- Counter resetting and halting of the count-up function controllable by software
- Automatic interrupt generation on detection of event counter overflow
- Use of module standby mode enables this module to be placed in standby mode independently when not used.

2. Block diagram

Figure 9.7 shows a block diagram of the asynchronous event counter.

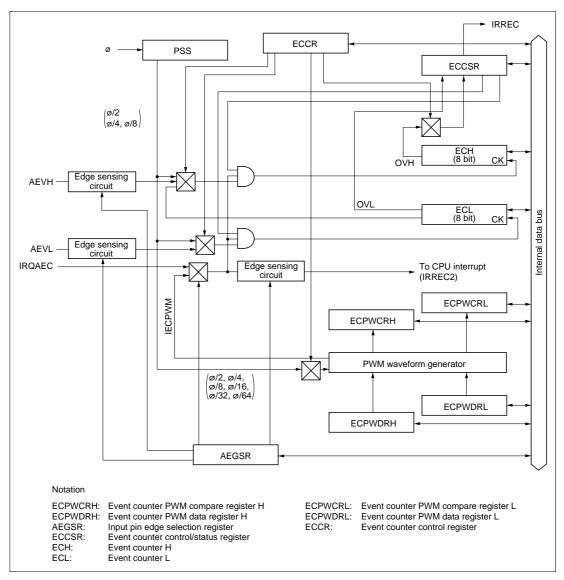


Figure 9.7 Block Diagram of Asynchronous Event Counter

3. Pin configuration

Table 9.7 shows the asynchronous event counter pin configuration.

Table 9.7Pin Configuration

Name	Abbrev.	I/O	Function
Asynchronous event input H	AEVH	Input	Event input pin for input to event counter H
Asynchronous event input L	AEVL	Input	Event input pin for input to event counter L
Event input enable interrupt input	IRQAEC	Input	Input pin for interrupt enabling event input

4. Register configuration

Table 9.8 shows the register configuration of the asynchronous event counter.

Table 9.8 Asynchronous Event Counter Registers

Name	Abbrev.	R/W	Initial Value	Address
Event counter PWM compare register H	ECPWCRH	R/W	H'FF	H'FF8C
Event counter PWM compare register L	ECPWCRL	R/W	H'FF	H'FF8D
Event counter PWM data register H	ECPWDRH	W	H'00	H'FF8E
Event counter PWM data register L	ECPWDRL	W	H'00	H'FF8F
Input pin edge selection register	AEGSR	R/W	H'00	H'FF92
Event counter control register	ECCR	R/W	H'00	H'FF94
Event counter control/status register	ECCSR	R/W	H'00	H'FF95
Event counter H	ECH	R	H'00	H'FF96
Event counter L	ECL	R	H'00	H'FF97
Clock stop register 2	CKSTP2	R/W	H'FF	H'FFFB

9.4.2 Register Configurations

1. Event counter PWM compare register H (ECPWCRH)

Bit	7	6	5	4	3	2	1	0
	ECPWCRH7	ECPWCRH6	ECPWCRH5	ECPWCRH4	ECPWCRH3	ECPWCRH2	ECPWCRH1	ECPWCRH0
Initial value	1	1	1	1	1	1	1	1
Read/Write	R/W							

Note: When ECPWME in AEGSR is 1, event counter PWM is operating and therefore ECPWCRH should not be modified.

When changing the conversion period, event counter PWM must be halted by clearing ECPWME to 0 in AEGSR before modifying ECPWCRH.

ECPWCRH is an 8-bit read/write register that sets the event counter PWM waveform conversion period.

2. Event counter PWM compare register L (ECPWCRL)

Bit	7	6	5	4	3	2	1	0
	ECPWCRL7	ECPWCRL6	ECPWCRL5	ECPWCRL4	ECPWCRL3	ECPWCRL2	ECPWCRL1	ECPWCRL0
Initial value	1	1	1	1	1	1	1	1
Read/Write	R/W							

Note: When ECPWME in AEGSR is 1, event counter PWM is operating and therefore ECPWCRL should not be modified.

When changing the conversion period, event counter PWM must be halted by clearing ECPWME to 0 in AEGSR before modifying ECPWCRL.

ECPWCRL is an 8-bit read/write register that sets the event counter PWM waveform conversion period.

3. Event counter PWM data register H (ECPWDRH)

Bit	7	6	5	4	3	2	1	0
	ECPWDRH7	ECPWDRH6	ECPWDRH5	ECPWDRH4	ECPWDRH3	ECPWDRH2	ECPWDRH1	ECPWDRH0
Initial value	0	0	0	0	0	0	0	0
Read/Write	W	W	W	W	W	W	W	W

Note: When ECPWME in AEGSR is 1, event counter PWM is operating and therefore ECPWDRH should not be modified. When changing the data, event counter PWM must be halted by clearing ECPWME to 0 in AEGSR before modifying ECPWDRH.

ECPWDRH is an 8-bit write-only register that controls event counter PWM waveform generator data.

4. Event counter PWM data register L (ECPWDRL)

Bit	7	6	5	4	3	2	1	0
	ECPWDRL7	ECPWDRL6	ECPWDRL5	ECPWDRL4	ECPWDRL3	ECPWDRL2	ECPWDRL1	ECPWDRL0
Initial value	0	0	0	0	0	0	0	0
Read/Write	W	W	W	W	W	W	W	W

Note: When ECPWME in AEGSR is 1, event counter PWM is operating and therefore ECPWDRL should not be modified. When changing the data, event counter PWM must be halted by clearing ECPWME to 0 in AEGSR before modifying ECPWDRL.

ECPWDRL is an 8-bit write-only register that controls event counter PWM waveform generator data.

5. Input pin edge selection register (AEGSR)

Bit	7	6	5	4	3	2	1	0
	AHEGS1	AHEGS0	ALEGS1	ALEGS0	AIEGS1	AIEGS0	ECPWME	—
Initial value	0	0	0	0	0	0	0	0
Read/Write	R/W	R/W						

AEGSR is an 8-bit read/write register that selects rising, falling, or both edge sensing for the AEVH, AEVL, and IRQAEC pins.

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Bits 7 and 6: AEC edge select H

Bits 7 and 6 select rising, falling, or both edge sensing for the AEVH pin.

Bit 7 AHEGS1	Bit 6 AHEGS0	Description	
0	0	Falling edge on AEVH pin is sensed	(initial value)
	1	Rising edge on AEVH pin is sensed	
1	0	Both edges on AEVH pin are sensed	
	1	Use prohibited	

Bits 5 and 4: AEC edge select L

Bits 5 and 4 select rising, falling, or both edge sensing for the AEVL pin.

Bit 5 ALEGS1	Bit 4 ALEGS0	Description	
0	0	Falling edge on AEVL pin is sensed	(initial value)
	1	Rising edge on AEVL pin is sensed	
1	0	Both edges on AEVL pin are sensed	
	1	Use prohibited	

Bits 3 and 2: IRQAEC edge select

Bits 3 and 2 select rising, falling, or both edge sensing for the IRQAEC pin.

Bit 3 AIEGS1	Bit 2 AIEGS0	Description	
0	0	Falling edge on IRQAEC pin is sensed	(initial value)
	1	Rising edge on IRQAEC pin is sensed	
1	0	Both edges on IRQAEC pin are sensed	
	1	Use prohibited	

Bit 1: Event counter PWM enable

Bit 1 controls enabling/disabling of event counter PWM and selection/deselection of IRQAEC.

Bit 1 ECPWME	Description	
0	AEC PWM halted, IRQAEC selected	(initial value)
1	AEC PWM operation enabled, IRQAEC deselected	

Bit 0: Reserved bit

Bit 0 is a readable/writable reserved bit. It is initialized to 0 by a reset.

Note: Do not set this bit to 1.

6. Event counter control register (ECCR)

Bit	7	6	5	4	3	2	1	0
	ACKH1	ACKH0	ACKL1	ACKL0	PWCK2	PWCK1	PWCK0	
Initial value	0	0	0	0	0	0	0	0
Read/Write	R/W	R/W						

ECCR performs counter input clock and IRQAEC/IECPWM control.

Bits 7 and 6: AEC clock select H (ACKH1, ACKH0)

Bits 7 and 6 select the clock used by ECH.

Bit 7 ACKH1	Bit 6 ACKH0	Description	
0	0	AEVH pin input	(initial value)
	1	ø/2	
1	0	ø/4	
	1	ø/8	

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Bits 5 and 4: AEC clock select L (ACKL1, ACKL0)

Bits 5 and 4 select the clock used by ECL.

Bit 5 ACKL1	Bit 4 ACKL0	Description	
0	0	AEVL pin input	(initial value)
	1	ø/2	
1	0	ø/4	
	1	ø/8	

Bits 3 to 1: Event counter PWM clock select (PWCK2, PWCK1, PWCK0)

Bits 3 to 1 select the event counter PWM clock.

Bit 3 PWCK2	Bit 2 PWCK1	Bit 1 PWCK0	Description	
0	0	0	ø/2	(initial value)
		1	ø/4	
	1	0	ø/8	
		1	ø/16	
1	*	0	ø/32	
		1	ø/64	

*: Don't care

Bit 0: Reserved bit

Bit 0 is a readable/writable reserved bit. It is initialized to 0 by a reset.

Note: Do not set this bit to 1.

7. Event counter control/status register (ECCSR)

Bit	7	6	5	4	3	2	1	0
	OVH	OVL	—	CH2	CUEH	CUEL	CRCH	CRCL
Initial Value	0	0	0	0	0	0	0	0
Read/Write	R/W*	R/W*	R/W	R/W	R/W	R/W	R/W	R/W

Note: * Bits 7 and 6 can only be written with 0, for flag clearing.

ECCSR is an 8-bit read/write register that controls counter overflow detection, counter resetting, and halting of the count-up function.

ECCSR is initialized to H'00 upon reset.

Bit 7: Counter overflow flag H (OVH)

Bit 7 is a status flag indicating that ECH has overflowed from H'FF to H'00. This flag is set when ECH overflows. It is cleared by software but cannot be set by software. OVH is cleared by reading it when set to 1, then writing 0.

When ECH and ECL are used as a 16-bit event counter with CH2 cleared to 0, OVH functions as a status flag indicating that the 16-bit event counter has overflowed from H'FFFF to H'0000.

Bit 7 OVH	Description	
0	ECH has not overflowed Clearing conditions: After reading OVH = 1, cleared by writing 0 to OVH	(initial value)
1	ECH has overflowed Setting conditions: Set when ECH overflows from H'FF to H'00	

Bit 6: Counter overflow flag L (OVL)

Bit 6 is a status flag indicating that ECL has overflowed from H'FF to H'00. This flag is set when ECL overflows. It is cleared by software but cannot be set by software. OVL is cleared by reading it when set to 1, then writing 0.

Bit 6 OVL	Description	
0	ECL has not overflowed Clearing conditions: After reading OVL = 1, cleared by writing 0 to OVL	(initial value)
1	ECL has overflowed Setting conditions: Set when ECL overflows from H'FF to H'00 while CH2 is set to 1	

Bit 5: Reserved bit

Bit 5 is reserved; it can be read and written, and is initialized to 0 upon reset.

Bit 4: Channel select (CH2)

Bit 4 selects whether ECH and ECL are used as a single-channel 16-bit event counter or as two independent 8-bit event counter channels. When CH2 is cleared to 0, ECH and ECL function as a 16-bit event counter which is incremented each time an event clock is input to the AEVL pin. In this case, the overflow signal from ECL is selected as the ECH input clock. When CH2 is set to 1, ECH and ECL function as independent 8-bit event counters which are incremented each time an event clock is input to the AEVH or AEVL pin, respectively.

Bit 4 CH2	Description				
0	ECH and ECL are used together as a single-channel 16-bit event counter				
	(initial value)				
1	ECH and ECL are used as two independent 8-bit event counter channels				

Bit 3: Count-up enable H (CUEH)

Bit 3 enables event clock input to ECH. When 1 is written to this bit, event clock input is enabled and increments the counter. When 0 is written to this bit, event clock input is disabled and the ECH value is held. The AEVH pin or the ECL overflow signal can be selected as the event clock source by bit CH2.

Bit 3 CUEH	Description	
0	ECH event clock input is disabled ECH value is held	(initial value)
1	ECH event clock input is enabled	

Bit 2: Count-up enable L (CUEL)

Bit 3 enables event clock input to ECL. When 1 is written to this bit, event clock input is enabled and increments the counter. When 0 is written to this bit, event clock input is disabled and the ECL value is held.

Bit 2 CUEL	Description	
0	ECL event clock input is disabled ECL value is held	(initial value)
1	ECL event clock input is enabled	

Bit 1: Counter reset control H (CRCH)

Bit 1 controls resetting of ECH. When this bit is cleared to 0, ECH is reset. When 1 is written to this bit, the counter reset is cleared and the ECH count-up function is enabled.

Bit 1 CRCH	Description	
0	ECH is reset	(initial value)
1	ECH reset is cleared and count-up function is enabled	

Bit 0: Counter reset control L (CRCL)

Bit 0 controls resetting of ECL. When this bit is cleared to 0, ECL is reset. When 1 is written to this bit, the counter reset is cleared and the ECL count-up function is enabled.

Bit 0 CRCL	Description	
0	ECL is reset	(initial value)
1	ECL reset is cleared and count-up function is enabled	

8. Event counter H (ECH)

Bit	7	6	5	4	3	2	1	0
	ECH7	ECH6	ECH5	ECH4	ECH3	ECH2	ECH1	ECH0
Initial Value	0	0	0	0	0	0	0	0
Read/Write	R	R	R	R	R	R	R	R

ECH is an 8-bit read-only up-counter that operates either as an independent 8-bit event counter or as the upper 8-bit up-counter of a 16-bit event counter configured in combination with ECL. Either the external asynchronous event AEVH pin or the overflow signal from lower 8-bit counter ECL can be selected as the input clock source. ECH can be cleared to H'00 by software, and is also initialized to H'00 upon reset.

9. Event counter L (ECL)

ECL is an 8-bit read-only up-counter that operates either as an independent 8-bit event counter or as the lower 8-bit up-counter of a 16-bit event counter configured in combination with ECH. The event clock from the external asynchronous event AEVL pin, or $\phi/2$, $\phi/4$, or $\phi/8$, is used as the input clock source. ECL can be cleared to H'00 by software, and is also initialized to H'00 upon reset.

Bit	7	6	5	4	3	2	1	0
	ECL7	ECL6	ECL5	ECL4	ECL3	ECL2	ECL1	ECL0
Initial Value	0	0	0	0	0	0	0	0
Read/Write	R	R	R	R	R	R	R	R

10. Clock stop register 2 (CKSTPR2)

Bit	7	6	5	4	3	2	1	0
		_	—	PW2CKSTP	AECKSTP	—	PW1CKSTP	LDCKSTP
Initial value	1	1	1	1	1	1	1	1
Read/Write	_	_	_	R/W	R/W	_	R/W	R/W

CKSTPR2 is an 8-bit read/write register that performs module standby mode control for peripheral modules. Only the bit relating to the asynchronous event counter is described here. For details of the other bits, see the sections on the relevant modules.

Bit 3: Asynchronous event counter module standby mode control (AECKSTP)

Bit 3 controls setting and clearing of module standby mode for the asynchronous event counter.

AECKSTP	Description	
0	Asynchronous event counter is set to module standby mode	
1	Asynchronous event counter module standby mode is cleared	(initial value)

9.4.3 Operation

1. 16-bit event counter operation

When bit CH2 is cleared to 0 in ECCSR, ECH and ECL operate as a 16-bit event counter.

Any of four input clock sources— $\emptyset/2$, $\emptyset/4$, $\emptyset/8$, or AEVL pin input—can be selected by means of bits ACKL1 and ACKL0 in ECCR.

When AEVL pin input is selected, input sensing is selected with bits ALEGS1 and ALEGS0.

The input clock is enabled only when IRQAEC is high or IECPWM is high. When IRQAEC is low or IECPWM is low, the input clock is not input to the counter, which therefore does not operate. Figure 9.8 shows an example of the software processing when ECH and ECL are used as a 16-bit event counter.

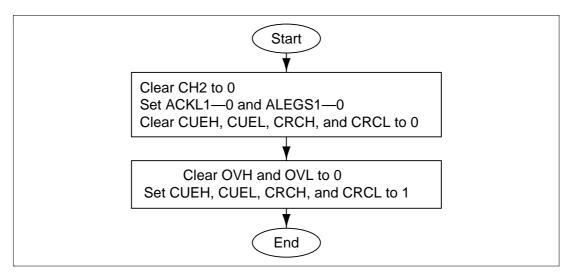


Figure 9.8 Example of Software Processing when Using ECH and ECL as 16-Bit Event Counter

As CH2 is cleared to 0 by a reset, ECH and ECL operate as a 16-bit event counter after a reset, and as ACKL1 and ACKL0 are cleared to 00, the operating clock is asynchronous event input from the AEVL pin (using falling edge sensing). When the next clock is input after the count value reaches H'FF in both ECH and ECL, ECH and ECL overflow from H'FFFF to H'0000, the OVH flag is set to 1 in ECCSR, the ECH and ECL count values each return to H'00, and counting up is restarted. When overflow occurs, the IRREC bit is set to 1 in IRR2. If the IENEC bit in IENR2 is 1 at this time, an interrupt request is sent to the CPU.

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2. 8-bit event counter operation

When bit CH2 is set to 1 in ECCSR, ECH and ECL operate as independent 8-bit event counters.

 $\phi/2$, $\phi/4$, $\phi/8$, or AEVH pin input can be selected as the input clock source for ECH by means of bits ACKH1 and ACKH0 in ECCR, and $\phi/2$, $\phi/4$, $\phi/8$, or AEVL pin input can be selected as the input clock source for ECL by means of bits ACKL1 and ACKL0 in ECCR.

Input sensing is selected with bits AHEGS1 and AHEGS0 when AEVH pin input is selected, and with bits ALEGS1 and ALEGS0 when AEVL pin input is selected.

The input clock is enabled only when IRQAEC is high or IECPWM is high. When IRQAEC is low or IECPWM is low, the input clock is not input to the counter, which therefore does not operate. Figure 9.9 shows an example of the software processing when ECH and ECL are used as 8-bit event counters.

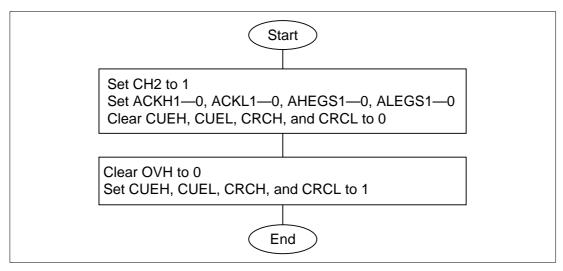


Figure 9.9 Example of Software Processing when Using ECH and ECL as 8-Bit Event Counters

ECH and ECL can be used as 8-bit event counters by carrying out the software processing shown in the example in figure 9.9. When the next clock is input after the ECH count value reaches H'FF, ECH overflows, the OVH flag is set to 1 in ECCSR, the ECH count value returns to H'00, and counting up is restarted. Similarly, when the next clock is input after the ECL count value reaches H'FF, ECL overflows, the OVL flag is set to 1 in ECCSR, the ECL count value returns to H'00, and counting up is restarted. When overflow occurs, the IRREC bit is set to 1 in IRR2. If the IENEC bit in IENR2 is 1 at this time, an interrupt request is sent to the CPU.

3. IRQAEC operation

When ECPWME in AEGSR is 0, the ECH and ECL input clocks are enabled only when IRQAEC is high. When IRQAEC is low, the input clocks are not input to the counters, and so ECH and ECL do not count. ECH and ECL count operations can therefore be controlled from outside by controlling IRQAEC. In this case, ECH and ECL cannot be controlled individually.

IRQAEC can also operate as an interrupt source. In this case the vector number is 6 and the vector addresses are H'000C and H'000D.

Interrupt enabling is controlled by IENEC2 in IENR1. When an IRQAEC interrupt is generated, IRR1 interrupt request flag IRREC2 is set to 1. If IENEC2 in IENR1 is set to 1 at this time, an interrupt request is sent to the CPU.

Rising, falling, or both edge sensing can be selected for the IRQAEC input pin with bits AIAGS1 and AIAGS0 in AEGSR.

4. Event counter PWM operation

When ECPWME in AEGSR is 1, the ECH and ECL input clocks are enabled only when event counter PWM output (IECPWM) is high. When IECPWM is low, the input clocks are not input to the counters, and so ECH and ECL do not count. ECH and ECL count operations can therefore be controlled cyclically from outside by controlling event counter PWM. In this case, ECH and ECL cannot be controlled individually.

IECPWM can also operate as an interrupt source. In this case the vector number is 6 and the vector addresses are H'000C and H'000D.

Interrupt enabling is controlled by IENEC2 in IENR1. When an IECPWM interrupt is generated, IRR1 interrupt request flag IRREC2 is set to 1. If IENEC2 in IENR1 is set to 1 at this time, an interrupt request is sent to the CPU.

Rising, falling, or both edge detection can be selected for IECPWM interrupt sensing with bits AIAGS1 and AIAGS0 in AEGSR.

Figure 9.10 and table 9.9 show examples of event counter PWM operation.

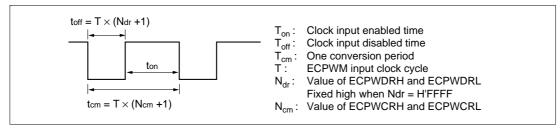


Figure 9.10 Event Counter Operation Waveform

Note: N_{dr} and N_{cm} above must be set so that $N_{dr} < N_{cm}$. If the settings do not satisfy this condition, do not set ECPWME to 1 in AEGSR.

Table 9.9 Examples of Event Counter PWM Operation

Conditions: $f_{osc} = 4$ MHz, $f_{\phi} = 2$ MHz, high-speed active mode, ECPWCR value (N_{cm}) = H'7A11, ECPWDR value (N_{dr}) = H'16E3

Clock Source Selection	Clock Source Cycle (T)*	ECPWMCR Value (N _{cm})	ECPWMDR Value (N _{dr})	$t_{off} = T \times (N_{dr} + 1)$	$t_{cm} = T \times (N_{cm} + 1)$	$\mathbf{t}_{on} = \mathbf{t}_{cm} - \mathbf{t}_{off}$
ø/2	1 µs	H'7A11	H'16E3 D'5859	5.86 ms	31.25 ms	25.39 ms
ø/4	2 µs	D'31249		11.72 ms	62.5 ms	50.78 ms
ø/8	4 µs	_		23.44 ms	125.0 ms	101.56 ms
ø/16	8 µs	_		46.88 ms	250.0 ms	203.12 ms
ø/32	16 µs	_		93.76 ms	500.0 ms	406.24 ms
ø/64	32 µs			187.52 ms	1000.0 ms	812.48 ms

Note: * $t_{\mbox{\tiny off}}$ minimum width

5. Clock Input Enable/Disable Function Operation

The clock input to the event counter can be controlled by the IRQAEC pin when ECPWME in AEGSR is 0, and by event counter PWM output IECPWM when ECPWME in AEGSR is 1. As this function forcibly terminates the clock input by each signal, a maximum error of one count will occur depending the IRQAEC or IECPWM timing.

Figure 9.11 shows an example of the operation of this function.

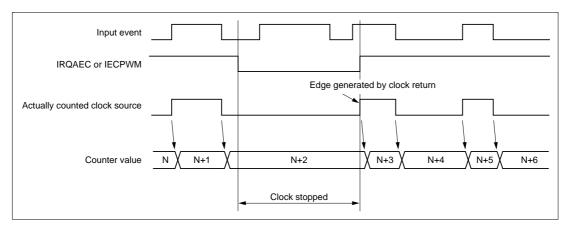


Figure 9.11 Example of Clock Control Operation

9.4.4 Asynchronous Event Counter Operation Modes

Asynchronous event counter operation modes are shown in table 9.10.

Table 9.10 Asynchronous Event Counter Operation Modes

Operation Mode	Reset	Active	Sleep	Watch	Subactive	Subsleep	Standby	Module Standby
AEGSR	Reset	Functions	Functions	Held*1	Functions	Functions	Held*1	Held
ECCR	Reset	Functions	Functions	Held*1	Functions	Functions	Held*1	Held
ECCSR	Reset	Functions	Functions	Held*1	Functions	Functions	Held*1	Held
ECH	Reset	Functions	Functions	Functions*1*2	Functions*2	Functions*2	Functions*1*2	Halted
ECL	Reset	Functions	Functions	Functions*1*2	Functions*2	Functions*2	Functions*1*2	Halted
IEQAEC	Reset	Functions	Functions	Held* ³	Functions	Functions	Held* ³	Held ^{*4}
Event counter PWM	Reset	Functions	Functions	Held	Held	Held	Held	Held

Notes: 1. When an asynchronous external event is input, the counter increments but the counter overflow H/L flags are not affected.

2. Operates when asynchronous mode external events are selected; halted and retained otherwise.

3. Clock control by IRQAEC operates, but interrupts do not.

4. As the clock is stopped in module standby mode, IRQAEC has no effect.

9.4.5 Application Notes

- 1. When reading the values in ECH and ECL, first clear bits CUEH and CUEL to 0 in ECCSR to prevent asynchronous event input to the counter. The correct value will not be returned if the event counter increments while being read.
- 2. Use a clock with a frequency of up to 16 MHz for input to the AEVH and AEVL pins, and ensure that the high and low widths of the clock are at least 30 ns. The duty cycle is immaterial.

Mode		Maximum AEVH/AEVL Pin Input Clock Frequency
Active (high-speed), sleep (high-speed)		16 MHz
Active (medium-speed), sleep (medium-speed) (ø/16)	2 · f _{OSC}
	(ø/32)	fosc
	(ø/64)	1/2 · f _{OSC}
f _{OSC} = 1 MHz to 4 MHz	(ø/128)	1/4 · f _{OSC}
Watch, subactive, subsleep, standby	(øw/2)	1000 kHz
	(øw/4)	500 kHz
øw = 32.768 kHz or 38.4 kHz	(øw/8)	250 kHz

- 3. When AEC uses with 16-bit mode, set CUEH in ECCSR to "1" first, set CRCH in ECCSR to "1" second, or set both CUEH and CRCH to "1" at same time before clock entry. While AEC is operating on 16-bit mode, do not change CUEH. Otherwise, ECH will be miscounted up.
- 4. When ECPWME in AEGSR is 1, event counter PWM is operating and therefore ECPWCRH, ECPWCRL, ECPWDRH, and ECPWDRL should not be modified. When changing the data, event counter PWM must be halted by clearing ECPWME to 0 in AEGSR before modifying these registers.
- 5. The event counter PWM data register and event counter PWM compare register must be set so that event counter PWM data register < event counter PWM compare register. If the settings do not satisfy this condition, do not set ECPWME to 1 in AEGSR.
- 6. As synchronization is established internally when an IRQAEC interrupt is generated, a maximum error of 1 tcyc will occur between clock halting and interrupt acceptance.

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Section 10 Serial Communication Interface

10.1 Overview

The H8/3802 Series is provided with one serial communication interface, SCI3.

Serial communication interface 3 (SCI3) can carry out serial data communication in either asynchronous or synchronous mode. It is also provided with a multiprocessor communication function that enables serial data to be transferred among processors.

10.1.1 Features

Features of SCI3 are listed below.

- Choice of asynchronous or synchronous mode for serial data communication
 - Asynchronous mode

Serial data communication is performed asynchronously, with synchronization provided character by character. In this mode, serial data can be exchanged with standard asynchronous communication LSIs such as a Universal Asynchronous Receiver/Transmitter (UART) or Asynchronous Communication Interface Adapter (ACIA). A multiprocessor communication function is also provided, enabling serial data communication among processors.

There is a choice of 16 data transfer formats.

Data length	7, 8, 5 bits
Stop bit length	1 or 2 bits
Parity	Even, odd, or none
Multiprocessor bit	1 or 0
Receive error detection	Parity, overrun, and framing errors
Break detection	Break detected by reading the RXD_{32} pin level directly when a framing error occurs

- Synchronous mode

Serial data communication is synchronized with a clock. In his mode, serial data can be exchanged with another LSI that has a synchronous communication function.

Data length	8 bits
Receive error detection	Overrun errors

• Full-duplex communication

Separate transmission and reception units are provided, enabling transmission and reception to be carried out simultaneously. The transmission and reception units are both double-buffered, allowing continuous transmission and reception.

- On-chip baud rate generator, allowing any desired bit rate to be selected
- Choice of an internal or external clock as the transmit/receive clock source
- Six interrupt sources: transmit end, transmit data empty, receive data full, overrun error, framing error, and parity error

10.1.2 Block diagram

Figure 10.1 shows a block diagram of SCI3.

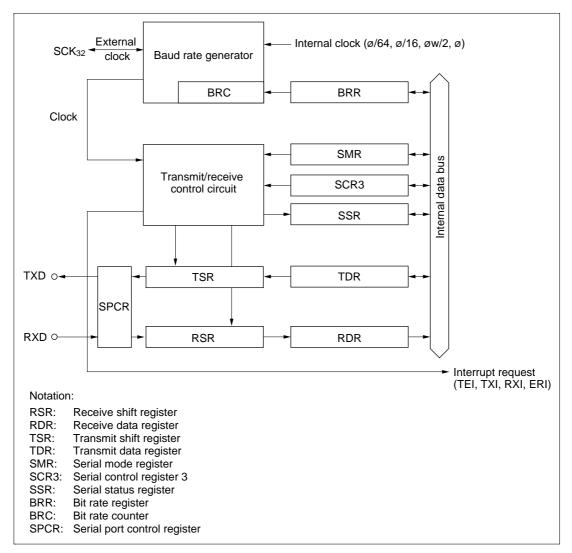


Figure 10.1 SCI3 Block Diagram

10.1.3 Pin configuration

Table 10.1 shows the SCI3 pin configuration.

Table 10.1 Pin Configuration

Name	Abbrev.	I/O	Function
SCI3 clock	SCK ₃₂	I/O	SCI3 clock input/output
SCI3 receive data input	RXD ₃₂	Input	SCI3 receive data input
SCI3 transmit data output	TXD ₃₂	Output	SCI3 transmit data output

10.1.4 Register configuration

Table 10.2 shows the SCI3 register configuration.

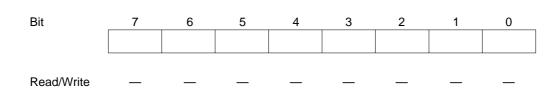
Table 10.2 Registers

Name	Abbrev.	R/W	Initial Value	Address
Serial mode register	SMR	R/W	H'00	H'FFA8
Bit rate register	BRR	R/W	H'FF	H'FFA9
Serial control register 3	SCR3	R/W	H'00	H'FFAA
Transmit data register	TDR	R/W	H'FF	H'FFAB
Serial data register	SSR	R/W	H'84	H'FFAC
Receive data register	RDR	R	H'00	H'FFAD
Transmit shift register	TSR	Protected		_
Receive shift register	RSR	Protected	_	_
Bit rate counter	BRC	Protected	_	
Clock stop register 1	CKSTPR1	R/W	H'FF	H'FFFA
Serial port control register	SPCR	R/W		H'FF91

10.2 Register Descriptions

Receive shift register (RSR)

10.2.1



RSR is a register used to receive serial data. Serial data input to RSR from the RXD_{32} pin is set in the order in which it is received, starting from the LSB (bit 0), and converted to parallel data. When one byte of data is received, it is transferred to RDR automatically.

RSR cannot be read or written directly by the CPU.

10.2.2 Receive data register (RDR)

Bit	7	6	5	4	3	2	1	0
	RDR7	RDR6	RDR5	RDR4	RDR3	RDR2	RDR1	RDR0
Initial value	0	0	0	0	0	0	0	0
Read/Write	R	R	R	R	R	R	R	R

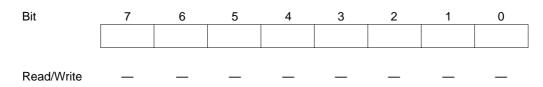
RDR is an 8-bit register that stores received serial data.

When reception of one byte of data is finished, the received data is transferred from RSR to RDR, and the receive operation is completed. RSR is then able to receive data. RSR and RDR are double-buffered, allowing consecutive receive operations.

RDR is a read-only register, and cannot be written by the CPU.

RDR is initialized to H'00 upon reset, and in standby, module standby or watch mode.

10.2.3 Transmit shift register (TSR)



TSR is a register used to transmit serial data. Transmit data is first transferred from TDR to TSR, and serial data transmission is carried out by sending the data to the TXD_{32} pin in order, starting from the LSB (bit 0). When one byte of data is transmitted, the next byte of transmit data is transferred to TDR, and transmission started, automatically. Data transfer from TDR to TSR is not performed if no data has been written to TDR (if bit TDRE is set to 1 in the serial status register (SSR)).

TSR cannot be read or written directly by the CPU.

10.2.4 Transmit data register (TDR)

Bit	7	6	5	4	3	2	1	0
	TDR7	TDR6	TDR5	TDR4	TDR3	TDR2	TDR1	TDR0
Initial value	1	1	1	1	1	1	1	1
Read/Write	R/W							

TDR is an 8-bit register that stores transmit data. When TSR is found to be empty, the transmit data written in TDR is transferred to TSR, and serial data transmission is started. Continuous transmission is possible by writing the next transmit data to TDR during TSR serial data transmission.

TDR can be read or written by the CPU at any time.

TDR is initialized to H'FF upon reset, and in standby, module standby, or watch mode.

10.2.5 Serial mode register (SMR)

Bit	7	6	5	4	3	2	1	0
	СОМ	CHR	PE	PM	STOP	MP	CKS1	CKS0
Initial value	0	0	0	0	0	0	0	0
Read/Write	R/W	R/W	R/W	R/W	R/W	R/W	R/W	R/W

SMR is an 8-bit register used to set the serial data transfer format and to select the clock source for the baud rate generator.

SMR can be read or written by the CPU at any time.

SMR is initialized to H'00 upon reset, and in standby, module standby, or watch mode.

Bit 7: Communication mode (COM)

Bit 7 selects whether SCI3 operates in asynchronous mode or synchronous mode.

Bit 7 COM	Description	
0	Asynchronous mode	(initial value)
1	Synchronous mode	

Bit 6: Character length (CHR)

Bit 6 selects either 7 or 8 bits as the data length to be used in asynchronous mode. In synchronous mode the data length is always 8 bits, irrespective of the bit 6 setting.

Bit 6 CHR	Description	
0	8-bit data/5-bit data*2	(initial value)
1	7-bit data ^{*1} /5-bit data ^{*2}	
Notes:	1. When 7-bit data is selected, the MSB (bit 7) of TDR is not transmitted.	

 When 5-bit data is selected, set both PE and MP to 1. The three most significant bits (bits 7, 6, and 5) of TDR are not transmitted.

Bit 5: Parity enable (PE)

Bit 5 selects whether a parity bit is to be added during transmission and checked during reception in asynchronous mode. In synchronous mode parity bit addition and checking is not performed, irrespective of the bit 5 setting.

Bit 5 PE	Description	
0	Parity bit addition and checking disabled*2	(initial value)
1	Parity bit addition and checking enabled ^{*1/*2}	
Notes: 1.	When PE is set to 1, even or odd parity, as designated by bit PM, is a data before it is sent, and the received parity bit is checked against the designated by bit PM.	

2. For the case where 5-bit data is selected, see table 10.11.

Bit 4: Parity mode (PM)

Bit 4 selects whether even or odd parity is to be used for parity addition and checking. The PM bit setting is only valid in asynchronous mode when bit PE is set to 1, enabling parity bit addition and checking. The PM bit setting is invalid in synchronous mode, and in asynchronous mode if parity bit addition and checking is disabled.

Bit 4 PM	Description	
0	Even parity ^{*1}	(initial value)
1	Odd parity ^{*2}	
Notes:	 When even parity is selected, a parity bit is added in a number of 1 bits in the transmit data plus the parity bit 	

number of 1 bits in the transmit data plus the parity bit is an even number; in reception, a check is carried out to confirm that the number of 1 bits in the receive data plus the parity bit is an even number.2. When odd parity is selected, a parity bit is added in transmission so that the total

2. When our party is selected, a party bit is added in transmission so that the total number of 1 bits in the transmit data plus the parity bit is an odd number; in reception, a check is carried out to confirm that the number of 1 bits in the receive data plus the parity bit is an odd number.

Bit 3: Stop bit length (STOP)

Bit 3 selects 1 bit or 2 bits as the stop bit length in asynchronous mode. The STOP bit setting is only valid in asynchronous mode. When synchronous mode is selected the STOP bit setting is invalid since stop bits are not added.

STOP	Description	
0	1 stop bit ^{*1}	(initial value)
1	2 stop bits ^{*2}	

Notes: 1. In transmission, a single 1 bit (stop bit) is added at the end of a transmit character.

2. In transmission, two 1 bits (stop bits) are added at the end of a transmit character.

In reception, only the first of the received stop bits is checked, irrespective of the STOP bit setting. If the second stop bit is 1 it is treated as a stop bit, but if 0, it is treated as the start bit of the next transmit character.

Bit 2: Multiprocessor mode (MP)

Bit 2 enables or disables the multiprocessor communication function. When the multiprocessor communication function is enabled, the parity settings in the PE and PM bits are invalid. The MP bit setting is only valid in asynchronous mode. When synchronous mode is selected the MP bit should be set to 0. For details on the multiprocessor communication function, see 10.3.4, Multiprocessor Communication Function.

Bit 2 MP	Description				
0	Multiprocessor communication function disabled*	(initial value)			
1	Multiprocessor communication function enabled*				

Note: * For the case where 5-bit data is selected, see table 10.11.

Bits 1 and 0: Clock select 1, 0 (CKS1, CKS0)

Bits 1 and 0 choose $\phi/64$, $\phi/16$, $\phi w/2$, or ϕ as the clock source for the baud rate generator.

For the relation between the clock source, bit rate register setting, and baud rate, see 8, Bit rate register (BRR).

Bit 1 CKS1	Bit 0 CKS0	Description	
0	0	ø clock	(initial value)
0	1	ø w/2 clock ^{*1} /ø w clock ^{*2}	
1	0	ø/16 clock	
1	1	ø/64 clock	

Notes: 1. ø w/2 clock in active (medium-speed/high-speed) mode and sleep mode

2. ø w clock in subactive mode and subsleep mode

3. In subactive or subsleep mode, SCI3 can be operated when CPU clock is øw/2 only.

10.2.6 Serial control register 3 (SCR3)

Bit	7	6	5	4	3	2	1	0
	TIE	RIE	TE	RE	MPIE	TEIE	CKE1	CKE0
Initial value	0	0	0	0	0	0	0	0
Read/Write	R/W	R/W	R/W	R/W	R/W	R/W	R/W	R/W

SCR3 is an 8-bit register for selecting transmit or receive operation, the asynchronous mode clock output, interrupt request enabling or disabling, and the transmit/receive clock source.

SCR3 can be read or written by the CPU at any time.

SCR3 is initialized to H'00 upon reset, and in standby, module standby or watch mode.

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Bit 7: Transmit interrupt enable (TIE)

Bit 7 selects enabling or disabling of the transmit data empty interrupt request (TXI) when transmit data is transferred from the transmit data register (TDR) to the transmit shift register (TSR), and bit TDRE in the serial status register (SSR) is set to 1.

TXI can be released by clearing bit TDRE or bit TIE to 0.

Bit 7		
TIE	Description	
0	Transmit data empty interrupt request (TXI) disabled	(initial value)
1	Transmit data empty interrupt request (TXI) enabled	

Bit 6: Receive interrupt enable (RIE)

Bit 6 selects enabling or disabling of the receive data full interrupt request (RXI) and the receive error interrupt request (ERI) when receive data is transferred from the receive shift register (RSR) to the receive data register (RDR), and bit RDRF in the serial status register (SSR) is set to 1. There are three kinds of receive error: overrun, framing, and parity.

RXI and ERI can be released by clearing bit RDRF or the FER, PER, or OER error flag to 0, or by clearing bit RIE to 0.

Bit 6 RIE	Description	
0	Receive data full interrupt request (RXI) and receive error interrupt request (ERI) disabled	(initial value)
1	Receive data full interrupt request (RXI) and receive error interrupt request (ERI) enabled	

Bit 5: Transmit enable (TE)

Bit 5 selects enabling or disabling of the start of transmit operation.

Bit 5 TE	Description	
0	Transmit operation disabled ^{*1} (TXD ₃₂ pin is transmit data pin)	(initial value)
1	Transmit operation enabled ^{*2} (TXD ₃₂ pin is transmit data pin)	

Notes: 1. Bit TDRE in SSR is fixed at 1.

 When transmit data is written to TDR in this state, bit TDR in SSR is cleared to 0 and serial data transmission is started. Be sure to carry out serial mode register (SMR) settings, and setting of bit SPC32 in SPCR, to decide the transmission format before setting bit TE to 1.

Bit 4: Receive enable (RE)

Bit 4 selects enabling or disabling of the start of receive operation.

Bit 4 RE	Description	
0	Receive operation disabled ^{*1} (RXD pin is I/O port)	(initial value)
1	Receive operation enabled ^{*2} (RXD pin is receive data pin)	
Notes: 1	. Note that the RDRF, FER, PER, and OER flags in SSR are not affected cleared to 0, and retain their previous state.	when bit RE is

 In this state, serial data reception is started when a start bit is detected in asynchronous mode or serial clock input is detected in synchronous mode. Be sure to carry out serial mode register (SMR) settings to decide the reception format before setting bit RE to 1.

Bit 3: Multiprocessor interrupt enable (MPIE)

Bit 3 selects enabling or disabling of the multiprocessor interrupt request. The MPIE bit setting is only valid when asynchronous mode is selected and reception is carried out with bit MP in SMR set to 1. The MPIE bit setting is invalid when bit COM is set to 1 or bit MP is cleared to 0.

Bit 3 MPIE	Description	
0	Multiprocessor interrupt request disabled (normal receive operation) Clearing conditions: When data is received in which the multiprocessor bit is set to 1	(initial value)
1	Multiprocessor interrupt request enabled*	
Note: * R	eceive data transfer from RSR to RDR, receive error detection, and setting	of the RDRF.

ote: * Receive data transfer from RSR to RDR, receive error detection, and setting of the RDRF, FER, and OER status flags in SSR is not performed. RXI, ERI, and setting of the RDRF, FER, and OER flags in SSR, are disabled until data with the multiprocessor bit set to 1 is received. When a receive character with the multiprocessor bit set to 1 is received, bit MPBR in SSR is set to 1, bit MPIE is automatically cleared to 0, and RXI and ERI requests (when bits TIE and RIE in serial control register 3 (SCR3) are set to 1) and setting of the RDRF, FER, and OER flags are enabled.

Bit 2: Transmit end interrupt enable (TEIE)

Bit 2 selects enabling or disabling of the transmit end interrupt request (TEI) if there is no valid transmit data in TDR when MSB data is to be sent.

Bit 2 TEIE	Description	
0	Transmit end interrupt request (TEI) disabled	(initial value)
1	Transmit end interrupt request (TEI) enabled*	

Note: * TEI can be released by clearing bit TDRE to 0 and clearing bit TEND to 0 in SSR, or by clearing bit TEIE to 0.

Bits 1 and 0: Clock enable 1 and 0 (CKE1, CKE0)

Bits 1 and 0 select the clock source and enabling or disabling of clock output from the SCK_{32} pin. The combination of CKE1 and CKE0 determines whether the SCK_{32} pin functions as an I/O port, a clock output pin, or a clock input pin.

The CKE0 bit setting is only valid in case of internal clock operation (CKE1 = 0) in asynchronous mode. In synchronous mode, or when external clock operation is used (CKE1 = 1), bit CKE0 should be cleared to 0.

After setting bits CKE1 and CKE0, set the operating mode in the serial mode register (SMR).

Bit 1	Bit 0		Description	
CKE1	CKE0	Communication Mode	Clock Source	SCK ₃₂ Pin Function
0	0	Asynchronous	Internal clock	I/O port ^{*1}
		Synchronous	Internal clock	Serial clock output*1
0	1	Asynchronous	Internal clock	Clock output*2
		Synchronous	Reserved	
1	0	Asynchronous	External clock	Clock input ^{*3}
		Synchronous	External clock	Serial clock input
1	1	Asynchronous	Reserved	
		Synchronous	Reserved	

For details on clock source selection, see table 10.9 in 10.3.1.

Notes: 1. Initial value

2. A clock with the same frequency as the bit rate is output.

3. Input a clock with a frequency 16 times the bit rate.

10.2.7 Serial status register (SSR)

Bit	7	6	5	4	3	2	1	0
	TDRE	RDRF	OER	FER	PER	TEND	MPBR	MPBT
Initial value	1	0	0	0	0	1	0	0
Read/Write	R/(W)*	R/(W)*	R/(W)*	R/(W)*	R/(W)*	R	R	R/W
Note: * Only a write of 0 for flag clearing is passible								

Note: * Only a write of 0 for flag clearing is possible.

SSR is an 8-bit register containing status flags that indicate the operational status of SCI3, and multiprocessor bits.

SSR can be read or written to by the CPU at any time, but 1 cannot be written to bits TDRE, RDRF, OER, PER, and FER.

Bits TEND and MPBR are read-only bits, and cannot be modified.

SSR is initialized to H'84 upon reset, and in standby, module standby, or watch mode.

Bit 7: Transmit data register empty (TDRE)

Bit 7 indicates that transmit data has been transferred from TDR to TSR.

Bit 7 TDRE	Description	
0	Transmit data written in TDR has not been transferred to TSR Clearing conditions: After reading TDRE = 1, cleared by writing 0 to TDRE When data is written to TDR by an instruction	
1	Transmit data has not been written to TDR, or transmit data written in TDR has been transferred to TSR Setting conditions: When bit TE in SCR3 is cleared to 0 When data is transferred from TDR to TSR	(initial value)

Bit 6: Receive data register full (RDRF)

Bit 6 indicates that received data is stored in RDR.

Bit 6 RDRF	Description
0	There is no receive data in RDR(initial value)Clearing conditions:After reading RDRF = 1, cleared by writing 0 to RDRFWhen RDR data is read by an instruction
1	There is receive data in RDR Setting conditions: When reception ends normally and receive data is transferred from RSR to RDR
Note:	If an error is detected in the receive data, or if the RE bit in SCR3 has been cleared to 0, RDR and bit RDRF are not affected and retain their previous state. Note that if data reception is completed while bit RDRF is still set to 1, an overrun error (OER) will result and the receive data will be lost.

Bit 5: Overrun error (OER)

Bit 5 indicates that an overrun error has occurred during reception.

Bit 5 OER	Description	
0	Reception in progress or completed ^{*1} Clearing conditions: After reading OER = 1, cleared by writing 0 to OER	(initial value)
1	An overrun error has occurred during reception ^{*2} Setting conditions: When reception is completed with RDRF set to 1	
Notes: 7	 When bit RE in SCR3 is cleared to 0, bit OER is not affected and retain state. 	ns its previous

^{2.} RDR retains the receive data it held before the overrun error occurred, and data received after the error is lost. Reception cannot be continued with bit OER set to 1, and in synchronous mode, transmission cannot be continued either.

Bit 4: Framing error (FER)

Bit 4 indicates that a framing error has occurred during reception in asynchronous mode.

Bit 4 FER	Description	
0	Reception in progress or completed ^{*1} Clearing conditions:	(initial value)
	After reading FER = 1, cleared by writing 0 to FER	
1	A framing error has occurred during reception Setting conditions:	
	When the stop bit at the end of the receive data is checked for a valor of 1 at the end of reception, and the stop bit is 0 ^{*2}	alue
Notes:	When bit RE in SCR3 is cleared to 0, bit FER is not affected and reta state.	ins its previous
	Note that, in 2-stop-bit mode, only the first stop bit is checked for a vase second stop bit is not checked. When a framing error occurs the rece transferred to RDR but bit RDRF is not set. Reception cannot be con FER set to 1. In synchronous mode, neither transmission nor reception when bit FER is set to 1.	eive data is itinued with bit

Bit 3: Parity error (PER)

Bit 3 indicates that a parity error has occurred during reception with parity added in asynchronous mode.

Bit 3 PER	Description
0	Reception in progress or completed*1(initial value)Clearing conditions:After reading PER = 1, cleared by writing 0 to PER
1	A parity error has occurred during reception ^{*2} Setting conditions: When the number of 1 bits in the receive data plus parity bit does not match the parity designated by bit PM in the serial mode register (SMR)
Notes:	 When bit RE in SCR3 is cleared to 0, bit PER is not affected and retains its previous state.
:	 Receive data in which it a parity error has occurred is still transferred to RDR, but bit RDRF is not set. Reception cannot be continued with bit PER set to 1. In synchronous mode, neither transmission nor reception is possible when bit FER is set to 1.

Bit 2: Transmit end (TEND)

Bit 2 indicates that bit TDRE is set to 1 when the last bit of a transmit character is sent.

Bit 2 is a read-only bit and cannot be modified.

Bit 2 TEND	Description
0	Transmission in progress Clearing conditions: After reading TDRE = 1, cleared by writing 0 to TDRE When data is written to TDR by an instruction
1	Transmission ended(initial value)Setting conditions:When bit TE in SCR3 is cleared to 0When bit TDRE is set to 1 when the last bit of a transmit character is sent

Bit 1: Multiprocessor bit receive (MPBR)

Bit 1 stores the multiprocessor bit in a receive character during multiprocessor format reception in asynchronous mode.

Bit 1 is a read-only bit and cannot be modified.

Bit 1 MPBR	Description	
0	Data in which the multiprocessor bit is 0 has been received*	(initial value)
1	Data in which the multiprocessor bit is 1 has been received	

Note: * When bit RE is cleared to 0 in SCR3 with the multiprocessor format, bit MPBR is not affected and retains its previous state.

Bit 0: Multiprocessor bit transfer (MPBT)

Bit 0 stores the multiprocessor bit added to transmit data when transmitting in asynchronous mode. The bit MPBT setting is invalid when synchronous mode is selected, when the multiprocessor communication function is disabled, and when not transmitting.

Bit 0 MPBT	Description	
0	A 0 multiprocessor bit is transmitted	(initial value)
1	A 1 multiprocessor bit is transmitted	

10.2.8 Bit rate register (BRR)

Bit	7	6	5	4	3	2	1	0
	BRR7	BRR6	BRR5	BRR4	BRR3	BRR2	BRR1	BRR0
Initial value	1	1	1	1	1	1	1	1
Read/Write	R/W							

BRR is an 8-bit register that designates the transmit/receive bit rate in accordance with the baud rate generator operating clock selected by bits CKS1 and CKS0 of the serial mode register (SMR).

BRR can be read or written by the CPU at any time.

BRR is initialized to H'FF upon reset, and in standby, module standby, or watch mode.

Table 10.3 shows examples of BRR settings in asynchronous mode. The values shown are for active (high-speed) mode.

	OSC															
	32.8 kHz			38.4 kHz				2 MHz		2.	2.4576 MHz				4 MHz	
B Bit Rate		Error		Error				Erro		Error	r		Error			
(bit/s)	n	Ν	(%)	n	Ν	(%)	n	Ν	(%)	n	Ν	(%)	n	Ν	(%)	
110	Ca	nnot k	be used,	_	_				_	2	21	-0.83	_			
150	as	error		0	3	0	2	12	0.16	3	3	0	2	25	0.16	
200	exc	ceeds	3%	0	2	0	0	155	0.16	3	2	0	_		_	
250				—	_	_	0	124	0	0	153	-0.26	0	249	0	
300	_			0	1	0	0	103	0.16	3	1	0	2	12	0.16	
600				0	0	0	0	51	0.16	3	0	0	0	103	0.16	
1200	_			_	_	_	0	25	0.16	2	1	0	0	51	0.16	
2400	-			_	_		0	12	0.16	2	0	0	0	25	0.16	
4800	-			_				_	_	0	7	0	0	12	0.16	
9600	_			_	_		—	_		0	3	0	—	_	_	
19200	-			_	_				_	0	1	0			_	
31250	_			_			0	0	0				0	1	0	
38400	_			_						0	0	0			_	

 Table 10.3
 Examples of BRR Settings for Various Bit Rates (Asynchronous Mode) (1)

	OSC					
		10 M	Hz		16 MI	Hz
B Bit Rate (bit/s)	n	N	Error (%)	n	N	Error (%)
110	2	88	-0.25	2	141	-0.02
150	2	64	0.16	2	103	0.16
200	2	48	-0.35	2	77	0.16
250	2	38	0.16	2	62	-0.79
300	_	_	_	2	51	0.16
600			—	2	25	0.16
1200	0	129	0.16	0	207	0.16
2400	0	64	0.16	0	103	0.16
4800			_	0	51	0.16
9600	_	_	_	0	25	0.16
19200	_	_	_	0	12	0.16
31250	0	4	0	0	7	0
38400	_	_		_		

 Table 10.3
 Examples of BRR Settings for Various Bit Rates (Asynchronous Mode) (2)

Notes: 1. The setting should be made so that the error is not more than 1%.

2. The value set in BRR is given by the following equation:

$$N = \frac{OSC}{(64 \times 2^{2n} \times B)} - 1$$

where

- B: Bit rate (bit/s)
- N: Baud rate generator BRR setting ($0 \le N \le 255$)

OSC: Value of ØOSC (Hz)

- n: Baud rate generator input clock number (n = 0, 2, or 3)(The relation between n and the clock is shown in table 10.4.)
- 3. The error in table 10.3 is the value obtained from the following equation, rounded to two decimal places.

$$Error (\%) = \frac{B (rate obtained from n, N, OSC) - R(bit rate in left-hand column in table 10.3.)}{R (bit rate in left-hand column in table 10.3.)} \times 100$$

		SMR Setting		
n	Clock	CKS1	CKS0	
0	Ø	0	0	
0	ø _w /2 ^{*1} /ø _w ^{*2}	0	1	
2	ø/16	1	0	
3	ø/64	1	1	

Table 10.4 Relation between n and Clock

Notes: 1. ø w/2 clock in active (medium-speed/high-speed) mode and sleep mode

2. ø w clock in subactive mode and subsleep mode

In subactive or subsleep mode, SCI3 can be operated when CPU clock is øw/2 only.

Table 10.5 shows the maximum bit rate for each frequency. The values shown are for active (high-speed) mode.

	Maximum Bit Rate	Setting		
OSC (MHz)	(bit/s)	n	Ν	
0.0384*	600	0	0	
2	31250	0	0	
2.4576	38400	0	0	
4	62500	0	0	
10	156250	0	0	
16	250000	0	0	

Note: * When SMR is set up to CKS1 = 0, CKS0 = 1.

Table 10.6 shows examples of BRR settings in synchronous mode. The values shown are for active (high-speed) mode.

					OSC				
B Bit Rate		38.4 k	Hz		2 MH	z		4 MH	z
(bit/s)	n	Ν	Error	n	Ν	Error	n	Ν	Error
200	0	23	0	_	_	_	_	_	_
250		—	_	—	_	_	2	124	0
300	2	0	0				_		_
500									
1k				0	249	0	—		_
2.5k				0	99	0	0	199	0
5k				0	49	0	0	99	0
10k				0	24	0	0	49	0
25k				0	9	0	0	19	0
50k				0	4	0	0	9	0
100k				_	_	_	0	4	0
250k			0	0	0	0	0	1	0
500k							0	0	0
1M									

 Table 10.6
 Examples of BRR Settings for Various Bit Rates (Synchronous Mode) (1)

	osc					
B Bit Rate		10 MH	lz		16 MF	lz
(bit/s)	n	Ν	Error	n	Ν	Error
200		_			_	_
250		—		3	124	0
300				_	_	
500				2	249	0
1k		—		2	124	0
2.5k				2	49	0
5k	0	249	0	2	24	0
10k	0	124	0	0	199	0
25k	0	49	0	0	79	0
50k	0	24	0	0	39	0
100k		_		0	19	0
250k	0	4	0	0	7	0
500k				0	3	0
1M				0	1	0

 Table 10.6
 Examples of BRR Settings for Various Bit Rates (Synchronous Mode) (2)

Blank: Cannot be set.

- : A setting can be made, but an error will result.

Notes: The value set in BRR is given by the following equation:

$$N = \frac{OSC}{(8 \times 2^{2n} \times B)} - 1$$

where

B: Bit rate (bit/s)

N: Baud rate generator BRR setting ($0 \le N \le 255$)

OSC: Value of ø_{OSC} (Hz)

n: Baud rate generator input clock number (n = 0, 2, or 3)(The relation between n and the clock is shown in table 10.7.)

		SMR Setting		
n	Clock	CKS1	CKS0	
0	Ø	0	0	
0	ø _w /2 ^{*1} /ø _w *2	0	1	
2	ø/16	1	0	
3	ø/64	1	1	

Table 10.7 Relation between n and Clock

Notes: 1. ø w/2 clock in active (medium-speed/high-speed) mode and sleep mode

2. ø w clock in subactive mode and subsleep mode

In subactive or subsleep mode, SCI3 can be operated when CPU clock is øw/2 only.

10.2.9 Clock stop register 1 (CKSTPR1)

Bit	7	6	5	4	3	2	1	0
		_	S32CKSTP	ADCKSTP	—	TFCKSTP	—	TACKSTP
Initial value	1	1	1	1	1	1	1	1
Read/Write	_	_	R/W	R/W	_	R/W	_	R/W

CKSTPR1 is an 8-bit read/write register that performs module standby mode control for peripheral modules. Only the bits relating to SCI3 are described here. For details of the other bits, see the sections on the relevant modules.

Bit 5: SCI3 module standby mode control (S32CKSTP)

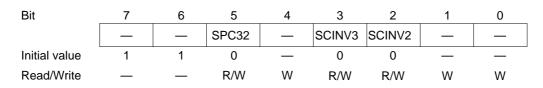
Bit 5 controls setting and clearing of module standby mode for SCI3.

S32CKSTP Description

0	SCI3 is set to module standby mode	
1	SCI3 module standby mode is cleared	(initial value)
NI 4		

Note: All SCI3 register is initialized in module standby mode.

10.2.10 Serial Port Control Register (SPCR)



SPCR is an 8-bit readable/writable register that performs RXD₃₂ and TXD₃₂ pin input/output data inversion switching.

Bits 7 and 6: Reserved bits

Bits 7 and 6 are reserved; they are always read as 1 and cannot be modified.

Bit 5: P4₂/TXD₃₂ pin function switch (SPC32)

This bit selects whether pin $P4_2/TXD_{32}$ is used as $P4_2$ or as TXD_{32} .

Bit 5 SPC32	Description	
0	Functions as P4 ₂ I/O pin	(initial value)
1	Functions as TXD ₃₂ output pin*	

Note: * Set the TE bit in SCR3 after setting this bit to 1.

Bits 4, 1, and 0: Reserved bits

Bits 4, 1, and 0 are reserved; only 0 can be written to these bits.

Bit 3 specifies whether or not TXD_{32} pin output data is to be inverted.

Bit 3 SCINV3	Description	
0	TXD ₃₂ output data is not inverted	(initial value)
1	TXD ₃₂ output data is inverted	

Bit 2: RXD₃₂ pin input data inversion switch

Bit 2 specifies whether or not RXD_{32} pin input data is to be inverted.

Bit 2 SCINV2	Description	
0	RXD ₃₂ input data is not inverted	(initial value)
1	RXD ₃₂ input data is inverted	

10.3 Operation

10.3.1 Overview

SCI3 can perform serial communication in two modes: asynchronous mode in which synchronization is provided character by character, and synchronous mode in which synchronization is provided by clock pulses. The serial mode register (SMR) is used to select asynchronous or synchronous mode and the data transfer format, as shown in table 10.8.

The clock source for SCI3 is determined by bit COM in SMR and bits CKE1 and CKE0 in SCR3, as shown in table 10.9.

1. Asynchronous mode

- Choice of 5-, 7-, or 8-bit data length
- Choice of parity addition, multiprocessor bit addition, and addition of 1 or 2 stop bits. (The combination of these parameters determines the data transfer format and the character length.)
- Framing error (FER), parity error (PER), overrun error (OER), and break detection during reception
- Choice of internal or external clock as the clock source

When internal clock is selected: SCI3 operates on the baud rate generator clock, and a clock with the same frequency as the bit rate can be output.

When external clock is selected: A clock with a frequency 16 times the bit rate must be input. (The on-chip baud rate generator is not used.)

- 2. Synchronous mode
- Data transfer format: Fixed 8-bit data length
- Overrun error (OER) detection during reception
- Choice of internal or external clock as the clock source When internal clock is selected: SCI3 operates on the baud rate generator clock, and a serial clock is output.

When external clock is selected: The on-chip baud rate generator is not used, and SCI3 operates on the input serial clock.

SMR						Data Transfer Format			
bit 7 COM	bit 6 CHR	bit 2 MP	bit 5 PE	bit 3 STOP	Mode	Data Length	Multiprocessor Bit	Parity Bit	Stop Bit Length
0	0	0	0	0	Asynchronous	8-bit data	No	No	1 bit
				1	mode				2 bits
			1	0	-			Yes	1 bit
				1	-				2 bits
	1	_	0	0	-	7-bit data	-	No	1 bit
				1	-				2 bits
			1	0	-			Yes	1 bit
				1	-				2 bits
	0	1	0	0	-	8-bit data	Yes	No	1 bit
				1	-				2 bits
			1	0	-	5-bit data	No		1 bit
				1	-				2 bits
	1	-	0	0	-	7-bit data	Yes		1 bit
				1	-				2 bits
			1	0	-	5-bit data	No	Yes	1 bit
				1	-				2 bits
1	*	0	*	*	Synchronous mode	8-bit data	No	No	No

Table 10.8 SMR Settings and Corresponding Data Transfer Formats

*: Don't care

SMR	SCR3				
bit 7	bit 1 bit 0				Transmit/Receive Clock
COM	CKE1	CKE0	Mode	Clock Source	SCK ₃₂ Pin Function
0	0	0	Asynchronous	Internal	I/O port (SCK ₃₂ pin not used)
		1	mode		Outputs clock with same frequency as bit rate
	1	0	-	External	Inputs clock with frequency 16 times bit rate
1	0	0	Synchronous	Internal	Outputs serial clock
	1	0	mode	External	Inputs serial clock
0	1	1	Reserved (Do r	not specify these	combinations)
1	0	1			
1	1	1			

 Table 10.9
 SMR and SCR3 Settings and Clock Source Selection

3. Interrupts and continuous transmission/reception

SCI3 can carry out continuous reception using RXI and continuous transmission using TXI. These interrupts are shown in table 10.10.

Interrupt	Flags	Interrupt Request Conditions	Notes
RXI	RDRF RIE	When serial reception is performed normally and receive data is transferred from RSR to RDR, bit RDRF is set to 1, and if bit RIE is set to 1 at this time, RXI is enabled and an interrupt is requested. (See figure 10.2 (a).)	The RXI interrupt routine reads the receive data transferred to RDR and clears bit RDRF to 0. Continuous reception can be performed by repeating the above operations until reception of the next RSR data is completed.
ТХІ	TDRE TIE	When TSR is found to be empty (on completion of the previous transmission) and the transmit data placed in TDR is transferred to TSR, bit TDRE is set to 1. If bit TIE is set to 1 at this time, TXI is enabled and an interrupt is requested. (See figure 10.2 (b).)	The TXI interrupt routine writes the next transmit data to TDR and clears bit TDRE to 0. Continuous transmission can be performed by repeating the above operations until the data transferred to TSR has been transmitted.
TEI	TEND TEIE	When the last bit of the character in TSR is transmitted, if bit TDRE is set to 1, bit TEND is set to 1. If bit TEIE is set to 1 at this time, TEI is enabled and an interrupt is requested. (See figure 10.2 (c).)	TEI indicates that the next transmit data has not been written to TDR when the last bit of the transmit character in TSR is sent.

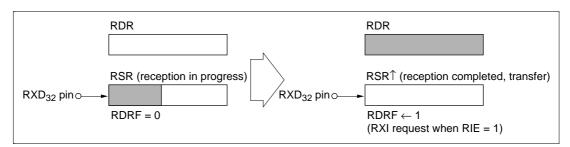


Figure 10.2 (a) RDRF Setting and RXI Interrupt

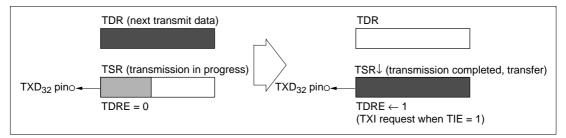


Figure 10.2 (b) TDRE Setting and TXI Interrupt

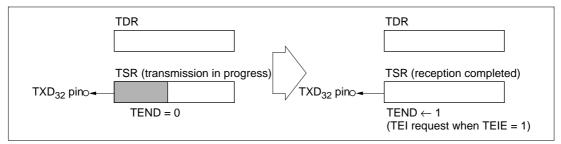


Figure 10.2 (c) TEND Setting and TEI Interrupt

10.3.2 Operation in Asynchronous Mode

In asynchronous mode, serial communication is performed with synchronization provided character by character. A start bit indicating the start of communication and one or two stop bits indicating the end of communication are added to each character before it is sent.

SCI3 has separate transmission and reception units, allowing full-duplex communication. As the transmission and reception units are both double-buffered, data can be written during transmission and read during reception, making possible continuous transmission and reception.

1. Data transfer format

The general data transfer format in asynchronous communication is shown in figure 10.3.

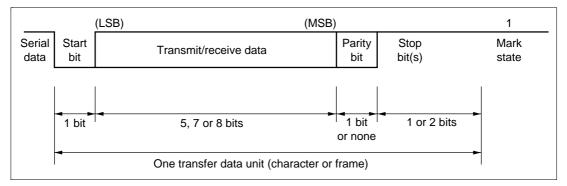


Figure 10.3 Data Format in Asynchronous Communication

In asynchronous communication, the communication line is normally in the mark state (high level). SCI3 monitors the communication line and when it detects a space (low level), identifies this as a start bit and begins serial data communication.

One transfer data character consists of a start bit (low level), followed by transmit/receive data (LSB-first format, starting from the least significant bit), a parity bit (high or low level), and finally one or two stop bits (high level).

In asynchronous mode, synchronization is performed by the falling edge of the start bit during reception. The data is sampled on the 8th pulse of a clock with a frequency 16 times the bit period, so that the transfer data is latched at the center of each bit.

Table 10.11 shows the 16 data transfer formats that can be set in asynchronous mode. The format is selected by the settings in the serial mode register (SMR).

	S	MR		Serial Data Transfer Format and Frame Length
CHR	PE	MP	STOP	1 2 3 4 5 6 7 8 9 10 11 12
0	0	0	0	S 8-bit data STOP
0	0	0	1	S 8-bit data STOP STOP
0	0	1	0	S 8-bit data MPB STOP
0	0	1	1	S 8-bit data MPB STOP STOP
0	1	0	0	S 8-bit data P STOP
0	1	0	1	S 8-bit data P STOP STOP
0	1	1	0	S 5-bit data STOP
0	1	1	1	S 5-bit data STOP STOP
1	0	0	0	S 7-bit data STOP
1	0	0	1	S 7-bit data STOP STOP
1	0	1	0	S 7-bit data MPB STOP
1	0	1	1	S 7-bit data MPB STOP STOP
1	1	0	0	S 7-bit data P STOP
1	1	0	1	S 7-bit data P STOP STOP
1	1	1	0	S 5-bit data P STOP
1	1	1	1	S 5-bit data P STOP STOP

 Table 10.11 Data Transfer Formats (Asynchronous Mode)

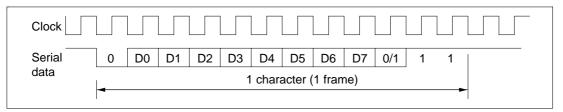
Notation: S: Start bit STOP: Stop bit P: Parity bit MPB: Multiprocessor bit

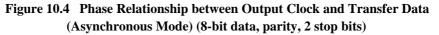
2. Clock

Either an internal clock generated by the baud rate generator or an external clock input at the SCK_{32} pin can be selected as the SCI3 transmit/receive clock. The selection is made by means of bit COM in SMR and bits SCE1 and CKE0 in SCR3. See table 10.9 for details on clock source selection.

When an external clock is input at the SCK_{32} pin, the clock frequency should be 16 times the bit rate.

When SCI3 operates on an internal clock, the clock can be output at the SCK_{32} pin. In this case the frequency of the output clock is the same as the bit rate, and the phase is such that the clock rises at the center of each bit of transmit/receive data, as shown in figure 10.4.





- 3. Data transfer operations
- SCI3 initialization

Before data is transferred on SCI3, bits TE and RE in SCR3 must first be cleared to 0, and then SCI3 must be initialized as follows.

Note: If the operation mode or data transfer format is changed, bits TE and RE must first be cleared to 0.

When bit TE is cleared to 0, bit TDRE is set to 1.

Note that the RDRF, PER, FER, and OER flags and the contents of RDR are retained when RE is cleared to 0.

When an external clock is used in asynchronous mode, the clock should not be stopped during operation, including initialization. When an external clock is used in synchronous mode, the clock should not be supplied during operation, including initialization.

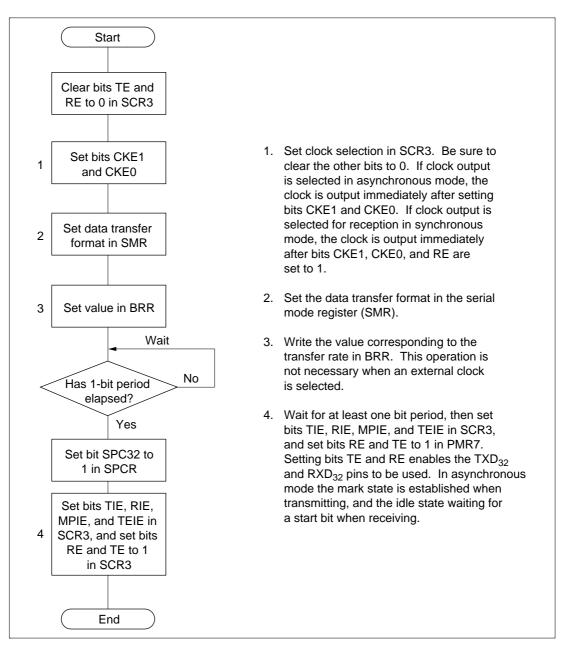


Figure 10.5 shows an example of a flowchart for initializing SCI3.

Figure 10.5 Example of SCI3 Initialization Flowchart

• Transmitting

Figure 10.6 shows an example of a flowchart for data transmission. This procedure should be followed for data transmission after initializing SCI3.

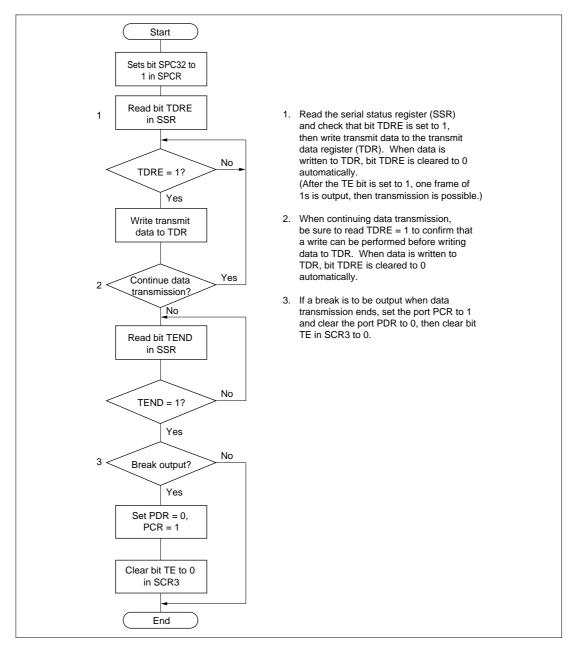


Figure 10.6 Example of Data Transmission Flowchart (Asynchronous Mode)

SCI3 operates as follows when transmitting data.

SCI3 monitors bit TDRE in SSR, and when it is cleared to 0, recognizes that data has been written to TDR and transfers data from TDR to TSR. It then sets bit TDRE to 1 and starts transmitting. If bit TIE in SCR3 is set to 1 at this time, a TXI request is made.

Serial data is transmitted from the TXD_{32} pin using the relevant data transfer format in table 10.11. When the stop bit is sent, SCI3 checks bit TDRE. If bit TDRE is cleared to 0, SCI3 transfers data from TDR to TSR, and when the stop bit has been sent, starts transmission of the next frame. If bit TDRE is set to 1, bit TEND in SSR bit is set to 1the mark state, in which 1s are transmitted, is established after the stop bit has been sent. If bit TEIE in SCR3 is set to 1 at this time, a TEI request is made.

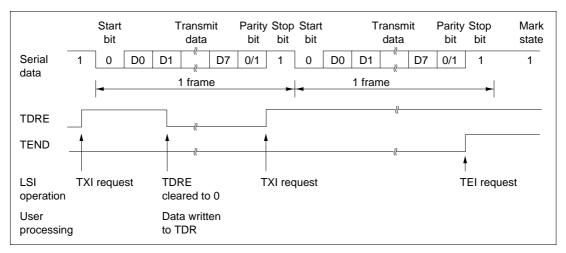
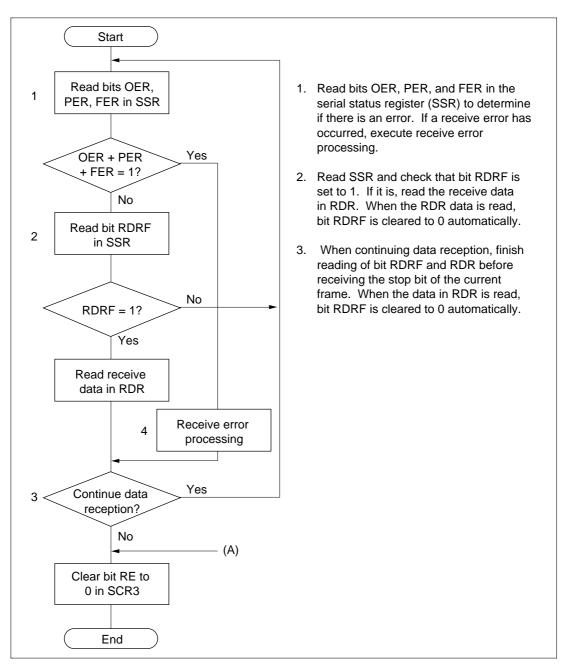


Figure 10.7 shows an example of the operation when transmitting in asynchronous mode.

Figure 10.7 Example of Operation when Transmitting in Asynchronous Mode (8-bit data, parity, 1 stop bit)

• Receiving

Figure 10.8 shows an example of a flowchart for data reception. This procedure should be followed for data reception after initializing SCI3.





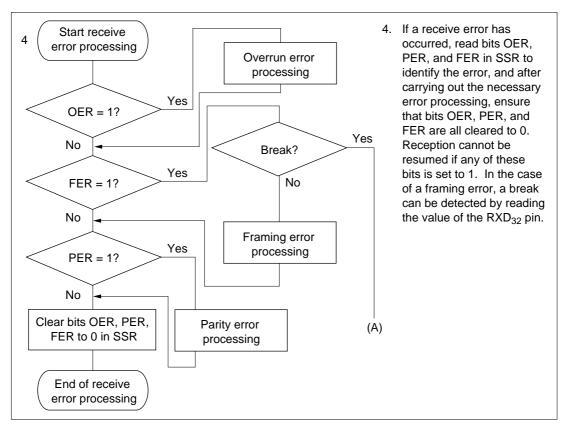


Figure 10.8 Example of Data Reception Flowchart (Asynchronous Mode) (cont)

SCI3 operates as follows when receiving data.

SCI3 monitors the communication line, and when it detects a 0 start bit, performs internal synchronization and begins reception. Reception is carried out in accordance with the relevant data transfer format in table 10.11. The received data is first placed in RSR in LSB-to-MSB order, and then the parity bit and stop bit(s) are received. SCI3 then carries out the following checks.

• Parity check

SCI3 checks that the number of 1 bits in the receive data conforms to the parity (odd or even) set in bit PM in the serial mode register (SMR).

• Stop bit check

SCI3 checks that the stop bit is 1. If two stop bits are used, only the first is checked.

Status check

SCI3 checks that bit RDRF is set to 0, indicating that the receive data can be transferred from RSR to RDR.

If no receive error is found in the above checks, bit RDRF is set to 1, and the receive data is stored in RDR. If bit RIE is set to 1 in SCR3, an RXI interrupt is requested. If the error checks identify a receive error, bit OER, PER, or FER is set to 1 depending on the kind of error. Bit RDRF retains its state prior to receiving the data. If bit RIE is set to 1 in SCR3, an ERI interrupt is requested.

Table 10.12 shows the conditions for detecting a receive error, and receive data processing.

Note: No further receive operations are possible while a receive error flag is set. Bits OER, FER, PER, and RDRF must therefore be cleared to 0 before resuming reception.

Table 10.12 Receive Error Detection Conditions and Receive Data Processing

Receive Error	Abbreviation	Detection Conditions	Receive Data Processing
Overrun error	OER	When the next date receive operation is completed while bit RDRF is still set to 1 in SSR	Receive data is not transferred from RSR to RDR
Framing error	FER	When the stop bit is 0	Receive data is transferred from RSR to RDR
Parity error	PER	When the parity (odd or even) set in SMR is different from that of the received data	Receive data is transferred from RSR to RDR

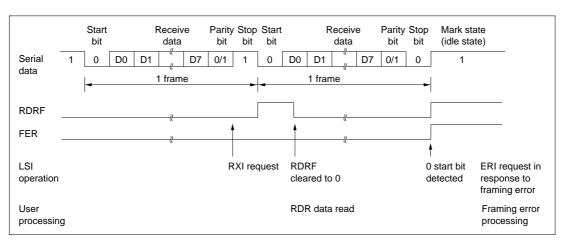
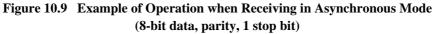


Figure 10.9 shows an example of the operation when receiving in asynchronous mode.



10.3.3 Operation in Synchronous Mode

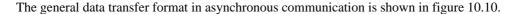
In synchronous mode, SCI3 transmits and receives data in synchronization with clock pulses. This mode is suitable for high-speed serial communication.

SCI3 has separate transmission and reception units, allowing full-duplex communication with a shared clock.

As the transmission and reception units are both double-buffered, data can be written during transmission and read during reception, making possible continuous transmission and reception.

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1. Data transfer format



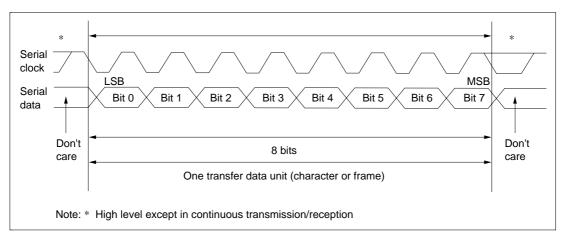


Figure 10.10 Data Format in Synchronous Communication

In synchronous communication, data on the communication line is output from one falling edge of the serial clock until the next falling edge. Data confirmation is guaranteed at the rising edge of the serial clock.

One transfer data character begins with the LSB and ends with the MSB. After output of the MSB, the communication line retains the MSB state.

When receiving in synchronous mode, SCI3 latches receive data at the rising edge of the serial clock.

The data transfer format uses a fixed 8-bit data length.

Parity and multiprocessor bits cannot be added.

2. Clock

Either an internal clock generated by the baud rate generator or an external clock input at the SCK_{32} pin can be selected as the SCI3 serial clock. The selection is made by means of bit COM in SMR and bits CKE1 and CKE0 in SCR3. See table 10.9 for details on clock source selection.

When SCI3 operates on an internal clock, the serial clock is output at the SCK₃₂ pin. Eight pulses of the serial clock are output in transmission or reception of one character, and when SCI3 is not transmitting or receiving, the clock is fixed at the high level.

- 3. Data transfer operations
- SCI3 initialization

Data transfer on SCI3 first of all requires that SCI3 be initialized as described in 10.3.2 3. SCI3 initialization, and shown in figure 10.5.

• Transmitting

Figure 10.11 shows an example of a flowchart for data transmission. This procedure should be followed for data transmission after initializing SCI3.

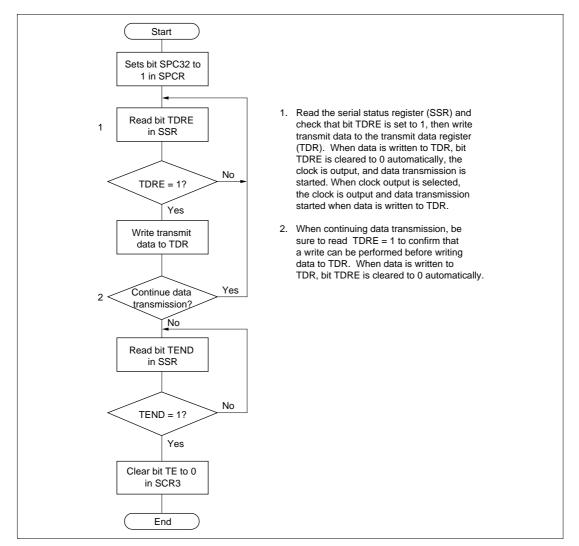


Figure 10.11 Example of Data Transmission Flowchart (Synchronous Mode)

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SCI3 operates as follows when transmitting data.

SCI3 monitors bit TDRE in SSR, and when it is cleared to 0, recognizes that data has been written to TDR and transfers data from TDR to TSR. It then sets bit TDRE to 1 and starts transmitting. If bit TIE in SCR3 is set to 1 at this time, a TXI request is made.

When clock output mode is selected, SCI3 outputs 8 serial clock pulses. When an external clock is selected, data is output in synchronization with the input clock.

Serial data is transmitted from the TXD32 pin in order from the LSB (bit 0) to the MSB (bit 7). When the MSB (bit 7) is sent, checks bit TDRE. If bit TDRE is cleared to 0, SCI3 transfers data from TDR to TSR, and starts transmission of the next frame. If bit TDRE is set to 1, SCI3 sets bit TEND to 1 in SSR, and after sending the MSB (bit 7), retains the MSB state. If bit TEIE in SCR3 is set to 1 at this time, a TEI request is made.

After transmission ends, the SCK pin is fixed at the high level.

Note: Transmission is not possible if an error flag (OER, FER, or PER) that indicates the data reception status is set to 1. Check that these error flags are all cleared to 0 before a transmit operation.

Figure 10.12 shows an example of the operation when transmitting in synchronous mode.

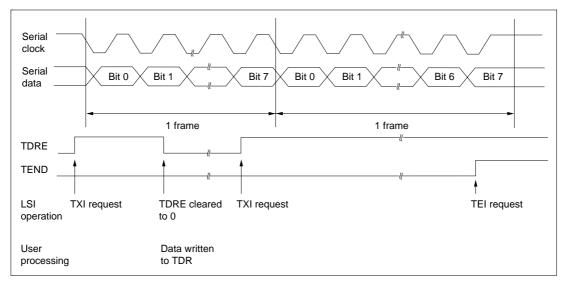


Figure 10.12 Example of Operation when Transmitting in Synchronous Mode

• Receiving

Figure 10.13 shows an example of a flowchart for data reception. This procedure should be followed for data reception after initializing SCI3.

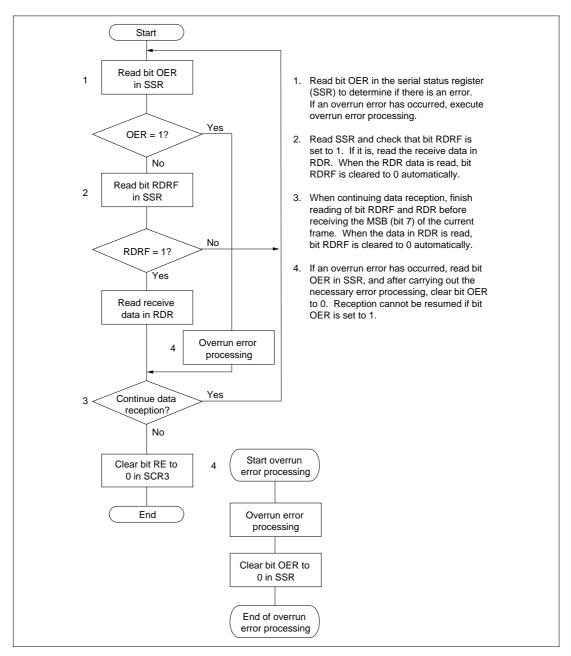


Figure 10.13 Example of Data Reception Flowchart (Synchronous Mode)

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SCI3 operates as follows when receiving data.

SCI3 performs internal synchronization and begins reception in synchronization with the serial clock input or output.

The received data is placed in RSR in LSB-to-MSB order.

After the data has been received, SCI3 checks that bit RDRF is set to 0, indicating that the receive data can be transferred from RSR to RDR.

If this check shows that there is no overrun error, bit RDRF is set to 1, and the receive data is stored in RDR. If bit RIE is set to 1 in SCR3, an RXI interrupt is requested. If the check identifies an overrun error, bit OER is set to 1.

Bit RDRF remains set to 1. If bit RIE is set to 1 in SCR3, an ERI interrupt is requested.

See table 10.12 for the conditions for detecting a receive error, and receive data processing.

Note: No further receive operations are possible while a receive error flag is set. Bits OER, FER, PER, and RDRF must therefore be cleared to 0 before resuming reception.

Figure 10.14 shows an example of the operation when receiving in synchronous mode.

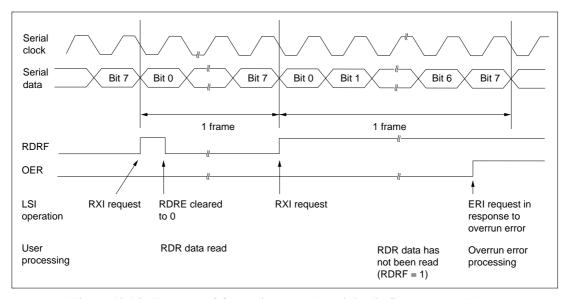


Figure 10.14 Example of Operation when Receiving in Synchronous Mode

• Simultaneous transmit/receive

Figure 10.15 shows an example of a flowchart for a simultaneous transmit/receive operation. This procedure should be followed for simultaneous transmission/reception after initializing SCI3.

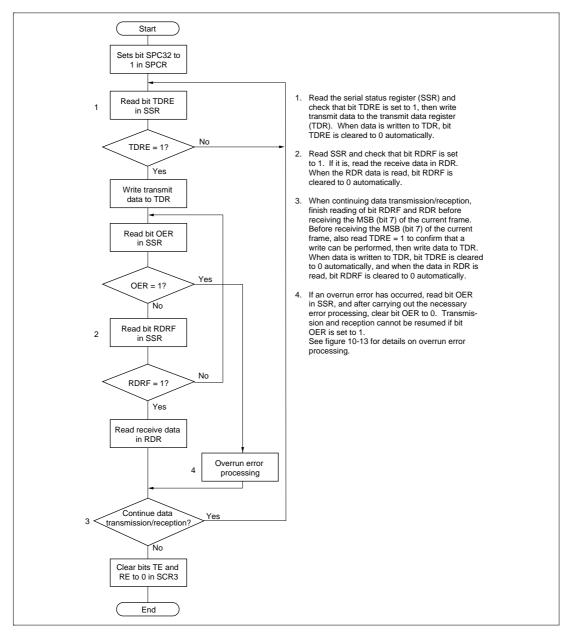


Figure 10.15 Example of Simultaneous Data Transmission/Reception Flowchart (Synchronous Mode)

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- Notes: 1. When switching from transmission to simultaneous transmission/reception, check that SCI3 has finished transmitting and that bits TDRE and TEND are set to 1, clear bit TE to 0, and then set bits TE and RE to 1 simultaneously.
 - 2. When switching from reception to simultaneous transmission/reception, check that SCI3 has finished receiving, clear bit RE to 0, then check that bit RDRF and the error flags (OER, FER, and PER) are cleared to 0, and finally set bits TE and RE to 1 simultaneously.

10.3.4 Multiprocessor Communication Function

The multiprocessor communication function enables data to be exchanged among a number of processors on a shared communication line. Serial data communication is performed in asynchronous mode using the multiprocessor format (in which a multiprocessor bit is added to the transfer data).

In multiprocessor communication, each receiver is assigned its own ID code. The serial communication cycle consists of two cycles, an ID transmission cycle in which the receiver is specified, and a data transmission cycle in which the transfer data is sent to the specified receiver. These two cycles are differentiated by means of the multiprocessor bit, 1 indicating an ID transmission cycle, and 0, a data transmission cycle.

The sender first sends transfer data with a 1 multiprocessor bit added to the ID code of the receiver it wants to communicate with, and then sends transfer data with a 0 multiprocessor bit added to the transmit data. When a receiver receives transfer data with the multiprocessor bit set to 1, it compares the ID code with its own ID code, and if they are the same, receives the transfer data sent next. If the ID codes do not match, it skips the transfer data until data with the multiprocessor bit set to 1 is sent again.

In this way, a number of processors can exchange data among themselves.

Figure 10.16 shows an example of communication between processors using the multiprocessor format.

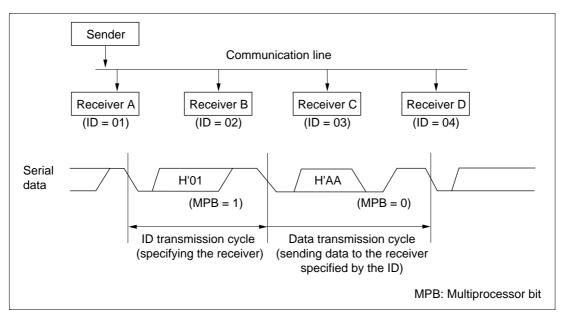


Figure 10.16 Example of Inter-Processor Communication Using Multiprocessor Format (Sending data H'AA to receiver A)

There is a choice of four data transfer formats. If a multiprocessor format is specified, the parity bit specification is invalid. See table 10.11 for details.

For details on the clock used in multiprocessor communication, see 10.3.2, Operation in Synchronous Mode.

• Multiprocessor transmitting

Figure 10.17 shows an example of a flowchart for multiprocessor data transmission. This procedure should be followed for multiprocessor data transmission after initializing SCI3.

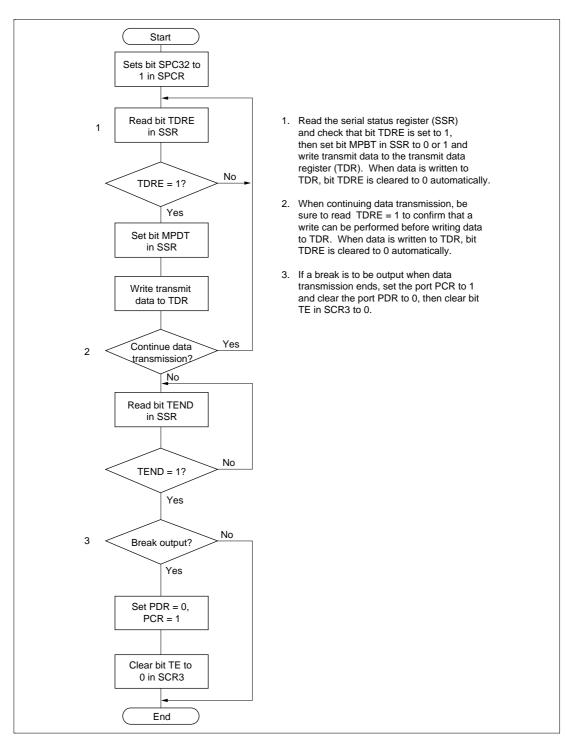


Figure 10.17 Example of Multiprocessor Data Transmission Flowchart

SCI3 operates as follows when transmitting data.

SCI3 monitors bit TDRE in SSR, and when it is cleared to 0, recognizes that data has been written to TDR and transfers data from TDR to TSR. It then sets bit TDRE to 1 and starts transmitting. If bit TIE in SCR3 is set to 1 at this time, a TXI request is made.

Serial data is transmitted from the TXD pin using the relevant data transfer format in table 10.11. When the stop bit is sent, SCI3 checks bit TDRE. If bit TDRE is cleared to 0, SCI3 transfers data from TDR to TSR, and when the stop bit has been sent, starts transmission of the next frame. If bit TDRE is set to 1 bit TEND in SSR bit is set to 1, the mark state, in which 1s are transmitted, is established after the stop bit has been sent. If bit TEIE in SCR3 is set to 1 at this time, a TEI request is made.

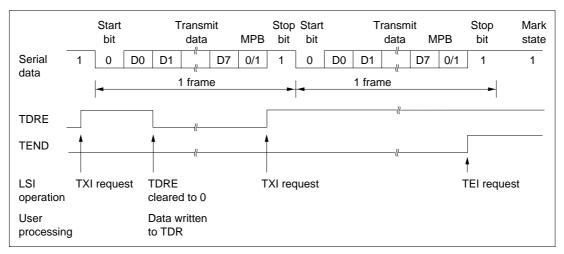
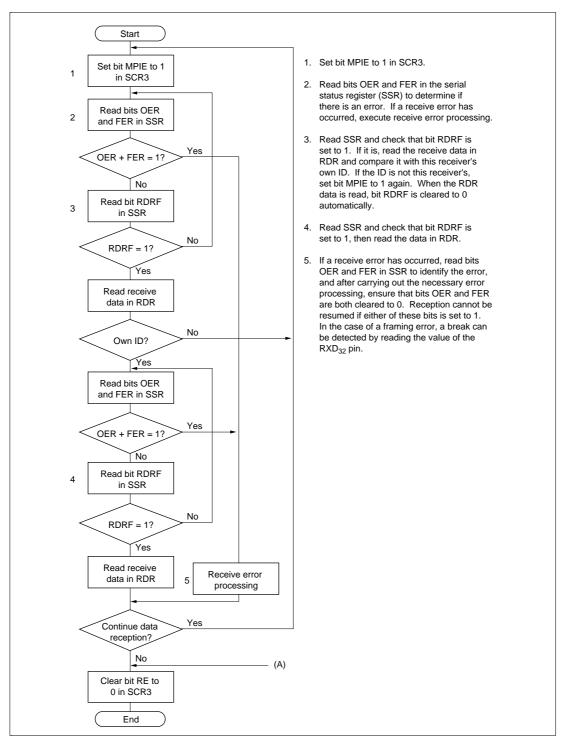


Figure 10.18 shows an example of the operation when transmitting using the multiprocessor format.

Figure 10.18 Example of Operation when Transmitting using Multiprocessor Format (8-bit data, multiprocessor bit, 1 stop bit)

· Multiprocessor receiving

Figure 10.19 shows an example of a flowchart for multiprocessor data reception. This procedure should be followed for multiprocessor data reception after initializing SCI3.





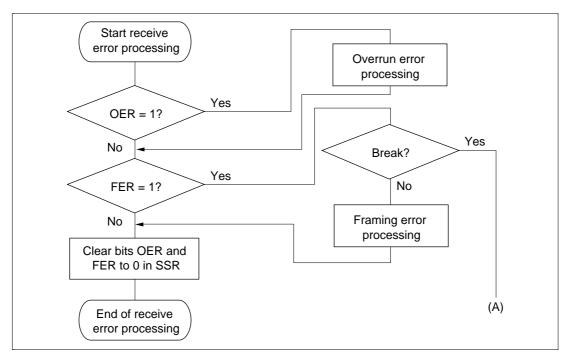


Figure 10.19 Example of Multiprocessor Data Reception Flowchart (cont)

Figure 10.20 shows an example of the operation when receiving using the multiprocessor format.

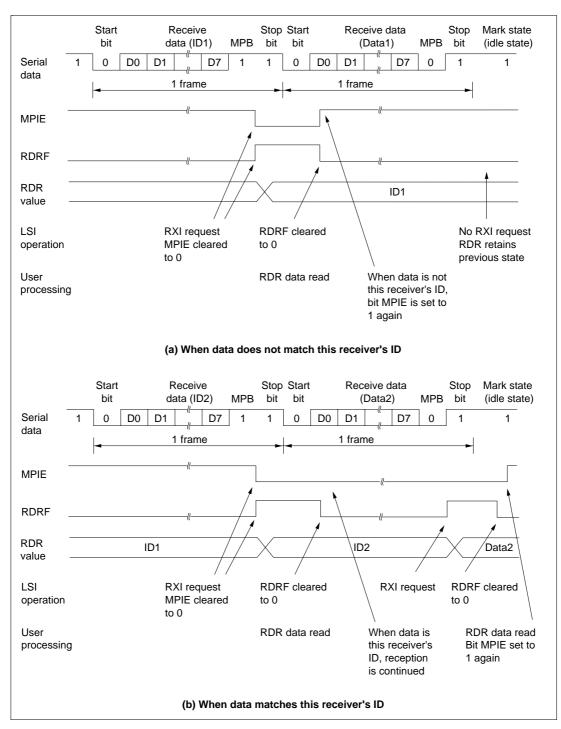


Figure 10.20 Example of Operation when Receiving using Multiprocessor Format (8-bit data, multiprocessor bit, 1 stop bit)

10.4 Interrupts

SCI3 can generate six kinds of interrupts: transmit end, transmit data empty, receive data full, and three receive error interrupts (overrun error, framing error, and parity error). These interrupts have the same vector address.

The various interrupt requests are shown in table 10.13.

Interrupt Abbreviation	Interrupt Request	Vector Address
RXI	Interrupt request initiated by receive data full flag (RDRF)	H'0024
ТХІ	Interrupt request initiated by transmit data empty flag (TDRE)	-
TEI	Interrupt request initiated by transmit end flag (TEND)	-
ERI	Interrupt request initiated by receive error flag (OER, FER, PER)	-

Table 10.13 SCI3 Interrupt Requests

Each interrupt request can be enabled or disabled by means of bits TIE and RIE in SCR3.

When bit TDRE is set to 1 in SSR, a TXI interrupt is requested. When bit TEND is set to 1 in SSR, a TEI interrupt is requested. These two interrupts are generated during transmission.

The initial value of bit TDRE in SSR is 1. Therefore, if the transmit data empty interrupt request (TXI) is enabled by setting bit TIE to 1 in SCR3 before transmit data is transferred to TDR, a TXI interrupt will be requested even if the transmit data is not ready.

Also, the initial value of bit TEND in SSR is 1. Therefore, if the transmit end interrupt request (TEI) is enabled by setting bit TEIE to 1 in SCR3 before transmit data is transferred to TDR, a TEI interrupt will be requested even if the transmit data has not been sent.

Effective use of these interrupt requests can be made by having processing that transfers transmit data to TDR carried out in the interrupt service routine.

To prevent the generation of these interrupt requests (TXI and TEI), on the other hand, the enable bits for these interrupt requests (bits TIE and TEIE) should be set to 1 after transmit data has been transferred to TDR.

When bit RDRF is set to 1 in SSR, an RXI interrupt is requested, and if any of bits OER, PER, and FER is set to 1, an ERI interrupt is requested. These two interrupt requests are generated during reception.

For further details, see 3.3, Interrupts.

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10.5 **Application Notes**

The following points should be noted when using SCI3.

1. Relation between writes to TDR and bit TDRE

Bit TDRE in the serial status register (SSR) is a status flag that indicates that data for serial transmission has not been prepared in TDR. When data is written to TDR, bit TDRE is cleared to 0 automatically. When SCI3 transfers data from TDR to TSR, bit TDRE is set to 1.

Data can be written to TDR irrespective of the state of bit TDRE, but if new data is written to TDR while bit TDRE is cleared to 0, the data previously stored in TDR will be lost of it has not yet been transferred to TSR. Accordingly, to ensure that serial transmission is performed dependably, you should first check that bit TDRE is set to 1, then write the transmit data to TDR once only (not two or more times).

2. Operation when a number of receive errors occur simultaneously

If a number of receive errors are detected simultaneously, the status flags in SSR will be set to the states shown in table 10.14. If an overrun error is detected, data transfer from RSR to RDR will not be performed, and the receive data will be lost.

SSR Status Flags				Receive Data Transfer	
RDRF*	OER	FER	PER	RSR o RDR	Receive Error Status
1	1	0	0	Х	Overrun error
0	0	1	0	0	Framing error
0	0	0	1	0	Parity error
1	1	1	0	Х	Overrun error + framing error
1	1	0	1	Х	Overrun error + parity error
0	0	1	1	0	Framing error + parity error
1	1	1	1	Х	Overrun error + framing error + parity error
0 0					

Table 10.14 SSR Status Flag States and Receive Data Transfer

O: Receive data is transferred from RSR to RDR.

X : Receive data is not transferred from RSR to RDR.

Note: * Bit RDRF retains its state prior to data reception. However, note that if RDR is read after an overrun error has occurred in a frame because reading of the receive data in the previous frame was delayed, RDRF will be cleared to 0.

3. Break detection and processing

When a framing error is detected, a break can be detected by reading the value of the RXD_{32} pin directly. In a break, the input from the RXD_{32} pin becomes all 0s, with the result that bit FER is set and bit PER may also be set.

SCI3 continues the receive operation even after receiving a break. Note, therefore, that even though bit FER is cleared to 0 it will be set to 1 again.

4. Mark state and break detection

When bit TE is cleared to 0, the TXD_{32} pin functions as an I/O port whose input/output direction and level are determined by PDR and PCR. This fact can be used to set the TXD_{32} pin to the mark state, or to detect a break during transmission.

To keep the communication line in the mark state (1 state) until bit TE is set to 1, set PCR = 1 and PDR = 1. Since bit TE is cleared to 0 at this time, the TXD_{32} pin functions as an I/O port and 1 is output.

To detect a break, clear bit TE to 0 after setting PCR = 1 and PDR = 0.

When bit TE is cleared to 0, the transmission unit is initialized regardless of the current transmission state, the TXD_{32} pin functions as an I/O port, and 0 is output from the TXD_{32} pin.

5. Receive error flags and transmit operation (synchronous mode only)

When a receive error flag (OER, PER, or FER) is set to 1, transmission cannot be started even if bit TDRE is cleared to 0. The receive error flags must be cleared to 0 before starting transmission.

Note also that receive error flags cannot be cleared to 0 even if bit RE is cleared to 0.

6. Receive data sampling timing and receive margin in asynchronous mode

In asynchronous mode, SCI3 operates on a basic clock with a frequency 16 times the transfer rate. When receiving, SCI3 performs internal synchronization by sampling the falling edge of the start bit with the basic clock. Receive data is latched internally at the 8th rising edge of the basic clock. This is illustrated in figure 10.21.

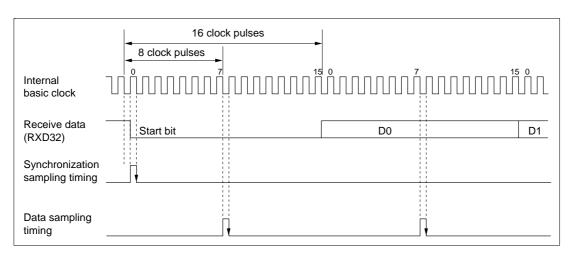


Figure 10.21 Receive Data Sampling Timing in Asynchronous Mode

Consequently, the receive margin in asynchronous mode can be expressed as shown in equation (1).

$$M = \{(0.5 - \frac{1}{2N}) - \frac{D - 0.5}{N} - (L - 0.5) F\} \times 100 [\%]$$
 Equation (1)

where

M: Receive margin (%)
N: Ratio of bit rate to clock (N = 16)
D: Clock duty (D = 0.5 to 1.0)
L: Frame length (L = 9 to 12)
F: Absolute value of clock frequency deviation

Substituting 0 for F (absolute value of clock frequency deviation) and 0.5 for D (clock duty) in equation (1), a receive margin of 46.875% is given by equation (2).

When D = 0.5 and F = 0, M = $\{0.5 - 1/(2 \times 16)\} \times 100$ [%] = 46.875% Equation (2)

However, this is only a computed value, and a margin of 20% to 30% should be allowed when carrying out system design.

7. Relation between RDR reads and bit RDRF

In a receive operation, SCI3 continually checks the RDRF flag. If bit RDRF is cleared to 0 when reception of one frame ends, normal data reception is completed. If bit RDRF is set to 1, this indicates that an overrun error has occurred.

When the contents of RDR are read, bit RDRF is cleared to 0 automatically. Therefore, if bit RDR is read more than once, the second and subsequent read operations will be performed while bit RDRF is cleared to 0. Note that, when an RDR read is performed while bit RDRF is cleared to 0, if the read operation coincides with completion of reception of a frame, the next frame of data may be read. This is illustrated in figure 10.22.

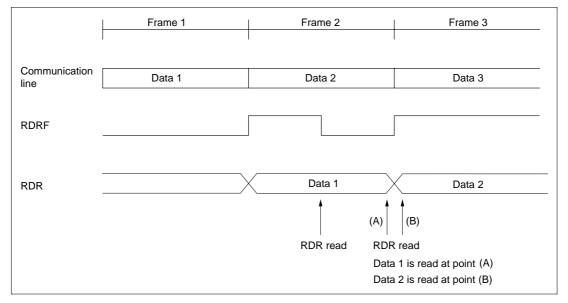


Figure 10.22 Relation between RDR Read Timing and Data

In this case, only a single RDR read operation (not two or more) should be performed after first checking that bit RDRF is set to 1. If two or more reads are performed, the data read the first time should be transferred to RAM, etc., and the RAM contents used. Also, ensure that there is sufficient margin in an RDR read operation before reception of the next frame is completed. To be precise in terms of timing, the RDR read should be completed before bit 7 is transferred in synchronous mode, or before the STOP bit is transferred in asynchronous mode.

8. Transmit and receive operations when making a state transition

Make sure that transmit and receive operations have completely finished before carrying out state transition processing.

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9. Switching SCK₃₂ function

If pin SCK₃₂ is used as a clock output pin by SCI3 in synchronous mode and is then switched to a general input/output pin (a pin with a different function), the pin outputs a low level signal for half a system clock (ϕ) cycle immediately after it is switched.

This can be prevented by either of the following methods according to the situation.

a. When an SCK₃₂ function is switched from clock output to non clock-output

When stopping data transfer, issue one instruction to clear bits TE and RE to 0 and to set bits CKE1 and CKE0 in SCR3 to 1 and 0, respectively. In this case, bit COM in SMR should be left 1. The above prevents SCK_{32} from being used as a general input/output pin. To avoid an intermediate level of voltage from being applied to SCK_{32} , the line connected to SCK_{32} should be pulled up to the V_{CC} level via a resistor, or supplied with output from an external device.

b. When an SCK₃₂ function is switched from clock output to general input/output

When stopping data transfer,

- (i) Issue one instruction to clear bits TE and RE to 0 and to set bits CKE1 and CKE0 in SCR3 to 1 and 0, respectively.
- (ii) Clear bit COM in SMR to 0
- (iii) Clear bits CKE1 and CKE0 in SCR3 to 0

Note that special care is also needed here to avoid an intermediate level of voltage from being applied to SCK₃₂.

10. Set up at subactive or subsleep mode

At subactive or subsleep mode, SCI3 becomes possible use only at CPU clock is øw/2.

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Section 11 10-Bit PWM

11.1 Overview

The H8/3802 Series is provided with two on-chip 10-bit PWMs (pulse width modulators), designated PWM1 and PWM2, with identical functions. The PWMs can be used as D/A converters by connecting a low-pass filter. In this section the suffix m (m = 1 or 2) is used with register names, etc., as in PWDRLm, which denotes the PWDRL registers for each PWM.

11.1.1 Features

Features of the 10-bit PWMs are as follows.

- Choice of four conversion periods Any of the following conversion periods can be chosen: 4,096/ø, with a minimum modulation width of 4/ø 2,048/ø, with a minimum modulation width of 2/ø 1,024/ø, with a minimum modulation width of 1/ø 512/ø, with a minimum modulation width of 1/2 ø
- Pulse division method for less ripple
- Use of module standby mode enables this module to be placed in standby mode independently when not used.

11.1.2 Block Diagram

Figure 11.1 shows a block diagram of the 10-bit PWM.

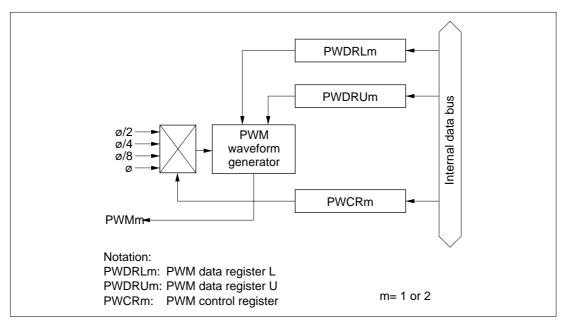


Figure 11.1 Block Diagram of the 10 bit PWM

11.1.3 Pin Configuration

Table 11.1 shows the output pin assigned to the 10-bit PWM.

Table 11.1 Pin Configuration

Name	Abbrev.	I/O	Function
PWM1 output pin	PWM1	Output	Pulse-division PWM waveform output (PWM1)
PWM2 output pin	PWM2	Output	Pulse-division PWM waveform output (PWM2)

11.1.4 Register Configuration

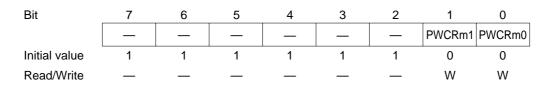
Table 11.2 shows the register configuration of the 10-bit PWM.

Table 11.2 Register Configuration

Name	Abbrev.	R/W	Initial Value	Address
PWM1 control register	PWCR1	W	H'FC	H'FFD0
PWM1 data register U	PWDRU1	W	H'FC	H'FFD1
PWM1 data register L	PWDRL1	W	H'00	H'FFD2
PWM2 control register	PWCR2	W	H'FC	H'FFCD
PWM2 data register U	PWDRU2	W	H'FC	H'FFCE
PWM2 data register L	PWDRL2	W	H'00	H'FFCF
Clock stop register 2	CKSTPR2	R/W	H'FF	H'FFFB

11.2 Register Descriptions

11.2.1 PWM Control Register (PWCRm)



PWCRm is an 8-bit write-only register for input clock selection.

Upon reset, PWCRm is initialized to H'FC.

Bits 7 to 2: Reserved bits

Bits 7 to 2 are reserved; they are always read as 1, and cannot be modified.

Bits 1 and 0: Clock select 1 and 0 (PWCRm1, PWCRm0)

Bits 1 and 0 select the clock supplied to the 10-bit PWM. These bits are write-only bits; they are always read as 1.

Bit 1 PWCRm1	Bit 0 PWCRm0	Description	
0	0	The input clock is \emptyset (t \emptyset * = 1/ \emptyset) The conversion period is 512/ \emptyset , with a minimum modulation width of 1/2 \emptyset	(initial value)
0	1	The input clock is $\emptyset/2$ (tø* = 2/ \emptyset) The conversion period is 1,024/ \emptyset , with a minimum modulation width of 1/ \emptyset	
1	0	The input clock is $\emptyset/4$ (t $\emptyset^* = 4/\emptyset$) The conversion period is 2,048/ \emptyset , with a minimum modulation width of 2/ \emptyset	
1	1	The input clock is $\emptyset/8$ (t $\emptyset^* = 8/\emptyset$) The conversion period is 4,096/ \emptyset , with a minimum modulation width of 4/ \emptyset	

Note: * Period of PWM input clock.

PWDRUm								
Bit	7	6	5	4	3	2	1	0
	—	—	—	—	_	_	PWDRUm ₁	PWDRUm ₀
Initial value	1	1	1	1	1	1	0	0
Read/Write	—	—	—	—	—	—	W	W
PWDRLm								
Bit	7	6	5	4	3	2	1	0
	PWDRLm7	PWDRLm ₆	PWDRLm5	PWDRLm ₄	PWDRLm ₃	PWDRLm ₂	PWDRLm ₁	PWDRLm0
Initial value	0	0	0	0	0	0	0	0
Read/Write	W	W	W	W	W	W	W	W

11.2.2 PWM Data Registers U and L (PWDRUm, PWDRLm)

PWDRUm and PWDRLm form a 10-bit write-only register, with the upper 2 bits assigned to PWDRUm and the lower 8 bits to PWDRLm. The value written to PWDRUm and PWDRLm gives the total high-level width of one PWM waveform cycle.

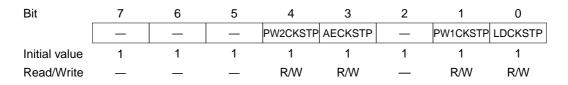
When 10-bit data is written to PWDRUm and PWDRLm, the register contents are latched in the PWM waveform generator, updating the PWM waveform generation data. The 10-bit data should always be written in the following sequence:

- 1. Write the lower 8 bits to PWDRLm.
- 2. Write the upper 2 bits to PWDRUm for the same channel.

PWDRUm and PWDRLm are write-only registers. If they are read, all bits are read as 1.

Upon reset, PWDRUm is initialized to H'FC, and PWDRLm to H'00.

11.2.3 Clock Stop Register 2 (CKSTPR2)



CKSTPR2 is an 8-bit read/write register that performs module standby mode control for peripheral modules. Only the bit relating to the PWM is described here. For details of the other bits, see the sections on the relevant modules.

Bits 4 and 1: PWM module standby mode control (PWmCKSTP)

Bits 4 and 1 control setting and clearing of module standby mode for the PWMm.

PWn	PWmCKSTP Description							
0	PWMm is set to module standby mode							
1	PWMm module standby mode is cleared	(initial value)						

11.3 Operation

11.3.1 Operation

When using the 10-bit PWM, set the registers in the following sequence.

- 1. Set PWM1 or PWM2 in PMR9 to 1 for the PWM channel to be used, so that pin P9₀/PWM1 or P9₁/PWM2 is designated as the PWM output pin.
- Set bits PWCRm1 and PWCRm0 in the PWM control register (PWCRm) to select a conversion period of 4,096/ø (PWCRm1 = 1, PWCRm0 = 1), 2,048/ø (PWCRm1 = 1, PWCRm0 = 0), 1,024/ø (PWCRm1 = 0, PWCRm0 = 1), or 512/ø (PWCRm1 = 0, PWCRm0 = 0).
- 3. Set the output waveform data in PWDRUm and PWDRLm. Be sure to write in the correct sequence, first PWDRLm then PWDRUm for the same channel. When data is written to PWDRUm, the data will be latched in the PWM waveform generator, updating the PWM waveform generation in synchronization with internal signals.

One conversion period consists of 4 pulses, as shown in figure 11.2. The total of the high-level pulse widths during this period (T_H) corresponds to the data in PWDRUm and PWDRLm. This relation can be represented as follows.

 $T_{\rm H}$ = (data value in PWDRUm and PWDRLm + 4) × t_o/2

where tø is the PWM input clock period: $1/\phi$ (PWCRm = H'0), $2/\phi$ (PWCRm = H'1), $4/\phi$ (PWCRm = H'2), or $8/\phi$ (PWCRm = H'3).

Example: Settings in order to obtain a conversion period of 1,024 μ s: When PWCRm1 = 0 and PWCRm0 = 0, the conversion period is 512/ø, so ø must be 0.5 MHz. In this case, tfn = 256 μ s, with 1/2ø (resolution) = 1.0 μ s. When PWCRm1 = 0 and PWCRm0 = 1, the conversion period is 1,024/ø, so ø must be 1 MHz. In this case, tfn = 256 μ s, with 1/ø (resolution) = 1.0 μ s. When PWCRm1 = 1 and PWCRm0 = 0, the conversion period is 2,048/ø, so ø must be 2 MHz. In this case, tfn = 256 μ s, with 2/ø (resolution) = 1.0 μ s. When PWCRm1 = 1 and PWCRm0 = 1, the conversion period is 4,096/ø, so ø must be 4 MHz. In this case, tfn = 256 μ s, with 4/ø (resolution) = 1.0 μ s. Accordingly, for a conversion period of 1,024 μ s, the system clock frequency (ø) must be 0.5 MHz, 1 MHz, 2 Mhz, or 4MHz.

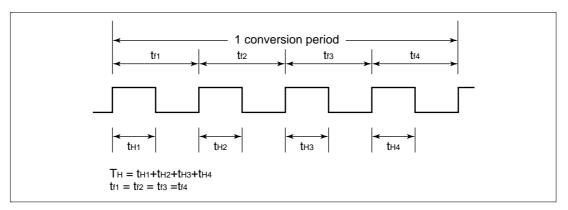


Figure 11.2 PWM Output Waveform

11.3.2 PWM Operation Modes

PWM operation modes are shown in table 11.3.

Table 11.3 PWM Operation Modes

Operation Mode	Reset	Active	Sleep	Watch	Subactive	Subsleep	Standby	Module Standby
PWCRm	Reset	Functions	Functions	Held	Held	Held	Held	Held
PWDRUm	Reset	Functions	Functions	Held	Held	Held	Held	Held
PWDRLm	Reset	Functions	Functions	Held	Held	Held	Held	Held

Section 12 A/D Converter

12.1 Overview

The H8/3802 Series includes on-chip a resistance-ladder-based successive-approximation analog-to-digital converter, and can convert up to 4 channels of analog input.

12.1.1 Features

The A/D converter has the following features.

- 10-bit resolution
- Four input channels
- Conversion time: approx. 12.4 µs per channel (at 5 MHz operation)
- Built-in sample-and-hold function
- Interrupt requested on completion of A/D conversion
- Use of module standby mode enables this module to be placed in standby mode independently when not used.

12.1.2 Block Diagram

Figure 12.1 shows a block diagram of the A/D converter.

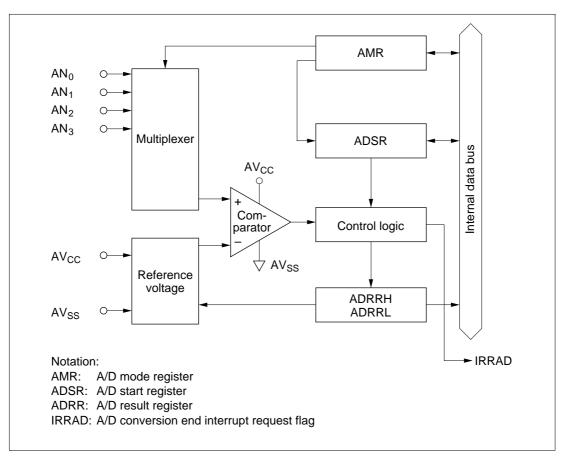


Figure 12.1 Block Diagram of the A/D Converter

12.1.3 Pin Configuration

Table 12.1 shows the A/D converter pin configuration.

Table 12.1 Pin Configuration

Abbrev.	I/O	Function
AV_{cc}	Input	Power supply and reference voltage of analog part
AV _{ss}	Input	Ground and reference voltage of analog part
AN ₀	Input	Analog input channel 0
AN ₁	Input	Analog input channel 1
AN ₂	Input	Analog input channel 2
AN ₃	Input	Analog input channel 3
	$ AV_{cc} AV_{ss} AN_0 AN_1 AN_2 $	AV _{cc} InputAV _{ss} InputAN ₀ InputAN ₁ InputAN ₂ Input

12.1.4 Register Configuration

Table 12.2 shows the A/D converter register configuration.

Table 12.2 Register Configuration

Name	Abbrev.	R/W	Initial Value	Address
A/D mode register	AMR	R/W	H'30	H'FFC6
A/D start register	ADSR	R/W	H'7F	H'FFC7
A/D result register H	ADRRH	R	Not fixed	H'FFC4
A/D result register L	ADRRL	R	Not fixed	H'FFC5
Clock stop register 1	CKSTPR1	R/W	H'FF	H'FFFA

12.2 Register Descriptions

Bit	7	6	5	4	3	2	1	0	7	6	5	4	3	2	1	0
	ADR9	ADR8	ADR7	ADR6	ADR5	ADR4	ADR3	ADR2	ADR1	ADR0	—	_	_	—	-	—
Initial value	Not	_	_	_	_	_	_									
	fixed															
Read/Write	R	R	R	R	R	R	R	R	R	R	—	—	—	—	—	—
					~								~			
				AD	RRH							AD	RRL			

12.2.1 A/D Result Registers (ADRRH, ADRRL)

ADRRH and ADRRL together comprise a 16-bit read-only register for holding the results of analog-to-digital conversion. The upper 8 bits of the data are held in ADRRH, and the lower 2 bits in ADRRL.

ADRRH and ADRRL can be read by the CPU at any time, but the ADRRH and ADRRL values during A/D conversion are not fixed. After A/D conversion is complete, the conversion result is stored as 10-bit data, and this data is held until the next conversion operation starts.

ADRRH and ADRRL are not cleared on reset.

12.2.2 A/D Mode Register (AMR)

Bit	7	6	5	4	3	2	1	0
	CKS	_	_	_	СНЗ	CH2	CH1	CH0
Initial value	0	0	1	1	0	0	0	0
Read/Write	R/W	R/W	_	_	R/W	R/W	R/W	R/W

AMR is an 8-bit read/write register for specifying the A/D conversion speed, external trigger option, and the analog input pins.

Upon reset, AMR is initialized to H'30.

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Bit 7: Clock select (CKS)

Bit 7 sets the A/D conversion speed.

Bit 7		Conversion Time					
CKS	Conversion Period	ø = 1 MHz	ø = 5 MHz				
0	62/ø (initial value)	62 µs	12.4 µs				
1	31/ø	31 µs	*				

Note: * Operation is not guaranteed if the conversion time is less than 12.4 µs. Set bit 7 for a value of at least 12.4 µs.

Bit 6: Reserved bit

Bit 6 is reserved; only 0 can be written to this bit.

Bits 5 and 4: Reserved bits

Bits 5 and 4 are reserved; they are always read as 1, and cannot be modified.

Bits 3 to 0: Channel select (CH3 to CH0)

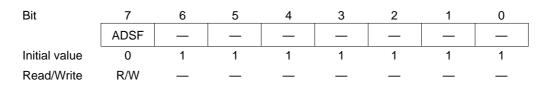
Bits 3 to 0 select the analog input channel.

The channel selection should be made while bit ADSF is cleared to 0.

Bit 3 CH3	Bit 2 CH2	Bit 1 CH1	Bit 0 CH0	Analog Input Channel	
0	0	*	*	No channel selected	(initial value)
0	1	0	0	AN _o	
0	1	0	1	AN ₁	
0	1	1	0	AN ₂	
0	1	1	1	AN ₃	
1	0	0	0	Setting prohibited	
1	0	0	1		
1	0	1	0		
1	0	1	1		

*: Don't care

12.2.3 A/D Start Register (ADSR)



The A/D start register (ADSR) is an 8-bit read/write register for starting and stopping A/D conversion.

A/D conversion is started by writing 1 to the A/D start flag (ADSF), which also sets ADSF to 1. When conversion is complete, the converted data is set in ADRRH and ADRRL, and at the same time ADSF is cleared to 0.

Bit 7: A/D start flag (ADSF)

Bit 7 controls and indicates the start and end of A/D conversion.

Bit 7 ADSF	Description	
0	Read: Indicates the completion of A/D conversion	(initial value)
	Write: Stops A/D conversion	
1	Read: Indicates A/D conversion in progress	
	Write: Starts A/D conversion	

Bits 6 to 0: Reserved bits

Bits 6 to 0 are reserved; they are always read as 1, and cannot be modified.

12.2.4 Clock Stop Register 1 (CKSTPR1)

Bit	7	6	5	4	3	2	1	0
	_	—	S32CKSTP	ADCKSTP	—	TFCKSTP	_	TACKSTP
Initial value	1	1	1	1	1	1	1	1
Read/Write	_	_	R/W	R/W	_	R/W	_	R/W

CKSTPR1 is an 8-bit read/write register that performs module standby mode control for peripheral modules. Only the bit relating to the A/D converter is described here. For details of the other bits, see the sections on the relevant modules.

Bit 4: A/D converter module standby mode control (ADCKSTP)

Bit 4 controls setting and clearing of module standby mode for the A/D converter.

ADCKSTP	Description	
0	A/D converter is set to module standby mode	
1	A/D converter module standby mode is cleared	(initial value)

12.3 Operation

12.3.1 A/D Conversion Operation

The A/D converter operates by successive approximations, and yields its conversion result as 10bit data.

A/D conversion begins when software sets the A/D start flag (bit ADSF) to 1. Bit ADSF keeps a value of 1 during A/D conversion, and is cleared to 0 automatically when conversion is complete.

The completion of conversion also sets bit IRRAD in interrupt request register 2 (IRR2) to 1. An A/D conversion end interrupt is requested if bit IENAD in interrupt enable register 2 (IENR2) is set to 1.

If the conversion time or input channel needs to be changed in the A/D mode register (AMR) during A/D conversion, bit ADSF should first be cleared to 0, stopping the conversion operation, in order to avoid malfunction.

12.3.2 A/D Converter Operation Modes

A/D converter operation modes are shown in table 12.3.

Table 12.3	A/D Converter Operation Modes
-------------------	-------------------------------

Operation Mode	Reset	Active	Sleep	Watch	Subactive	Subsleep	Standby	Module Standby
AMR	Reset	Functions	Functions	Held	Held	Held	Held	Held
ADSR	Reset	Functions	Functions	Held	Held	Held	Held	Held
ADRRH	Held*	Functions	Functions	Held	Held	Held	Held	Held
ADRRL	Held*	Functions	Functions	Held	Held	Held	Held	Held

Note: * Undefined in a power-on reset.

12.4 Interrupts

When A/D conversion ends (ADSF changes from 1 to 0), bit IRRAD in interrupt request register 2 (IRR2) is set to 1.

A/D conversion end interrupts can be enabled or disabled by means of bit IENAD in interrupt enable register 2 (IENR2).

For further details see 3.3, Interrupts.

12.5 Typical Use

An example of how the A/D converter can be used is given below, using channel 1 (pin AN1) as the analog input channel. Figure 12.2 shows the operation timing.

- Bits CH3 to CH0 of the A/D mode register (AMR) are set to 0101, making pin AN₁ the analog input channel. A/D interrupts are enabled by setting bit IENAD to 1, and A/D conversion is started by setting bit ADSF to 1.
- 2. When A/D conversion is complete, bit IRRAD is set to 1, and the A/D conversion result is stored is stored in ADRRH and ADRRL. At the same time ADSF is cleared to 0, and the A/D converter goes to the idle state.
- 3. Bit IENAD = 1, so an A/D conversion end interrupt is requested.
- 4. The A/D interrupt handling routine starts.
- 5. The A/D conversion result is read and processed.
- 6. The A/D interrupt handling routine ends.

If ADSF is set to 1 again afterward, A/D conversion starts and steps 2 through 6 take place.

Figures 12.3 and 12.4 show flow charts of procedures for using the A/D converter.

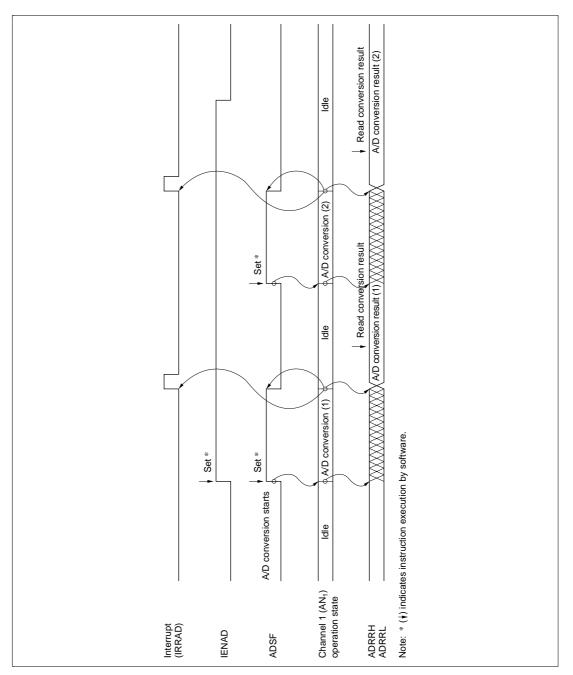


Figure 12.2 Typical A/D Converter Operation Timing

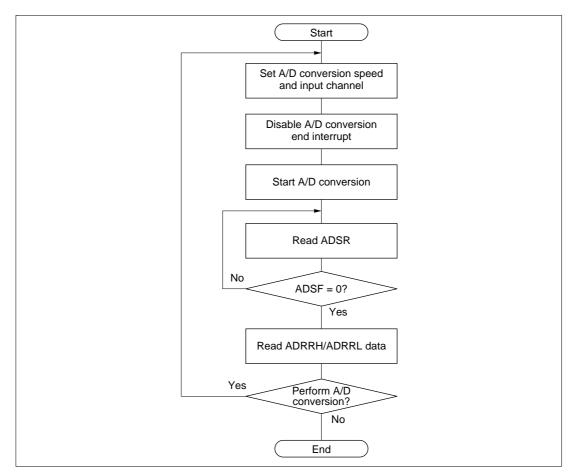


Figure 12.3 Flow Chart of Procedure for Using A/D Converter (Polling by Software)

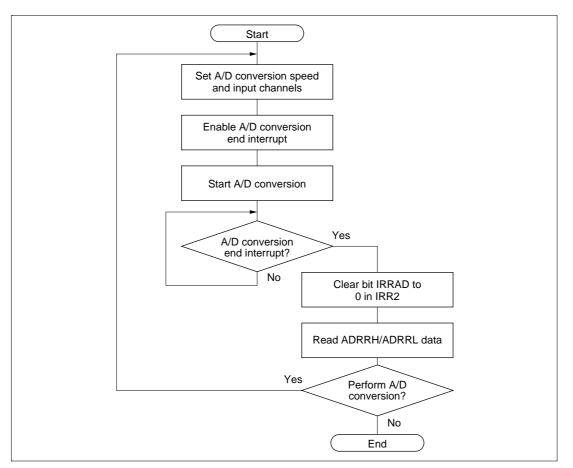


Figure 12.4 Flow Chart of Procedure for Using A/D Converter (Interrupts Used)

12.6 Application Notes

- Data in ADRRH and ADRRL should be read only when the A/D start flag (ADSF) in the A/D start register (ADSR) is cleared to 0.
- Changing the digital input signal at an adjacent pin during A/D conversion may adversely affect conversion accuracy.
- When A/D conversion is started after clearing module standby mode, wait for 10 ø clock cycles before starting.
- In active mode and sleep mode, the analog power supply current (AI_{STOP1}) flows in the ladder resistance even when the A/D converter is on standby. Therefore, if the A/D converter is not used, it is recommended that AV_{CC} be connected to the system power supply and the ADCKSTP (A/D converter module standby mode control) bit be cleared to 0 in clock stop register 1 (CKSTPR1).

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Section 13 LCD Controller/Driver

13.1 Overview

The H8/3802 Series has an on-chip segment type LCD control circuit, LCD driver, and power supply circuit, enabling it to directly drive an LCD panel.

13.1.1 Features

1. Features

Features of the LCD controller/driver are given below.

• Display capacity

Duty Cycle	Internal Driver
Static	25 seg
1/2	25 seg
1/3	25 seg
1/4	25 seg

• LCD RAM capacity

8 bits \times 13 bytes (104 bits)

- Word access to LCD RAM
- All four segment output pins can be used individually as port pins.
- Common output pins not used because of the duty cycle can be used for common doublebuffering (parallel connection).
- Display possible in operating modes other than standby mode
- Choice of 11 frame frequencies
- Built-in power supply split-resistance, supplying LCD drive power
- Use of module standby mode enables this module to be placed in standby mode independently when not used.
- A or B waveform selectable by software

13.1.2 Block Diagram

Figure 13.1 shows a block diagram of the LCD controller/driver.

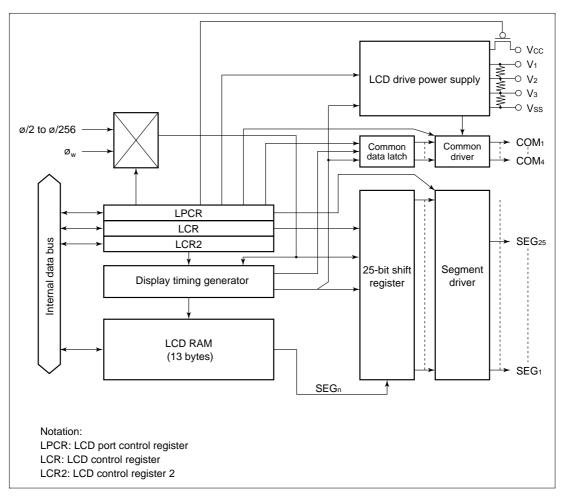


Figure 13.1 Block Diagram of LCD Controller/Driver

13.1.3 Pin Configuration

Table 13.1 shows the LCD controller/driver pin configuration.

Table 13.1 Pin Configuration

Name	Abbrev.	I/O	Function
Segment output pins	SEG_{25} to SEG_1	Output	LCD segment drive pins All pins are multiplexed as port pins (setting programmable)
Common output pins	COM ₄ to COM ₁	Output	LCD common drive pins Pins can be used in parallel with static or 1/2 duty
LCD power supply pins	V ₁ , V ₂ , V ₃		Used when a bypass capacitor is connected externally, and when an external power supply circuit is used

13.1.4 Register Configuration

Table 13.2 shows the register configuration of the LCD controller/driver.

Table 13.2 LCD Controller/Driver Registers

Name	Abbrev.	R/W	Initial Value	Address
LCD port control register	LPCR	R/W	—	H'FFC0
LCD control register	LCR	R/W	H'80	H'FFC1
LCD control register 2	LCR2	R/W	_	H'FFC2
LCD RAM		R/W	Undefined	H'F740 to H'F74C
Clock stop register 2	CKSTPR2	R/W	H'FF	H'FFFB

13.2 Register Descriptions

Bit	7	6	5	4	3	2	1	0
	DTS1	DTS0	CMX		SGS3	SGS2	SGS1	SGS0
Initial value	0	0	0		0	0	0	0
Read/Write	R/W	R/W	R/W	W	R/W	R/W	R/W	R/W

13.2.1 LCD Port Control Register (LPCR)

LPCR is an 8-bit read/write register which selects the duty cycle and LCD driver pin functions.

Bits 7 to 5: Duty cycle select 1 and 0 (DTS1, DTS0), common function select (CMX)

The combination of DTS1 and DTS0 selects static, 1/2, 1/3, or 1/4 duty. CMX specifies whether or not the same waveform is to be output from multiple pins to increase the common drive power when not all common pins are used because of the duty setting.

Bit 7 DTS1	Bit 6 DTS0	Bit 5 CMX	Duty Cycle	Common Drivers	Notes
0	0	0	Static	COM ₁ (initial value)	Do not use COM_4 , COM_3 , and COM_2 .
		1	_	COM ₄ to COM ₁	COM_4 , COM_3 , and COM_2 output the same waveform as COM_1 .
0	1	0	1/2 duty	COM ₂ to COM ₁	Do not use COM_4 and COM_3 .
		1	_	COM ₄ to COM ₁	COM_4 outputs the same waveform as COM_3 , and COM_2 outputs the same waveform as COM_1 .
1	0	0	1/3 duty	COM ₃ to COM ₁	Do not use COM ₄ .
		1		COM ₄ to COM ₁	Do not use COM ₄ .
1	1	0	1/4 duty	COM_4 to COM_1	_
		1			

Bit 4: Reserved bit

Bit 4 is reserved; only 0 can be written to this bit.

Bits 3 to 0: Segment driver select 3 to 0 (SGS3 to SGS0)

Bits 3 to 0 select the segment drivers to be used.

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											-
Bit 3 SGS3	Bit 2 SGS2	Bit 1 SGS1	Bit 0 SGS0	SEG ₂₅	SEG_{24} to SEG_{21}	SEG ₂₀ to SEG ₁₇	SEG ₁₆ to SEG ₁₃	SEG ₁₂ to SEG ₉	SEG_8 to SEG ₅	SEG₄ to SEG₁	Notes
0	0	0	0	Port	Port	Port	Port	Port	Port	Port	(Initial value)
			1	Port	Port	Port	Port	Port	Port	SEG	
		1	0	Port	Port	Port	Port	Port	SEG	SEG	_
			1	Port	Port	Port	Port	SEG	SEG	SEG	_
	1	0	0	Port	Port	Port	SEG	SEG	SEG	SEG	_
			1	Port	Port	SEG	SEG	SEG	SEG	SEG	
		1	0	Port	SEG	SEG	SEG	SEG	SEG	SEG	_
			1	SEG	SEG	SEG	SEG	SEG	SEG	SEG	_
1	0	0	0	SEG	SEG	SEG	SEG	SEG	SEG	SEG	
			1	SEG	SEG	SEG	SEG	SEG	SEG	Port	_
		1	0	SEG	SEG	SEG	SEG	SEG	Port	Port	_
			1	SEG	SEG	SEG	SEG	Port	Port	Port	
	1	0	0	SEG	SEG	SEG	Port	Port	Port	Port	
			1	SEG	SEG	Port	Port	Port	Port	Port	_
		1	0	SEG	Port	Port	Port	Port	Port	Port	_
			1	Port	Port	Port	Port	Port	Port	Port	

Function of Pins SEG₂₅ to SEG₁

13.2.2 LCD Control Register (LCR)

Bit	7	6	5	4	3	2	1	0
	_	PSW	ACT	DISP	CKS3	CKS2	CKS1	CKS0
Initial value	1	0	0	0	0	0	0	0
Read/Write	_	R/W	R/W	R/W	R/W	R/W	R/W	R/W

LCR is an 8-bit read/write register which performs LCD drive power supply on/off control and display data control, and selects the frame frequency.

LCR is initialized to H'80 upon reset.

Bit 7: Reserved bit

Bit 7 is reserved; it is always read as 1 and cannot be modified.

Bit 6: LCD drive power supply on/off control (PSW)

Bit 6 can be used to turn the LCD drive power supply off when LCD display is not required in a power-down mode, or when an external power supply is used. When the ACT bit is cleared to 0, or in standby mode, the LCD drive power supply is turned off regardless of the setting of this bit.

Bit 6 PSW	Description	
0	LCD drive power supply off	(initial value)
1	LCD drive power supply on	

Bit 5: Display function activate (ACT)

Bit 5 specifies whether or not the LCD controller/driver is used. Clearing this bit to 0 halts operation of the LCD controller/driver. The LCD drive power supply is also turned off, regardless of the setting of the PSW bit. However, register contents are retained.

Bit 5 ACT	Description	
0	LCD controller/driver operation halted	(initial value)
1	LCD controller/driver operates	

Bit 4: Display data control (DISP)

Bit 4 specifies whether the LCD RAM contents are displayed or blank data is displayed regardless of the LCD RAM contents.

Bit 4 DISP	Description	
0	Blank data is displayed	(initial value)
1	LCD RAM data is display	

Bits 3 to 0: Frame frequency select 3 to 0 (CKS3 to CKS0)

Bits 3 to 0 select the operating clock and the frame frequency. In subactive mode, watch mode, and subsleep mode, the system clock (ϕ) is halted, and therefore display operations are not performed if one of the clocks from $\phi/2$ to $\phi/256$ is selected. If LCD display is required in these modes, ϕ w, ϕ w/2, or ϕ w/4 must be selected as the operating clock.

Bit 3	Bit 2	Bit 1	Bit 0	Frame Frequency ^{*2}		
CKS3	CKS2	CKS1	CKS0	Operating Clock	ø = 2 MHz	ø = 250 kHz ^{*1}
0	*	0	0	ØW	128 Hz ^{*3} (initia	l value)
0	*	0	1	øw/2	64 Hz ^{*3}	
0	*	1	*	øw/4	32 Hz ^{*3}	
1	0	0	0	ø/2	_	244 Hz
1	0	0	1	ø/4	977 Hz	122 Hz
1	0	1	0	ø/8	488 Hz	61 Hz
1	0	1	1	ø/16	244 Hz	30.5 Hz
1	1	0	0	ø/32	122 Hz	
1	1	0	1	ø/64	61 Hz	
1	1	1	0	ø/128	30.5 Hz	—
1	1	1	1	ø/256		

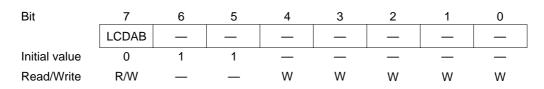
*: Don't care

Notes: 1. This is the frame frequency in active (medium-speed, ϕ osc/16) mode when ϕ = 2 MHz.

2. When 1/3 duty is selected, the frame frequency is 4/3 times the value shown.

3. This is the frame frequency when $\omega w = 32.768 \text{ kHz}$.

13.2.3 LCD Control Register 2 (LCR2)



LCR2 is an 8-bit read/write register which controls switching between the A waveform and B waveform.

Bit 7: A waveform/B waveform switching control (LCDAB)

Bit 7 specifies whether the A waveform or B waveform is used as the LCD drive waveform.

Bit 7 LCDAB	Description	
0	Drive using A waveform	(initial value)
1	Drive using B waveform	

Bits 6 and 5: Reserved bits

Bits 6 and 5 are reserved; they are always read as 1 and cannot be modified.

Bits 4 to 0: Reserved bits

Bits 4 to 0 are reserved; only 0 can be written to these bits.

13.2.4 Clock Stop Register 2 (CKSTPR2)

Bit	7	6	5	4	3	2	1	0
	_	—		PW2CKSTP	AECKSTP	_	PW1CKSTP	LDCKSTP
Initial value	1	1	1	1	1	1	1	1
Read/Write	_	—	_	R/W	R/W	_	R/W	R/W

CKSTPR2 is an 8-bit read/write register that performs module standby mode control for peripheral modules. Only the bit relating to the LCD controller/driver is described here. For details of the other bits, see the sections on the relevant modules.

Bit 0: LCD controller/driver module standby mode control (LDCKSTP)

Bit 0 controls setting and clearing of module standby mode for the LCD controller/driver.

Bit 0 LDCKSTP	Description	
0	LCD controller/driver is set to module standby mode	
1	LCD controller/driver module standby mode is cleared	(initial value)

13.3 Operation

13.3.1 Settings up to LCD Display

To perform LCD display, the hardware and software related items described below must first be determined.

1. Hardware settings

- a. Using 1/2 duty
 - When 1/2 duty is used, interconnect pins V₂ and V₃ as shown in figure 13.2.

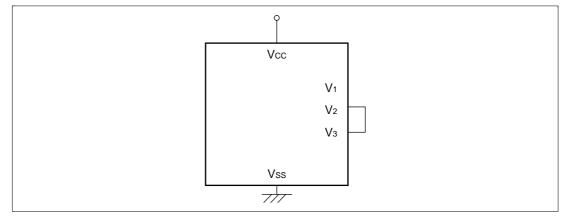


Figure 13.2 Handling of LCD Drive Power Supply when Using 1/2 Duty

b. Large-panel display

As the impedance of the built-in power supply split-resistance is large, it may not be suitable for driving a large panel. If the display lacks sharpness when using a large panel, refer to section 13.3.4, Boosting the LCD Drive Power Supply. When static or 1/2 duty is selected, the common output drive capability can be increased. Set CMX to 1 when selecting the duty cycle. In this mode, with a static duty cycle pins COM_4 to COM_1 output the same waveform, and with 1/2 duty the COM_1 waveform is output from pins COM_2 and COM_1 , and the COM_2 waveform is output from pins COM_4 and COM_3 .

- 2. Software settings
 - a. Duty selection
 Any of four duty cycles—static, 1/2 duty, 1/3 duty, or 1/4 duty—can be selected with bits
 DTS1 and DTS0.
 - b. Segment selection

The segment drivers to be used can be selected with bits SGS₃ to SGS₀.

c. Frame frequency selection

The frame frequency can be selected by setting bits CKS_3 to CKS_0 . The frame frequency should be selected in accordance with the LCD panel specification. For the clock selection method in watch mode, subactive mode, and subsleep mode, see 13.3.3, Operation in Power-Down Modes.

d. A or B waveform selection

Either the A or B waveform can be selected as the LCD waveform to be used by means of LCDAB.

13.3.2 Relationship between LCD RAM and Display

The relationship between the LCD RAM and the display segments differs according to the duty cycle. LCD RAM maps for the different duty cycles are shown in figures 13.3 to 13.6.

After setting the registers required for display, data is written to the part corresponding to the duty using the same kind of instruction as for ordinary RAM, and display is started automatically when turned on. Word- or byte-access instructions can be used for RAM setting.

	bit7	bit6	bit5	bit4	bit3	bit2	bit1	bit0
H'F740	SEG ₂	SEG ₂	SEG ₂	SEG ₂	SEG1	SEG1	SEG1	SEG1
H'F74C					SEG ₂₅	SEG ₂₅	SEG ₂₅	SEG ₂₅
					¢ COM			
	COM ₄	COM ₃	COM ₂	COM ₁	COM ₄	COM ₃	COM ₂	COM ₁

Figure 13.3 LCD RAM Map (1/4 Duty)

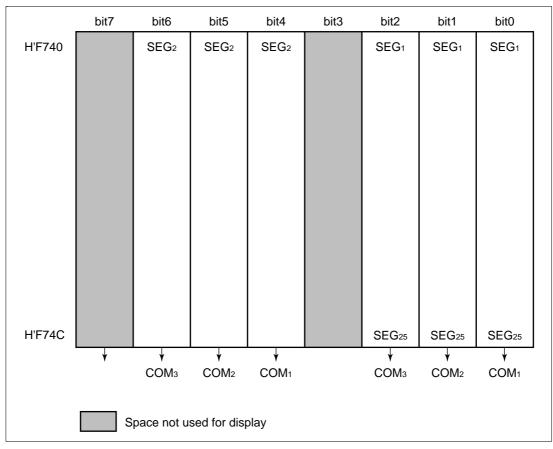
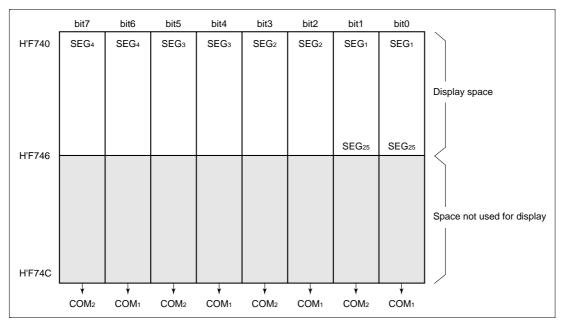
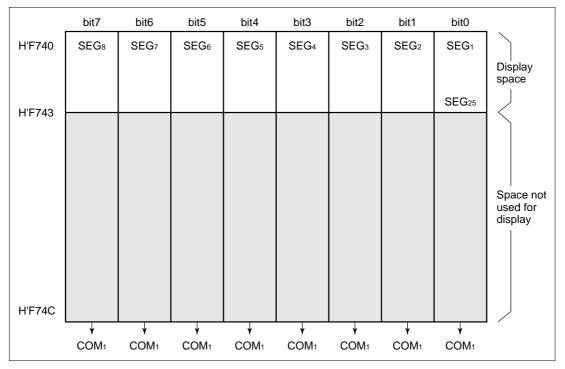


Figure 13.4 LCD RAM Map (1/3 Duty)









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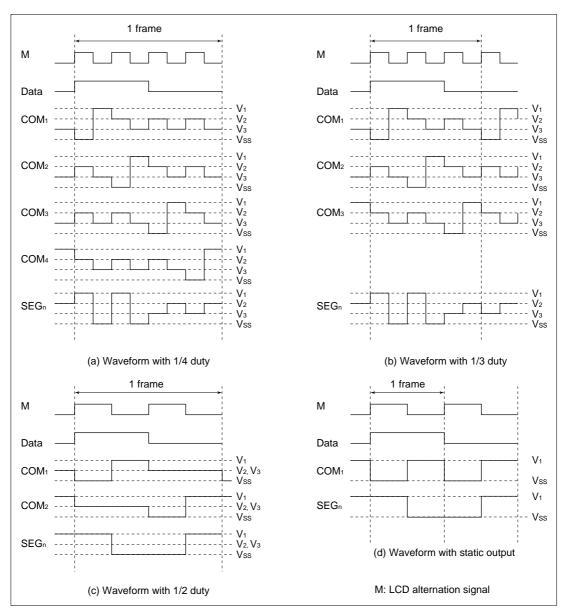


Figure 13.7 Output Waveforms for Each Duty Cycle (A Waveform)

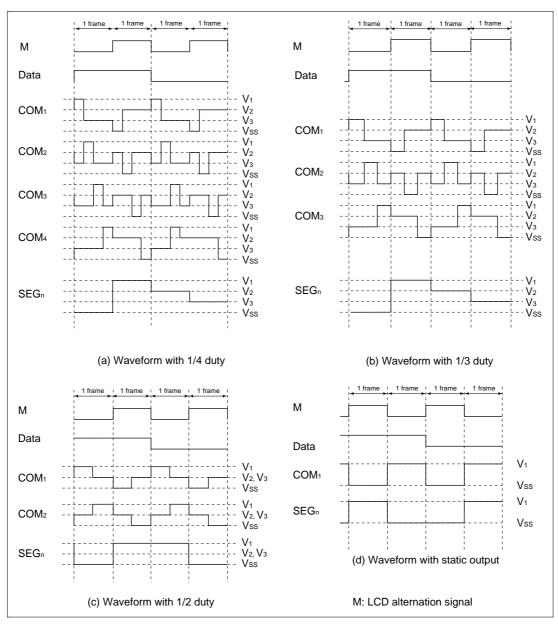


Figure 13.8 Output Waveforms for Each Duty Cycle (B Waveform)

Table 13.3 Output Levels

Data		0	0	1	1
Μ		0	1	0	1
Static	Common output	V ₁	V _{SS}	V ₁	V _{SS}
	Segment output	V ₁	V _{SS}	V _{SS}	V ₁
1/2 duty	Common output	V ₂ , V ₃	V ₂ , V ₃	V ₁	V _{SS}
	Segment output	V ₁	V _{SS}	V _{SS}	V ₁
1/3 duty	Common output	V ₃	V ₂	V ₁	V _{SS}
	Segment output	V ₂	V ₃	V _{SS}	V ₁
1/4 duty	Common output	V ₃	V ₂	V ₁	V _{SS}
	Segment output	V ₂	V ₃	V _{SS}	V ₁

M: LCD alternation signal

13.3.3 Operation in Power-Down Modes

In the H8/3802 Series, the LCD controller/driver can be operated even in the power-down modes. The operating state of the LCD controller/driver in the power-down modes is summarized in table 13.4.

In subactive mode, watch mode, and subsleep mode, the system clock oscillator stops, and therefore, unless ϕ w, ϕ w/2, or ϕ w/4 has been selected by bits CKS3 to CKS0, the clock will not be supplied and display will halt. Since there is a possibility that a direct current will be applied to the LCD panel in this case, it is essential to ensure that ϕ w, ϕ w/2, or ϕ w/4 is selected. In active (medium-speed) mode, the system clock is switched, and therefore CKS3 to CKS0 must be modified to ensure that the frame frequency does not change.

Mode		Reset	Active	Sleep	Watch	Sub- active	Sub- sleep	Standby	Module Standby
Clock	Ø	Runs	Runs	Runs	Stops	Stops	Stops	Stops	Stops*4
	ØW	Runs	Runs	Runs	Runs	Runs	Runs	Stops*1	Stops*4
Display	ACT = 0	Stops	Stops	Stops	Stops	Stops	Stops	Stops*2	Stops
operation	ACT = 1	Stops	Functions	Functions	Functions*	³ Functions ^{**}	³ Functions [*]	³ Stops ^{*2}	Stops

Table 13.4 Power-Down Modes and Display Operation

Notes: 1. The subclock oscillator does not stop, but clock supply is halted.

2. The LCD drive power supply is turned off regardless of the setting of the PSW bit.

- 3. Display operation is performed only if øw, øw/2, or øw/4 is selected as the operating clock.
- 4. The clock supplied to the LCD stops.

13.3.4 Boosting the LCD Drive Power Supply

When a large panel is driven, the on-chip power supply capacity may be insufficient. If the power supply capacity is insufficient when V_{CC} is used as the power supply, the power supply impedance must be reduced. This can be done by connecting bypass capacitors of around 0.1 to 0.3 μ F to pins V₁ to V₃, as shown in figure 13.9, or by adding a split-resistance externally.

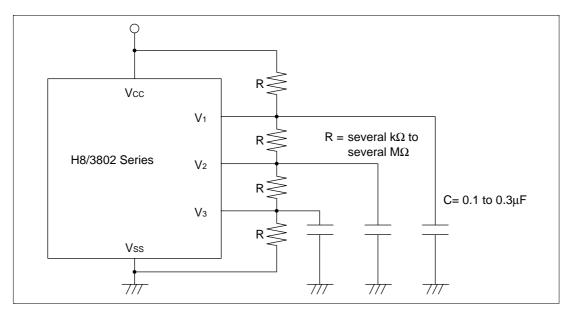


Figure 13.9 Connection of External Split-Resistance

Section 14 Electrical Characteristics

14.1 H8/3802 Series Absolute Maximum Ratings

Table 14.1 lists the absolute maximum ratings.

Table 14.1 Absolute Maximum Ratings

ltem		Symbol	Value	Unit
Power supply v	oltage	V _{CC}	-0.3 to +7.0	V
Analog power s	upply voltage	AV _{CC}	-0.3 to +7.0	V
Programming v	oltage	V _{PP}	-0.3 to +13.0	V
Input voltage	Ports other than Port B, IRQAEC	V _{in}	–0.3 to V _{CC} +0.3	V
	Port B	AV _{in}	–0.3 to AV _{CC} +0.3	V
	IRQAEC	HV _{in}	-0.3 to +7.3	V
Operating temperature		T _{opr}	-20 to +75	°C
Storage temper	ature	T _{stg}	-55 to +125	°C
			-55 to +125	°C

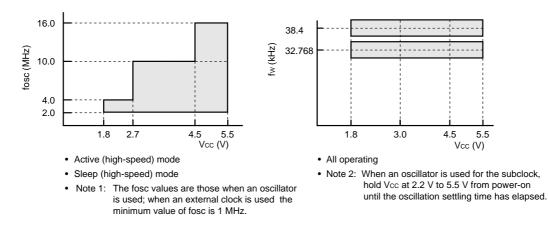
Note: Permanent damage may occur to the chip if maximum ratings are exceeded. Normal operation should be under the conditions specified in Electrical Characteristics. Exceeding these values can result in incorrect operation and reduced reliability.

14.2 H8/3802 Series Electrical Characteristics

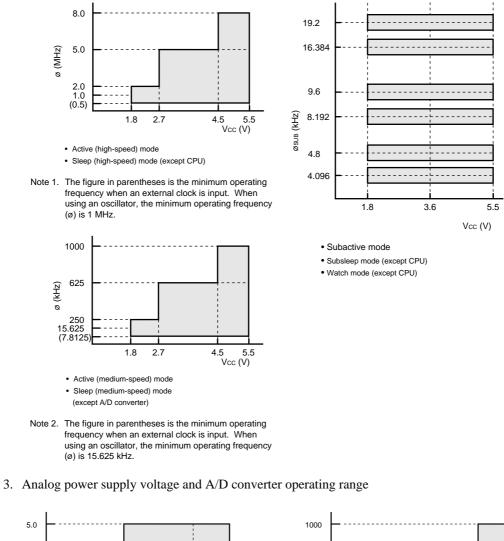
14.2.1 Power Supply Voltage and Operating Range

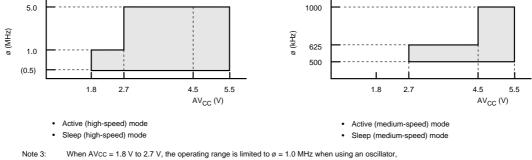
The power supply voltage and operating range are indicated by the shaded region in the figures.

1. Power supply voltage and oscillator frequency range



2. Power supply voltage and operating frequency range





Note 3: When AVCC = 1.8 V to 2.7 V, the operating range is limited to Ø = 1.0 MHz when using an oscillator, and is Ø = 0.5 MHz to 1.0 MHz when using an external clock.

14.2.2 DC Characteristics

Table 14.2 lists the DC characteristics of the H8/3802.

Table 14.2DC Characteristics

 $V_{CC} = 1.8 \text{ V}$ to 5.5 V, $AV_{CC} = 1.8 \text{ V}$ to 5.5 V, $V_{SS} = AV_{SS} = 0.0 \text{ V}$, $T_a = -20^{\circ}\text{C}$ to $+75^{\circ}\text{C}$ (including subactive mode) unless otherwise indicated.

				Value	es		
ltem	Symbol	Applicable Pins	Min	Тур	Max	Unit	Test Condition Notes
Input high	V _{IH}	RES,	0.8 V _{CC}		Vcc + 0.3	V	V_{CC} = 4.0 V to 5.5 V
voltage		$\label{eq:wkp_0} \begin{array}{l} \overline{WKP}_0 \text{ to } \overline{WKP}_7, \\ \overline{IRQ}_0, \overline{IRQ}_1, \\ AEVL, AEVH, \\ SCK_{32} \end{array}$	0.9 V _{CC}	—	V _{CC} + 0.3	_	Except the above
		RXD ₃₂	$0.7 V_{CC}$		V _{CC} + 0.3	V	$V_{CC} = 4.0 \text{ V to } 5.5 \text{ V}$
			0.8 V _{CC}		V _{CC} + 0.3	_	Except the above
		OSC ₁	0.8 V _{CC}		V _{CC} + 0.3	V	V_{CC} = 4.0 V to 5.5 V
			0.9 V _{CC}		V _{CC} + 0.3		Except the above
		X ₁	0.9 V _{CC}		V _{CC} + 0.3	V	V_{CC} = 1.8 V to 5.5 V
		P3 ₁ to P3 ₇ ,	0.7 V _{CC}		V _{CC} + 0.3	V	V_{CC} = 4.0 V to 5.5 V
		$P4_0$ to $P4_3$, $P5_0$ to $P5_7$, $P6_0$ to $P6_7$, $P7_0$ to $P7_7$, $P8_0$, PA_0 to PA_3	0.8 V _{CC}	_	V _{CC} + 0.3		Except the above
		PB ₀ to PB ₃	0.7 V _{CC}		AV _{CC} + 0.3	_	$V_{\rm CC} = 4.0 \text{ V to } 5.5 \text{ V}$
			0.8 V _{CC}		$AV_{CC} + 0.3$	-	Except the above
		IRQAEC	0.8 V _{CC}	_	7.3	V	$V_{CC} = 4.0 \text{ V to } 5.5 \text{ V}$
			0.9 V _{CC}		7.3		Except the above

Note: Connect the TEST pin to V_{SS} .

		_	Values						
ltem	Symbol	Applicable Pins	Min	Тур	Max	Unit	Test Condition	Notes	
Input low	V _{IL}	RES,	-0.3	—	$0.2 V_{CC}$	V	V_{CC} = 4.0 V to 5.5 V		
voltage		$\label{eq:wkp_0} \begin{array}{l} \overline{\text{WKP}}_0 \text{ to } \overline{\text{WKP}}_7, \\ \overline{\text{IRQ}}_0, \overline{\text{IRQ}}_1, \\ \text{IRQAEC, AEVL,} \\ \text{AEVH, SCK}_{32} \end{array}$	-0.3	_	0.1 V _{CC}		Except the above		
		RXD ₃₂ ,	-0.3	_	0.3 V _{CC}	V	V_{CC} = 4.0 V to 5.5 V		
			-0.3	_	$0.2 \ V_{CC}$		Except the above		
		OSC ₁	-0.3	_	0.2 V _{CC}	V	$V_{CC} = 4.0 \text{ V to } 5.5 \text{ V}$		
			-0.3	—	0.1 V _{CC}		Except the above		
		X ₁	-0.3	_	0.1 V _{CC}	V	$V_{\rm CC}$ = 1.8 V to 5.5 V		
		P3 ₁ to P3 ₇ ,	-0.3	_	0.3 V _{CC}	V	V_{CC} = 4.0 V to 5.5 V		
		$P4_0 \text{ to } P4_3,$ $P5_0 \text{ to } P5_7,$ $P6_0 \text{ to } P6_7,$ $P7_0 \text{ to } P7_7,$ $P8_0,$ $PA_0 \text{ to } PA_3,$ $PB_0 \text{ to } PB_3$	-0.3	_	0.2 V _{CC}		Except the above		
Output high voltage	V _{OH}	$P3_1$ to $P3_7$, $P4_0$ to $P4_2$,	V _{CC} – 1.0	_		V	$V_{CC} = 4.0 \text{ V to 5.5 V}$ $-I_{OH} = 1.0 \text{ mA}$	_	
		P50 to P57, P60 to P67, P70 to P77,	$V_{CC} - 0.5$	—	—		$V_{CC} = 4.0 \text{ V to } 5.5 \text{ V}$ $-I_{OH} = 0.5 \text{ mA}$	_	
		P80, PA0 to PA3	V _{CC} – 0.3	—	_		-I _{OH} = 0.1 mA	•	
Output low voltage	V _{OL}	P4 ₀ to P4 ₂	_	_	0.6	V	$V_{CC} = 4.0 \text{ V to } 5.5 \text{ V}$ $I_{OL} = 1.6 \text{ mA}$		
			_	—	0.5		I _{OL} = 0.4 mA		
		$\begin{array}{l} {\sf P5}_0 \text{ to } {\sf P5}_7, {\sf P6}_0 \\ {\sf to } {\sf P6}_7, {\sf P7}_0 \text{ to} \\ {\sf P7}_7, {\sf P8}_0, \\ {\sf PA}_0 \text{ to } {\sf PA}_3 \end{array}$	_		0.5		I _{OL} = 0.4 mA		
		P3 ₁ to P3 ₇			1.5		$V_{CC} = 4.0 \text{ V to } 5.5 \text{ V}$ $I_{OL} = 10 \text{ mA}$		
			_		0.6	_	$V_{CC} = 4.0 \text{ V to } 5.5 \text{ V}$ $I_{OL} = 1.6 \text{ mA}$		
			_	_	0.5	_	I _{OL} = 0.4 mA		

				Value	s			
ltem	Symbol	Applicable Pins	Min	Тур	Max	Unit	Test Condition	Notes
Output low voltage	V _{OL}	P9 ₀ to P9 ₂	_	—	0.5	V	V_{CC} = 2.2 V to 5.5 V, I_{OL} = 25 mA	
							I _{OL} = 15 mA	-
							I _{OL} = 10 mA	*6
		P9 ₃ to P9 ₅	_	_	0.5	_	I _{OL} = 10 mA	
Input/output	I _{IL}	RES, P4 ₃	_	_	20.0	μA	V _{IN} = 0.5 V to	*2
leakage			_	_	1.0		$V_{CC} - 0.5 V$	*1
current		$\begin{array}{c} OSC_{1}, X_{1}, \\ P3_{1} \mbox{to} \ P3_{7}, \\ P4_{0} \mbox{to} \ P4_{2}, \\ P5_{0} \mbox{to} \ P5_{7}, \\ P6_{0} \mbox{to} \ P6_{7}, \\ P7_{0} \mbox{to} \ P7_{7}, \\ P8_{0}, \mbox{IRQAEC}, \\ P9_{0} \mbox{to} \ P9_{5}, \\ PA_{0} \mbox{to} \ PA_{3} \end{array}$	_		1.0	μA	$V_{IN} = 0.5 V to$ $V_{CC} - 0.5 V$	_
		PB ₀ to PB ₃	_	_	1.0		$V_{IN} = 0.5 V \text{ to}$ AV _{CC} - 0.5 V	-
Pull-up MOS	-I _p	$P3_1$ to $P3_7$, $P5_0$ to $P5_7$,	50.0	_	300.0	μA	V _{CC} = 5 V, V _{IN} = 0 V	
current		P6 ₀ to P6 ₇	_	35.0	_		V _{CC} = 2.7 V, V _{IN} = 0 V	Reference value
Input capacitance	C _{IN}	All input pins except power supply, $\overline{\text{RES}}$, IRQAEC, P4 ₃ , PB ₀ to PB ₃	_		15.0	pF	f = 1 MHz, $V_{IN} = 0 \text{ V},$ $T_a = 25^{\circ}\text{C}$	
		IRQAEC	_		30.0			
		RES		_	80.0			*2
			_	_	15.0			*1
		P4 ₃			50.0			*2
			_	_	15.0			*1
		PB ₀ to PB ₃			15.0			

				Value	es			
ltem	Symbol	Applicable Pins	Min	Тур	Max	Unit	Test Condition	Notes
Active mode current dissipation	I _{OPE1}	V _{CC}	_	7.0	10.0	mA	Active (high-speed) mode $V_{CC} = 5 V$, $f_{OSC} = 10 MHz$	*3 *4
	I _{OPE2}	V _{cc}		2.2	3.0	mA	Active (medium- speed) mode $V_{CC} = 5 V$, $f_{OSC} = 10 MHz$ $ø_{osc}/128$	*3 *4
Sleep mode current dissipation	I _{SLEEP}	V _{cc}		3.8	5.0	mA	V _{CC} =5 V, f _{OSC} = 10 MHz	*3 *4
Subactive mode current dissipation	I _{SUB}	V _{cc}	_	15.0	30.0	μA	$V_{CC} = 2.7 V,$ LCD on 32 kHz crystal oscillator ($\phi_{SUB} = \phi_w/2$)	*3 *4
			_	8.0	_	μA	$V_{CC} = 2.7 V,$ LCD on 32 kHz crystal oscillator ($\sigma_{SUB} = \sigma_w/8$)	*3 *4 Reference value
Subsleep mode current dissipation	I _{SUBSP}	V _{cc}	_	7.5	16.0	μA	$V_{CC} = 2.7 V,$ LCD on 32 kHz crystal oscillator ($\phi_{SUB} = \phi_w/2$)	*3 *4
Watch mode current dissipation	I _{WATCH}	V _{cc}	_	3.8	6.0	μA	V _{CC} = 2.7 V 3 2 kHz crystal oscillator LCD not used	*2 *3 *4
				2.8	_			*1 *3 *4
Standby mode current dissipation	I _{STBY}	V _{cc}	_	1.0	5.0	μA	32 kHz crystal oscillator not used	*3 *4
RAM data retaining voltage	V _{RAM}	V _{cc}	1.5		_	V		

				Value	es			
ltem	Symbol	Applicable Pins	Min	Тур	Max	Unit	Test Condition	Notes
Allowable output low current	I _{OL}	Output pins except port 3 and 9	_		2.0	mA	V_{CC} = 4.0 V to 5.5 V	
(per pin)		Port 3	_	_	10.0		$V_{\rm CC}$ = 4.0 V to 5.5 V	
		Output pins except port 9	_		0.5			
Allowable output low	I _{OL}	P9 ₀ to P9 ₂			25.0	mA	$V_{CC} = 2.2 \text{ V to } 5.5 \text{ V}$	*5
current per pin)		_		15.0				
				_	10.0			
		P9 ₃ to P9 ₅	_	_	10.0	_		
Allowable output low current	ΣI_{OL}	Output pins except ports 3 and 9	_		40.0	mA	V_{CC} = 4.0 V to 5.5 V	
(total)		Port 3	_		80.0		$V_{CC} = 4.0 \text{ V to } 5.5 \text{ V}$	
		Output pins except port 9	_	_	20.0			-
		Port 9	_	_	80.0			
Allowable	–I _{ОН}	All output pins		_	2.0	mA	$V_{CC} = 4.0 \text{ V to } 5.5 \text{ V}$	
output high current (per pin)			_		0.2		Except the above	
Allowable	$\Sigma - I_{OH}$	All output pins	_		15.0	mA	$V_{CC} = 4.0 \text{ V to } 5.5 \text{ V}$	
output high			_	_	10.0		Except the above	

Notes: 1. Applies to the Mask ROM products.

2. Applies to the HD6473802.

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3. Pin states during current measurement.

Mode	RES Pin	Internal State	Other Pins	LCD Power Supply	Oscillator Pins
Active (high-speed) mode (I _{OPE1})	V _{CC}	Operates	V _{CC}	Halted	System clock oscillator: crystal
Active (medium- speed) mode (I _{OPE2})	_				Subclock oscillator: Pin X ₁ = GND
Sleep mode	V _{CC}	Only timers operate	V _{CC}	Halted	_
Subactive mode	V _{CC}	Operates	V _{CC}	Halted	System clock oscillator:
Subsleep mode	V_{CC}	Only timers operate, CPU stops	V _{CC}	Halted	crystal Subclock oscillator:
Watch mode	V_{CC}	Only time base operates, CPU stops	V _{CC}	Halted	crystal
Standby mode	V _{CC}	CPU and timers both stop	V _{CC}	Halted	System clock oscillator: crystal Subclock oscillator: Pin $X_1 = GND$

4. Excludes current in pull-up MOS transistors and output buffers.

5. When the PIOFF bit in the port mode register 9 is 0.

6. When the PIOFF bit in the port mode register 9 is 1.

14.2.3 AC Characteristics

Table 14.3 lists the control signal timing, and tables 14.4 lists the serial interface timing of the H8/3802.

Table 14.3 Control Signal Timing

 $V_{CC} = 1.8 \text{ V}$ to 5.5 V, $AV_{CC} = 1.8 \text{ V}$ to 5.5 V, $V_{SS} = AV_{SS} = 0.0 \text{ V}$, $T_a = -20^{\circ}\text{C}$ to +75°C (including subactive mode) unless otherwise indicated.

		Applicable		Values	6			Reference
Item	Symbol	Pins	Min	Тур	Мах	Unit	Test Condition	Figure
System clock	f _{OSC}	OSC_1, OSC_2	2.0	_	16.0	MHz	V_{CC} = 4.5 V to 5.5 V	
oscillation			2.0	_	10.0	_	$V_{\rm CC}$ = 2.7 V to 5.5 V	
frequency			2.0		4.0	_	Except the above	
OSC clock (ø _{OSC}) cycle time	t _{osc}	OSC ₁ , OSC ₂	62.5	_	500 (1000)	ns	$V_{CC} = 4.5 \text{ V to } 5.5 \text{ V}$	Figure 14.1
			100	_	500 (1000)	_	$V_{\rm CC}$ = 2.7 V to 5.5 V	*2
			250	_	500 (1000)	_	Except the above	
System clock (ø)	t _{cyc}		2	_	128	t _{OSC}		
cycle time			_		128	μs	-	
Subclock oscillation frequency	f _W	X ₁ , X ₂		32.768 or 38.4	_	kHz		
Watch clock (ø _W) cycle time	t _W	X ₁ , X ₂		30.5 or 26.0	—	μs		Figure 14.1
Subclock (ø _{SUB}) cycle time	t _{subcyc}		2	_	8	t _W		*1
Instruction cycle time			2	_	_	t _{cyc} t _{subcyc}		
Oscillation stabilization time	t _{rc}	OSC_1, OSC_2	_	20	45	μs	Figure 14.7 V _{CC} = 2.2 V to 5.5 V	Figure 14.7
			_	—	50	ms	Except the above	Figure 14.7

		Applicable		Values	6			Reference
Item	Symbol	Pins	Min	Тур	Max	Unit	Test Condition	Figure
Oscillation stabilization time	t _{rc}	X ₁ , X ₂	—	—	2.0	S	V_{CC} = 2.7 V to 5.5 V	*3
			_	_	10.0		$V_{CC} = 2.7 \text{ V to } 5.5 \text{ V}$	
External clock high	t _{CPH}	OSC ₁	25	_	_	ns	V_{CC} = 4.5 V to 5.5 V	Figure 14.1
width			40			_	$V_{\rm CC}$ = 2.7 V to 5.5 V	•
			100	_		_	Except the above	Figure 14.1
		X ₁	_	15.26 or 13.02	_	μs		-
External clock low	t _{CPL}	OSC ₁	25	_	_	ns	$V_{CC} = 4.5 \text{ V to } 5.5 \text{ V}$	Figure 14.1
width			40			_	$V_{\rm CC}$ = 2.7 V to 5.5 V	
			100	_	_	_	Except the above	Figure 14.1
		X ₁	_	15.26 or 13.02	_	μs		
External clock rise	t _{CPr}	OSC ₁	_		6	ns	$V_{CC} = 4.5 \text{ V to } 5.5 \text{ V}$	Figure 14.1
time			_		10	_	$V_{\rm CC}$ = 2.7 V to 5.5 V	•
			_	_	25	_	Except the above	Figure 14.1
		X ₁			55.0	ns		-
External clock fall	t _{CPf}	OSC1			6	ns	$V_{CC} = 4.5 \text{ V to } 5.5 \text{ V}$	Figure 14.1
time			_		10	_	$V_{\rm CC}$ = 2.7 V to 5.5 V	-
			_		25	_	Except the above	Figure 14.1
		X ₁			55.0	ns		
Pin $\overline{\text{RES}}$ low width	t _{REL}	RES	10	_	_	t _{cyc}		Figure 14.2
Input pin high width	t _{IH}	$\label{eq:relation} \begin{array}{l} \overline{IRQ}_0, \overline{IRQ}_1, \\ IRQAEC, \\ \overline{WKP}_0 \text{ to } \overline{WKP}_7 \end{array}$	2	_	_	t _{cyc} t _{subcyc}		Figure 14.3
		AEVL, AEVH	0.5	_		t _{osc}		
Input pin low width	t _{IL}	$\label{eq:relation} \begin{array}{l} \overline{IRQ}_0 \text{ to } \overline{IRQ}_1, \\ IRQAEC, \\ \overline{WKP}_0 \text{ to } \overline{WKP}_7 \end{array}$	2	—	-	t _{cyc} t _{subcyc}		Figure 14.3
		AEVL, AEVH	0.5	_		t _{osc}		-

Notes: 1. Selected with SA1 and SA0 of system clock control register 2 (SYSCR2).

2. The figure in parentheses applies when an external clock is used.

3. After powering on, hold $\rm V_{cc}$ at 2.2 V to 5.5 V until the chip's oscillation settling time has elapsed.

Table 14.4 Serial Interface (SCI3) Timing

 $V_{CC} = 1.8$ V to 5.5 V, $AV_{CC} = 1.8$ V to 5.5 V, $V_{SS} = AV_{SS} = 0.0$ V, $T_a = -20^{\circ}C$ to $+75^{\circ}C$ (including subactive mode) unless otherwise indicated.

				Values	5			Reference
ltem		Symbol	Min	Тур	Мах	Unit	Test Conditions	Figure
Input clock	Asynchronous	t _{scyc}	4	—		$t_{\rm cyc} {\rm or}$		Figure 14.4
cycle	Synchronous		6	—	_	t _{subcyc}		
Input clock p	oulse width	t _{SCKW}	0.4	—	0.6	t _{scyc}		Figure 14.4
Transmit dat	a delay time	t _{TXD}	—	—	1	t _{cyc} or	V_{CC} = 4.0 V to 5.5 V	Figure 14.5
(synchronou	s)		_	_	1	t _{subcyc}	Except the above	
Receive data	a setup time	t _{RXS}	200.0	—		ns	V_{CC} = 4.0 V to 5.5 V	Figure 14.5
(synchronou	s)		400.0	_	_	_	Except the above	Figure 14.5
Receive data	a hold time	t _{RXH}	200.0		_	ns	V_{CC} = 4.0 V to 5.5 V	Figure 14.5
(synchronou	s)		400.0			_	Except the above	Figure 14.5

14.2.4 A/D Converter Characteristics

Table 14.5 shows the A/D converter characteristics of the H8/3802.

Table 14.5 A/D Converter Characteristics

 V_{CC} = 1.8 V to 5.5 V, V_{SS} = AV_{SS} = 0.0 V, T_a = –20°C to +75°C unless otherwise indicated.

		Applicable		Val	ues			Reference
Item	Symbol	Pins	Min	Тур	Max	Unit	Test Condition	Figure
Analog power supply voltage	AV_{CC}	AV _{CC}	1.8	_	5.5	V		*1
Analog input voltage	AV _{IN}	AN_0 to AN_3	- 0.3	_	AV _{CC} + 0.3	V		
Analog power	Al _{OPE}	AV _{CC}	_	_	1.5	mA	$AV_{CC} = 5 V$	
supply current	AI _{STOP1}	AV _{CC}	_	600		μA		*2
								Reference value
	AI _{STOP2}	AV _{CC}	_	_	5	μA		*3
Analog input capacitance	C _{AIN}	AN ₀ to AN ₃	_	_	15.0	рF		
Allowable signal source impedance	R _{AIN}				10.0	kΩ		
Resolution (data length)			_	_	10	bit		
Nonlinearity error			—	_	±2.5	LSB	$AV_{CC} = 2.7 V \text{ to } 5.5 V$ $V_{CC} = 2.7 V \text{ to } 5.5 V$	
			_		±5.5	_	AV _{CC} = 2.0 V to 5.5 V V _{CC} = 2.0 V to 5.5 V	_
			_		±7.5	_	Except the above	*4
Quantization error			_	_	±0.5	LSB		

		Applicable		Va	lues			Reference
ltem	Symbol	Pins	Min	Тур	Мах	Unit	Test Condition	Figure
Absolute accuracy			_	_	±3.0	LSB	$AV_{CC} = 2.7 V \text{ to } 5.5 V$ $V_{CC} = 2.7 V \text{ to } 5.5 V$	
			_		±6.0		$AV_{CC} = 2.0 V \text{ to } 5.5 V$ $V_{CC} = 2.0 V \text{ to } 5.5 V$	-
			_	_	±8.0		Except the above	*4
Conversion time			12.4	_	124	μs	$AV_{CC} = 2.7 V \text{ to } 5.5 V$ $V_{CC} = 2.7 V \text{ to } 5.5 V$	
			62		124		Except the above	_

Notes: 1. Set $AV_{CC} = V_{CC}$ when the A/D converter is not used.

2. AI_{STOP1} is the current in active and sleep modes while the A/D converter is idle.

3. AI_{STOP2} is the current at reset and in standby, watch, subactive, and subsleep modes while the A/D converter is idle.

4. Conversion time 62 µs

14.2.5 LCD Characteristics

Table 14.6 shows the LCD characteristics.

Table 14.6 LCD Characteristics

 $V_{CC} = 1.8 \text{ V}$ to 5.5 V, $AV_{CC} = 1.8 \text{ V}$ to 5.5 V, $V_{SS} = AV_{SS} = 0.0 \text{ V}$, $T_a = -20^{\circ}\text{C}$ to $+75^{\circ}\text{C}$ (including subactive mode) unless otherwise specified.

		Applicable	Test		Va	lues		Reference
ltem	Symbol	Pins	Conditions	Min	Тур	Max	Unit	Figure
Segment driver drop voltage	V_{DS}	SEG ₁ to SEG ₂₅	$I_D = 2 \ \mu A$ V ₁ = 2.7 V to 5.5 V	_	_	0.6	V	*1
Common driver drop voltage	V _{DC}	COM ₁ to COM ₄	$I_D = 2 \ \mu A$ V ₁ = 2.7 V to 5.5 V		_	0.3	V	*1
LCD power supply split-resistance	R _{LCD}		Between V_1 and V_{SS}	0.5	3.0	9.0	MΩ	
Liquid crystal display voltage	V _{LCD}	V ₁		2.2		5.5	V	*2

Notes: 1. The voltage drop from power supply pins V₁, V₂, V₃, and VSS to each segment pin or common pin.

2. When the liquid crystal display voltage is supplied from an external power source, ensure that the following relationship is maintained: $V_{CC} \ge V_1 \ge V_2 \ge V_3 \ge V_{SS}$.

14.3 Operation Timing

Figures 14.1 to 14.5 show timing diagrams.

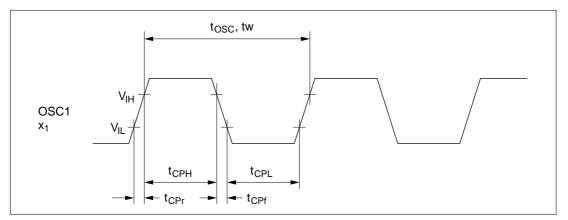


Figure 14.1 Clock Input Timing

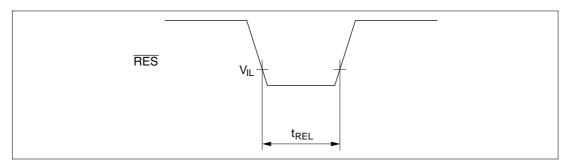


Figure 14.2 **RES** Low Width

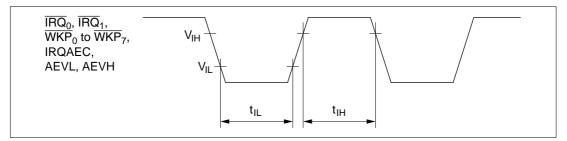
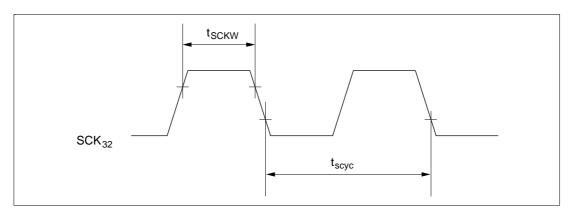
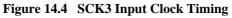


Figure 14.3 Input Timing





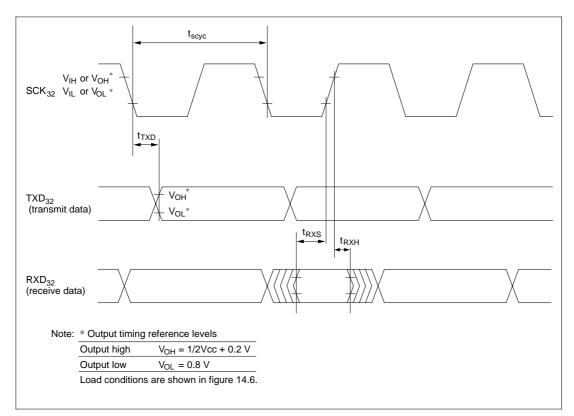


Figure 14.5 SCI3 Synchronous Mode Input/Output Timing

14.4 Output Load Circuit

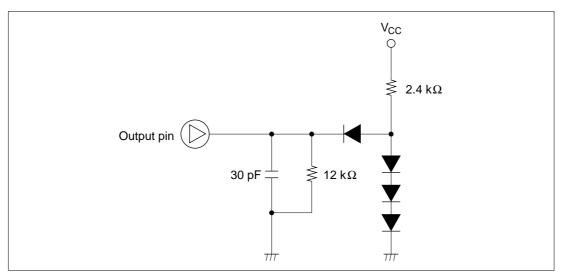


Figure 14.6 Output Load Condition

14.5 Resonator Equivalent Circuit

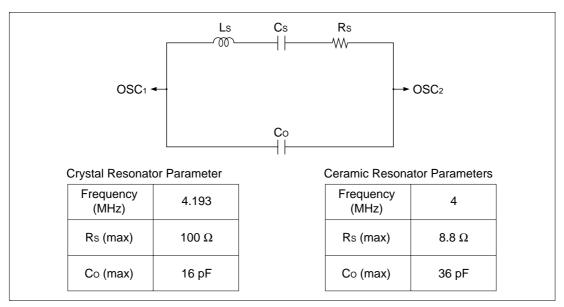


Figure 14.7 Resonator Equivalent Circuit

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14.6 Usage Note

The ZTAT and mask ROM versions both satisfy the electrical characteristics shown in this manual, but actual electrical characteristic values, operating margins, noise margins, and other properties may vary due to differences in manufacturing process, on-chip ROM, layout patterns, and so on.

When system evaluation testing is carried out using the ZTAT version, the same evaluation testing should also be conducted for the mask ROM version when changing over to that version.

Appendix A CPU Instruction Set

A.1 Instructions

Operation Notation

Rd8/16	General register (destination) (8 or 16 bits)
Rs8/16	General register (source) (8 or 16 bits)
Rn8/16	General register (8 or 16 bits)
CCR	Condition code register
N	N (negative) flag in CCR
Z	Z (zero) flag in CCR
V	V (overflow) flag in CCR
С	C (carry) flag in CCR
PC	Program counter
SP	Stack pointer
#xx: 3/8/16	Immediate data (3, 8, or 16 bits)
d: 8/16	Displacement (8 or 16 bits)
@aa: 8/16	Absolute address (8 or 16 bits)
+	Addition
-	Subtraction
×	Multiplication
÷	Division
^	Logical AND
V	Logical OR
\oplus	Exclusive logical OR
\rightarrow	Move
	Logical complement

Condition Code Notation

Symbol

\uparrow	Modified according to the instruction result
*	Not fixed (value not guaranteed)
0	Always cleared to 0
—	Not affected by the instruction execution result

Table A.1 lists the H8/300L CPU instruction set.

Table A.1Instruction Set

			In				sing Ler				s)							
	Operand Size		x: 8/16		@Rn	@(d:16, Rn)	@-Rn/@Rn+	@aa: 8/16	@(d:8, PC)	@ aa	mplied	С	one	ditio	on (Cod	le	of States
Mnemonic	õ	Operation	:xx#	R	0	0	ġ	0	0	0	<u></u>	I	н	Ν	z	v	С	°.
MOV.B #xx:8, Rd	В	$\text{\#xx:8} \rightarrow \text{Rd8}$	2									—		\updownarrow	\updownarrow	0	_	2
MOV.B Rs, Rd	В	$\text{Rs8} \rightarrow \text{Rd8}$		2								—		\updownarrow	\updownarrow	0	_	2
MOV.B @Rs, Rd	В	$@Rs16 \to Rd8$			2							—		\updownarrow	\updownarrow	0	_	4
MOV.B @(d:16, Rs), Rd	В	@(d:16, Rs16)→ Rd8				4						—		\updownarrow	\updownarrow	0	_	6
MOV.B @Rs+, Rd	в	$@$ Rs16 \rightarrow Rd8 Rs16+1 \rightarrow Rs16					2						—	\$	\$	0	-	6
MOV.B @aa:8, Rd	в	@aa:8 \rightarrow Rd8						2				—	—	\updownarrow	\uparrow	0	_	4
MOV.B @aa:16, Rd	в	@aa:16 \rightarrow Rd8						4				—	—	\$	\$	0	-	6
MOV.B Rs, @Rd	в	$Rs8 \rightarrow @Rd16$			2							—	—	\$	\$	0	-	4
MOV.B Rs, @(d:16, Rd)	в	$Rs8 \rightarrow @(d:16, Rd16)$				4						—	—	\$	\$	0	-	6
MOV.B Rs, @-Rd	в	Rd16–1 → Rd16 Rs8 → @Rd16					2							\$	\$	0	-	6
MOV.B Rs, @aa:8	в	$Rs8 \rightarrow @aa:8$						2				_	_	\$	\$	0	_	4
MOV.B Rs, @aa:16	в	$Rs8 \rightarrow @aa:16$						4				—		\$	\$	0	_	6
MOV.W #xx:16, Rd	W	#xx:16 → Rd	4									—		\$	\$	0	_	4
MOV.W Rs, Rd	W	$Rs16 \rightarrow Rd16$		2								—	—	\$	\$	0	—	2
MOV.W @Rs, Rd	W	$@Rs16 \to Rd16$			2							—		\updownarrow	\uparrow	0	_	4
MOV.W @(d:16, Rs), Rd	W	$@(\texttt{d:16}, \texttt{Rs16}) \rightarrow \texttt{Rd16}$				4						—	—	\updownarrow	\uparrow	0	_	6
MOV.W @Rs+, Rd	w	@Rs16 → Rd16 Rs16+2 → Rs16					2							\$	\$	0	-	6
MOV.W @aa:16, Rd	W	$@aa:16 \rightarrow Rd16 \\$						4				—		\updownarrow	\uparrow	0	_	6
MOV.W Rs, @Rd	W	$Rs16 \to @Rd16$			2							—		\updownarrow	\updownarrow	0	_	4
MOV.W Rs, @(d:16, Rd)	W	$\text{Rs16} \rightarrow @(\text{d:16}, \text{Rd16})$				4						—		\updownarrow	1	0	_	6
MOV.W Rs, @-Rd	W	$\begin{array}{l} \text{Rd162} \rightarrow \text{Rd16} \\ \text{Rs16} \rightarrow @\text{Rd16} \end{array}$					2							\$	\$	0	-	6
MOV.W Rs, @aa:16	W	$Rs16 \rightarrow @aa:16$						4				_	_	\updownarrow	\updownarrow	0		6
POP Rd	W	$\begin{array}{l} @SP \to Rd16 \\ SP+2 \to SP \end{array}$					2					_		\$	\$	0		6
PUSH Rs	W	$\begin{array}{l} SP-2 \rightarrow SP \\ Rs16 \rightarrow @SP \end{array}$					2					_		\$	\$	0		6

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			In					g M ngth			s)							
	Operand Size		k: 8/16		@Rn	@(d:16, Rn)	@-Rn/@Rn+	@aa: 8/16	@(d:8, PC)	@aa	Implied	с	ond	litio	on (Coc	le	. of States
Mnemonic	g	Operation	:XX#	Rn	0	0	9	0	0	0	<u></u>	I	н	Ν	z	v	С	Š
ADD.B #xx:8, Rd	В	$Rd8\text{+}\#xx:8\toRd8$	2									—	\updownarrow	\updownarrow	\updownarrow	\updownarrow	\updownarrow	2
ADD.B Rs, Rd	В	$Rd8\text{+}Rs8\toRd8$		2								—	\updownarrow	\updownarrow	\updownarrow	\updownarrow	\updownarrow	2
ADD.W Rs, Rd	W	$Rd16\text{+}Rs16 \rightarrow Rd16$		2								—	(1)	\updownarrow	\updownarrow	\updownarrow	\updownarrow	2
ADDX.B #xx:8, Rd	В	$Rd8\text{+}\#xx:8\ \text{+}C \to Rd8$	2									—	\updownarrow	\updownarrow	(2)	\updownarrow	\updownarrow	2
ADDX.B Rs, Rd	В	$Rd8\text{+}Rs8\text{+}C\rightarrowRd8$		2								—	\updownarrow	\updownarrow	(2)	\updownarrow	\updownarrow	2
ADDS.W #1, Rd	W	$Rd16+1 \rightarrow Rd16$		2								_	—	_	_		—	2
ADDS.W #2, Rd	w	$Rd16+2 \rightarrow Rd16$		2									—	_	—		—	2
INC.B Rd	В	$Rd8+1 \rightarrow Rd8$		2								—	—	\updownarrow	\updownarrow	\updownarrow	—	2
DAA.B Rd	В	Rd8 decimal adjust \rightarrow Rd8		2								—	*	\updownarrow	\updownarrow	*	(3)	2
SUB.B Rs, Rd	В	$Rd8Rs8\toRd8$		2								—	\updownarrow	\updownarrow	\updownarrow	\updownarrow	\updownarrow	2
SUB.W Rs, Rd	W	$Rd16Rs16 \rightarrow Rd16$		2								—	(1)	\updownarrow	\updownarrow	\updownarrow	\updownarrow	2
SUBX.B #xx:8, Rd	В	$Rd8\text{-}\#xx:8\text{-}C\toRd8$	2									—	\updownarrow	\updownarrow	(2)	\Leftrightarrow	\updownarrow	2
SUBX.B Rs, Rd	В	$Rd8\text{-}Rs8\text{-}C\toRd8$		2								—	\updownarrow	\updownarrow	(2)	\updownarrow	\updownarrow	2
SUBS.W #1, Rd	W	$Rd16-1 \rightarrow Rd16$		2								_	—	_	_	_	—	2
SUBS.W #2, Rd	w	$Rd16-2 \rightarrow Rd16$		2									—	_	—		—	2
DEC.B Rd	В	$Rd8-1 \rightarrow Rd8$		2								—	—	\updownarrow	€	↕	—	2
DAS.B Rd	в	Rd8 decimal adjust \rightarrow Rd8		2								—	*	\updownarrow	↕	*	—	2
NEG.B Rd	в	$0-Rd \rightarrow Rd$		2								—	\$	\updownarrow	\$	↕	\uparrow	2
CMP.B #xx:8, Rd	в	Rd8–#xx:8	2									_	\updownarrow	\updownarrow	↕	\updownarrow	\updownarrow	2
CMP.B Rs, Rd	В	Rd8–Rs8		2								—	\updownarrow	\updownarrow	\updownarrow	\updownarrow	\updownarrow	2
CMP.W Rs, Rd	W	Rd16–Rs16		2								—	(1)	\updownarrow	\updownarrow	\updownarrow	\updownarrow	2

			In					g M ngtł			s)							
	Operand Size		#xx: 8/16	_	@Rn	@(d:16, Rn)	@-Rn/@Rn+	@ aa: 8/16	@(d:8, PC)	@aa	Implied	С	one	ditio	on (Coc	le	of States
Mnemonic	õ	Operation	¥	Rn	8	0	0	0	0	0	<u>=</u>	I	н	Ν	z	v	С	Ś
MULXU.B Rs, Rd	В	Rd8 imes Rs8 ightarrow Rd16		2								_	_	_	_	_	_	14
DIVXU.B Rs, Rd	В	Rd16÷Rs8 → Rd16 (RdH: remainder, RdL: quotient)		2											(6)	_	_	14
AND.B #xx:8, Rd	В	$Rd8 {\scriptstyle \wedge} \#xx: 8 \rightarrow Rd8$	2									—	—	↕	\$	0	_	2
AND.B Rs, Rd	В	$Rd8 {\scriptscriptstyle \wedge} Rs8 \rightarrow Rd8$		2								—	—	↕	\$	0	_	2
OR.B #xx:8, Rd	В	$Rd8 \lor \#xx:8 \rightarrow Rd8$	2									_	_	↕	↕	0	_	2
OR.B Rs, Rd	В	$Rd8 \lor Rs8 \rightarrow Rd8$		2								_	_	\updownarrow	↕	0	_	2
XOR.B #xx:8, Rd	В	$Rd8 \oplus \#xx: 8 \rightarrow Rd8$	2									_		\updownarrow	↕	0	_	2
XOR.B Rs, Rd	в	$Rd8{\oplus}Rs8 \to Rd8$		2								_	—	↕	↕	0	_	2
NOT.B Rd	в	$\overline{Rd} \to Rd$		2								_		↕	\$	0	_	2
SHAL.B Rd	В			2								_		\$	\$	\$	\$	2
SHAR.B Rd	В			2										\$	\$	0	\$	2
SHLL.B Rd	В			2								_		\$	\$	0	\$	2
SHLR.B Rd	В	$0 \rightarrow \boxed[b_7 \\ b_0 \\ b_0 \\ c_1 \\ c_2 \\ c_3 \\ c_4 \\ c_5 \\ c_6 $		2										0	\$	0	\$	2
ROTXL.B Rd	В	C b7 b0		2										\$	\$	0	\$	2
ROTXR.B Rd	В	▶		2										\$	\$	0	\$	2

			In				sing Ler	•			s)							
	Operand Size		x: 8/16		@Rn	@(d:16, Rn)	@-Rn/@Rn+	@aa: 8/16	@(d:8, PC)	@ aa	mplied	с	one	ditio	on (Cod	le	. of States
Mnemonic	ő	Operation	:xx#	Rn	0	0	9	0	0	0	<u></u>	I	н	Ν	z	v	С	No.
ROTL.B Rd	В			2										\$	\$	0	\$	2
ROTR.B Rd	В	▶		2										\$	\$	0	\$	2
BSET #xx:3, Rd	В	(#xx:3 of Rd8) ← 1		2								—	—	—	_	_	—	2
BSET #xx:3, @Rd	в	(#xx:3 of @Rd16) ← 1			4							_	_	_	_	_	_	8
BSET #xx:3, @aa:8	в	(#xx:3 of @aa:8) ← 1						4				_	_	_	_	_	_	8
BSET Rn, Rd	в	(Rn8 of Rd8) ← 1		2								_		_	_	_	_	2
BSET Rn, @Rd	в	(Rn8 of @Rd16) ← 1			4							—		_	_	_	_	8
BSET Rn, @aa:8	В	(Rn8 of @aa:8) ← 1						4				—	—	_	_	_	_	8
BCLR #xx:3, Rd	В	(#xx:3 of Rd8) ← 0		2								—	—	—	—	—	—	2
BCLR #xx:3, @Rd	В	(#xx:3 of @Rd16) ← 0			4							—	—	—	_	—	—	8
BCLR #xx:3, @aa:8	В	(#xx:3 of @aa:8) ← 0						4				—	—	—	—	—	—	8
BCLR Rn, Rd	В	(Rn8 of Rd8) ← 0		2								—		—	_	—	—	2
BCLR Rn, @Rd	В	(Rn8 of @Rd16) ← 0			4							—		—	_	—	—	8
BCLR Rn, @aa:8	В	(Rn8 of @aa:8) ← 0						4				—	—	—	_	_	—	8
BNOT #xx:3, Rd	В	(#xx:3 of Rd8) ← (#xx:3 of Rd8)		2												_	_	2
BNOT #xx:3, @Rd	В	(#xx:3 of @Rd16) ← (#xx:3 of @Rd16)			4							_		_	_	_	—	8
BNOT #xx:3, @aa:8	в	(#xx:3 of @aa:8) ← (#xx:3 of @aa:8)						4				_	_	—	_	_	_	8
BNOT Rn, Rd	в	(Rn8 of Rd8) ← (Rn8 of Rd8)		2								_				—	—	2
BNOT Rn, @Rd	в	(Rn8 of @Rd16) ← (Rn8 of @Rd16)			4							—		—	—	—	—	8
BNOT Rn, @aa:8	в	(Rn8 of @aa:8) ← (Rn8 of @aa:8)						4				—					_	8

			In		dd Joti						s)							
	Operand Size		#xx: 8/16	Rn	@Rn	@(d:16, Rn)	@-Rn/@Rn+	@aa: 8/16	@(d:8, PC)	@ @aa	Implied		on					No. of States
		Operation	#	<u> </u>		•	•		•	9	-	I	н	Ν	Z ↑	V	С	
BTST #xx:3, Rd	B	$(\#xx:3 \text{ of } Rd8) \rightarrow Z$		2								_		_	\$	_		2
BTST #xx:3, @Rd	B	$(\#xx:3 \text{ of } @ \text{Rd16}) \rightarrow Z$			4							_		_	\$	_		6
BTST #xx:3, @aa:8	В	$(\#xx:3 \text{ of } @aa:8) \rightarrow Z$						4				_		_	\$	-		6
BTST Rn, Rd	В	$(Rn8 of Rd8) \rightarrow Z$		2								_		_	\$	_		2
BTST Rn, @Rd	В	(Rn8 of @Rd16) \rightarrow Z			4							_	—	_	\$	-		6
BTST Rn, @aa:8	В	(Rn8 of @aa:8) \rightarrow Z						4				_		_	\$	_		6
BLD #xx:3, Rd	В	(#xx:3 of Rd8) \rightarrow C		2								_	—	_	—	—	↕	2
BLD #xx:3, @Rd	В	(#xx:3 of @Rd16) \rightarrow C			4							—	—	—	—	_	\uparrow	6
BLD #xx:3, @aa:8	В	(#xx:3 of @aa:8) \rightarrow C						4				—	—	—	_	_	\uparrow	6
BILD #xx:3, Rd	В	$(\overline{\#xx:3 \text{ of } Rd8}) \to C$		2								—	—	—	_	_	$\left \uparrow \right $	2
BILD #xx:3, @Rd	В	$(\overline{\#xx:3 \text{ of } @ Rd16}) \to C$			4							—	—	—	_	_	\uparrow	6
BILD #xx:3, @aa:8	В	$(\overline{\#xx:3 \text{ of } @aa:8}) \rightarrow C$						4				—	—	_	—	—	\uparrow	6
BST #xx:3, Rd	в	$C \rightarrow$ (#xx:3 of Rd8)		2								—	_	—	—		—	2
BST #xx:3, @Rd	В	$C \rightarrow (\#xx:3 \text{ of } @Rd16)$			4							—	—	_	—	—	$\left -\right $	8
BST #xx:3, @aa:8	В	$C \rightarrow (\#xx:3 \text{ of } @aa:8)$						4				—	—	_	—	—	$\left - \right $	8
BIST #xx:3, Rd	в	$\overline{C} \rightarrow$ (#xx:3 of Rd8)		2								_		_	_	_	-	2
BIST #xx:3, @Rd	в	$\overline{C} \rightarrow (\#xx:3 \text{ of } @Rd16)$			4							_		_	_	_	-	8
BIST #xx:3, @aa:8	в	$\overline{C} \rightarrow$ (#xx:3 of @aa:8)						4				_		_	_	_	-	8
BAND #xx:3, Rd	в	$C_{\wedge}(\#xx:3 \text{ of } Rd8) \rightarrow C$		2								_	_	_	_	_	\$	2
BAND #xx:3, @Rd	в	$C_{\wedge}(\#xx:3 \text{ of } @Rd16) \rightarrow C$			4							_	_	_	_	_	\uparrow	6
BAND #xx:3, @aa:8	в	$C_{\wedge}(\#xx:3 \text{ of } @aa:8) \rightarrow C$						4				_	_	_	_	_	\uparrow	6
BIAND #xx:3, Rd	в	$C \land (\overline{\#xx:3 \text{ of } Rd8}) \rightarrow C$		2										_		_	\$	2
BIAND #xx:3, @Rd	в	$C_{\wedge}(\overline{\#xx:3 \text{ of } @ \text{Rd16}}) \rightarrow C$			4							_	_	_	_	_	\uparrow	6
BIAND #xx:3, @aa:8	в	$C_{\wedge}(\overline{\#xx:3 \text{ of } @aa:8}) \rightarrow C$						4				_	_	_	_	_	\uparrow	6
BOR #xx:3, Rd	В	C∨(#xx:3 of Rd8) → C		2								_	_	_	_	_	\uparrow	2
BOR #xx:3, @Rd	в	$C_{\vee}(\#xx:3 \text{ of } @Rd16) \rightarrow C$			4							_	_	_	_	_	\$	6
BOR #xx:3, @aa:8	в	$C \lor (\#xx:3 \text{ of } @aa:8) \rightarrow C$						4				_	_	_	_	_	\$	6
BIOR #xx:3, Rd	в	$C \lor (\overline{\#xx:3 \text{ of } Rd8}) \rightarrow C$		2								_	_	_	_	_	\$	2
BIOR #xx:3, @Rd	в	$C \lor (\#xx:3 \text{ of } @ Rd16) \rightarrow C$			4							_	_	_	_	_	\$	6

				In						ode 1 (b		s)							
	Operand Size		Branching	c: 8/16		٨n	@(d:16, Rn)	@-Rn/@Rn+	@ aa: 8/16	@(d:8, PC)	@aa	mplied	С	one	ditie	on (Coc	de	. of States
Mnemonic	do	Operation	Condition	:xx#	R	@Rn	0	8	8	0	0	Ē	I	н	Ν	z	v	С	No.
BIOR #xx:3, @aa:8	В	C∨(#xx:3 of	$eaa:8) \rightarrow C$						4				—	—	_	—	—	\$	6
BXOR #xx:3, Rd	В	C⊕(#xx:3 o	f Rd8) \rightarrow C		2								_		_	_	_	\$	2
BXOR #xx:3, @Rd	в	C⊕(#xx:3 o	f @Rd16) \rightarrow C			4							_		_	_	_	\$	6
BXOR #xx:3, @aa:8	В	C⊕(#xx:3 o	f @aa:8) \rightarrow C						4				_	_	_	_	_	\$	6
BIXOR #xx:3, Rd	В	C⊕(#xx:3 o	$\overline{f Rd8}) \rightarrow C$		2								_	_	_	_	_	\$	2
BIXOR #xx:3, @Rd	В	C⊕(#xx:3 o	f @Rd16) \rightarrow C			4							_	_	_	_	_	\$	6
BIXOR #xx:3, @aa:8	В	C⊕(#xx:3 o	f @aa:8) \rightarrow C						4				_	_	_	_	_	\$	6
BRA d:8 (BT d:8)	_	$PC \leftarrow PC + c$	d:8							2			_	_	_	_	_	_	4
BRN d:8 (BF d:8)	_	$PC \leftarrow PC + 2$	2							2			_	_	_	_	_	_	4
BHI d:8	_	lf	$C \lor Z = 0$							2			_	_		_	_	_	4
BLS d:8	_	condition	C ∨ Z = 1							2			_		_	_	_	_	4
BCC d:8 (BHS d:8)	_	is true then	C = 0							2			_		_	_	_	_	4
BCS d:8 (BLO d:8)	_	$PC \leftarrow$	C = 1							2			_			_	_	_	4
BNE d:8	_	PC+d:8 else next:	Z = 0							2			_			_	_	_	4
BEQ d:8	_		Z = 1							2			_			_	_	_	4
BVC d:8	_		V = 0							2			_			_	_	_	4
BVS d:8	_		V = 1							2			_			_	_	_	4
BPL d:8	_		N = 0							2			_			_		_	4
BMI d:8	_		N = 1							2			_	_		_	_	_	4
BGE d:8	_		N⊕V = 0							2			_	_	_	_	_	_	4
BLT d:8	_		N⊕V = 1							2			_		_	_	_	_	4
BGT d:8	_		Z ∨ (N⊕V) = 0							2			_			_	_	_	4
BLE d:8	_		Z ∨ (N⊕V) = 1							2			_			_	_	_	4
JMP @Rn	_	$PC \leftarrow Rn16$	3			2							_	_	_	_	_	_	4
JMP @aa:16	_	$PC \leftarrow aa:16$	6						4				_	_	_	_	—	_	6
JMP @@aa:8		PC ← @aa	:8								2		_		_	_	_	_	8
BSR d:8	_	$\begin{array}{c} SP-2 \rightarrow SF \\ PC \rightarrow @SF \\ PC \leftarrow PC+c \end{array}$)							2								_	6

			In				sin Lei				s)							
Mnemonic	Operand Size	Operation	#xx: 8/16	Rn	@Rn	@(d:16, Rn)	@-Rn/@Rn+	@aa: 8/16	@(d:8, PC)	@ @aa	Implied	C	on	ditio N	on (Coc	le C	No. of States
JSR @Rn	-	$\begin{array}{l} SP-2 \rightarrow SP \\ PC \rightarrow @SP \\ PC \leftarrow Rn16 \end{array}$			2								_		_	_		6
JSR @aa:16	-	$\begin{array}{l} SP-2 \to SP \\ PC \to @ SP \\ PC \leftarrow aa: 16 \end{array}$						4								_	-	8
JSR @@aa:8	_	$\begin{array}{c} SP-2 \to SP \\ PC \to @ SP \\ PC \leftarrow @ aa:8 \end{array}$								2							_	8
RTS	-	$PC \leftarrow @SP$ $SP+2 \rightarrow SP$									2		_		—	—		8
RTE	-	$CCR \leftarrow @SP$ SP+2 \rightarrow SP PC $\leftarrow @SP$ SP+2 \rightarrow SP									2	\$	\$	\$	\$	\$	\$	10
SLEEP	-	Transit to sleep mode.									2		_	—	_	_	—	2
LDC #xx:8, CCR	В	$\#xx:8 \rightarrow CCR$	2									\updownarrow	\$	\$	\$	€	\updownarrow	2
LDC Rs, CCR	В	$Rs8 \rightarrow CCR$		2								\updownarrow	\$	\$	\$	\$	\updownarrow	2
STC CCR, Rd	В	$CCR \rightarrow Rd8$		2								—	—	—	—	—	—	2
ANDC #xx:8, CCR	В	CCR_{A} #xx:8 $\rightarrow CCR$	2									\uparrow	\$	\$	\$	\$	\$	2
ORC #xx:8, CCR	В	$CCR \lor \#xx:8 \rightarrow CCR$	2									\uparrow	\$	\$	\$	↕	\$	2
XORC #xx:8, CCR	В	$CCR \oplus \#xx:8 \rightarrow CCR$	2									\updownarrow	\$	\$	\$	€	\updownarrow	2
NOP	_	$PC \gets PC+2$									2		—	—			—	2
EEPMOV		if R4L \pm 0 Repeat @R5 \rightarrow @R6 R5+1 \rightarrow R5 R6+1 \rightarrow R6 R4L-1 \rightarrow R4L Until R4L=0 else next;									4							(4)

Notes: (1) Set to 1 when there is a carry or borrow from bit 11; otherwise cleared to 0.

(2) If the result is zero, the previous value of the flag is retained; otherwise the flag is cleared to 0.

(3) Set to 1 if decimal adjustment produces a carry; otherwise retains value prior to arithmetic operation.

(4) The number of states required for execution is 4n + 9 (n = value of R4L).

(5) Set to 1 if the divisor is negative; otherwise cleared to 0.

(6) Set to 1 if the divisor is zero; otherwise cleared to 0.

A.2 Operation Code Map

Table A.2 is an operation code map. It shows the operation codes contained in the first byte of the instruction code (bits 15 to 8 of the first instruction word).

Instruction when first bit of byte 2 (bit 7 of first instruction word) is 0. Instruction when first bit of byte 2 (bit 7 of first instruction word) is 1.

High	0	-	2	3	4	5	9	7	8	6	А	В	С	D	Е	ш
0	NOP	SHLL			BRA	MULXU		BSEI								
-	SLEEP	SHLR F			BRN	DIVXU		BNOI								
7	STC	ROTXL			BHI			BCLK								
3	LDC	ROTXR			BLS		BTST BO									
4	ORC	OR			BCC	RTS		BOR BIOR								
5	XORC	XOR			BCS	BSR		BXOR BIXOR								
9	ANDC	AND			BNE	RTE		BAND BIAND								
7	LDC	NOT NEG			BEQ		BST BIST	BLD	AL	AD	C	SU	0	X	AN	W
8	ADD	SUB		2	BVC				ADD	ADDX	CMP	SUBX	OR	XOR	AND	MOV
0	Q	۵			BVS			NOM								
٨	INC	DEC			BPL	dWL	WO									
В	ADDS	SUBS			BMI			EEPMOV								
U	MOV	Ö			BGE		* NOM									
۵	٨C	CMP			ВLT	-		-manipulatic								
ш	ADDX	SUBX			BGT	JSR		Bit-manipulation instructions								
Ŀ	DAA	DAS			BLE											

Note: * The PUSH and POP instructions are identical in machine language to MOV instructions.

Table A.2 Operation Code Map

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A.3 Number of Execution States

The tables here can be used to calculate the number of states required for instruction execution. Table A.4 indicates the number of states required for each cycle (instruction fetch, read/write, etc.), and table A.3 indicates the number of cycles of each type occurring in each instruction. The total number of states required for execution of an instruction can be calculated from these two tables as follows:

 $Execution \ states = I \times S_I + J \times S_J + K \times S_K + L \times S_L + M \times S_M + N \times S_N$

Examples: When instruction is fetched from on-chip ROM, and an on-chip RAM is accessed.

BSET #0, @FF00 From table A.4: I = L = 2, J = K = M = N = 0From table A.3: $S_I = 2$, $S_L = 2$ Number of states required for execution $= 2 \times 2 + 2 \times 2 = 8$ When instruction is fetched from on-chip ROM, branch address is read from on-chip ROM, and on-chip RAM is used for stack area.

JSR @@ 30 From table A.4: I = 2, J = K = 1, L = M = N = 0From table A.3: $S_I = S_J = S_K = 2$ Number of states required for execution $= 2 \times 2 + 1 \times 2 + 1 \times 2 = 8$

Execution Status			Access Location					
(instruction cycle)		On-Chip Memory	On-Chip Peripheral Module					
Instruction fetch	SI	2	_					
Branch address read	SJ							
Stack operation	S _K							
Byte data access	SL		2 or 3*					
Word data access	S _M		_					
Internal operation	S _N	1						

Table A.3 Number of Cycles in Each Instruction

Note: * Depends on which on-chip module is accessed. See 2.9.1, Notes on Data Access for details.

Instruction	Mnemonic	Instruction Fetch I	Branch Addr. Read J	Stack Operation K	Byte Data Access L	Word Data Access M	Internal Operation N
			J	ĸ	L	IVI	N
ADD	ADD.B #xx:8, Rd	1					
	ADD.B Rs, Rd	1					
4000	ADD.W Rs, Rd	1					
ADDS	ADDS.W #1, Rd	1					
	ADDS.W #2, Rd	1					
ADDX	ADDX.B #xx:8, Rd	1					
	ADDX.B Rs, Rd	1					
AND	AND.B #xx:8, Rd	1					
	AND.B Rs, Rd	1					
ANDC	ANDC #xx:8, CCR	1					
BAND	BAND #xx:3, Rd	1					
	BAND #xx:3, @Rd	2			1		
	BAND #xx:3, @aa:8	2			1		
Bcc	BRA d:8 (BT d:8)	2			0		
	BRN d:8 (BF d:8)	2					
	BHI d:8	2					
	BLS d:8	2					
	BCC d:8 (BHS d:8)	2					
	BCS d:8 (BLO d:8)	2					
	BNE d:8	2					
	BEQ d:8	2					
	BVC d:8	2					
	BVS d:8	2					
	BPL d:8	2					
	BMI d:8	2					
	BGE d:8	2					
	BLT d:8	2					
	BGT d:8	2					
	BLE d:8	2					
BCLR	BCLR #xx:3, Rd	1			1		
	BCLR #xx:3, @Rd	2			2		
	BCLR #xx:3, @aa:8	2			2		
	BCLR Rn, Rd	1					
	BCLR Rn, @Rd	2			2		
	BCLR Rn, @aa:8	2			2		
BIAND	BIAND #xx:3, Rd	1					
	BIAND #xx:3, @Rd	2			1		
	BIAND #xx:3, @aa:8	2			1		

Table A.4 Number of Cycles in Each Instruction

Instruction	Mnemonic	Instruction Fetch I	Branch Addr. Read J	Stack Operation K	Byte Data Access L	Word Data Access M	Internal Operation N
BILD	BILD #xx:3, Rd	1	5	ĸ	-	IVI	N
DILU	BILD #xx:3, @Rd	2			1		
	BILD #xx:3, @Ru BILD #xx:3, @aa:8	2			1		
					1		
BIOR	BIOR #xx:3, Rd	1					
	BIOR #xx:3, @Rd	2			1		
	BIOR #xx:3, @aa:8	2			1		
BIST	BIST #xx:3, Rd	1					
	BIST #xx:3, @Rd	2			2		
	BIST #xx:3, @aa:8	2			2		
BIXOR	BIXOR #xx:3, Rd	1					
	BIXOR #xx:3, @Rd	2			1		
	BIXOR #xx:3, @aa:8	2			1		
BLD	BLD #xx:3, Rd	1					
	BLD #xx:3, @Rd	2			1		
	BLD #xx:3, @aa:8	2			1		
BNOT	BNOT #xx:3, Rd	1					
	BNOT #xx:3, @Rd	2			2		
	BNOT #xx:3, @aa:8	2			2		
	BNOT Rn, Rd	1					
	BNOT Rn, @Rd	2			2		
	BNOT Rn, @aa:8	2			2		
BOR	BOR #xx:3, Rd	1					
	BOR #xx:3, @Rd	2			1		
	BOR #xx:3, @aa:8	2			1		
BSET	BSET #xx:3, Rd	1					
	BSET #xx:3, @Rd	2			2		
	BSET #xx:3, @aa:8	2			2		
	BSET Rn, Rd	1					
	BSET Rn, @Rd	2			2		
	BSET Rn, @aa:8	2			2		
BSR	BSR d:8	2		1	.1		
BST	BST #xx:3, Rd	1					
DOT	BST #xx:3, @Rd	2			2		
	BST #xx:3, @aa:8	2			2		
BTST	BTST #xx:3, Rd	1			<u>د</u>		
1010		2			1		
	BTST #xx:3, @Rd	2					
	BTST #xx:3, @aa:8	2			1		
	BTST Rn, Rd				1		
	BTST Rn, @Rd	2			1		

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Instruction	Mnemonic	Instruction Fetch I	Branch Addr. Read J	Stack Operation K	Byte Data Access L	Word Data Access M	Internal Operation N
BTST	BTST Rn, @aa:8	2	•	i.	1	m	n
BXOR	BXOR #xx:3, Rd			-	· · ·		
BROR	BXOR #xx:3, @Rd	2			1		
	BXOR #xx:3, @aa:8	2			1		
CMP	CMP. B #xx:8, Rd	1					
0	CMP. B Rs, Rd	1					
	CMP.W Rs, Rd	1					
DAA	DAA.B Rd	1					
DAS	DAS.B Rd	1					
DEC	DEC.B Rd	1					
DIVXU	DIVXU.B Rs, Rd	1					12
EEPMOV		2			2n+2*		12
	EEPMOV				211+2*		
INC	INC.B Rd	1					
JMP	JMP @Rn	2					
	JMP @aa:16	2	4				2
100	JMP @@aa:8	2	1				2
JSR	JSR @Rn	2		1			
	JSR @aa:16	2	4	1			2
	JSR @@aa:8	2	1	1			
LDC	LDC #xx:8, CCR	1					
	LDC Rs, CCR	1					
MOV	MOV.B #xx:8, Rd	1					
	MOV.B Rs, Rd	1					
	MOV.B @Rs, Rd	1			1		
	MOV.B @(d:16, Rs), Rd	2			1		
	MOV.B @Rs+, Rd	1			1		2
	MOV.B @aa:8, Rd	1			1		
	MOV.B @aa:16, Rd	2			1		
	MOV.B Rs, @Rd	1			1		
	MOV.B Rs, @(d:16, Rd)	2			1		0
	MOV.B Rs, @–Rd MOV.B Rs, @aa:8	1 1			1 1		2
	MOV.B Rs, @aa:16	1 2			1		
	MOV.B RS, @aa.16 MOV.W #xx:16, Rd	2			I		
	MOV.W #xx.16, Rd	2					
	MOV.W @Rs, Rd	1				1	
	MOV.W @(d:16, Rs), Rd	2				1	
	MOV.W @(d.10, 1(3), 1(d) MOV.W @Rs+, Rd	1				1	2
	MOV.W @aa:16, Rd	2				1	-

Note: * n: Initial value in R4L. The source and destination operands are accessed n + 1 times each.

Instruction	Mnemonic	Instruction Fetch I	Branch Addr. Read J	Stack Operation K	Byte Data Access L	Word Data Access M	Internal Operation N
MOV	MOV.W Rs, @Rd	1	•	N	-	1	N
	MOV.W Rs, @(d:16, Rd)	2				1	
	MOV.W Rs, @-Rd	1				1	2
	MOV.W Rs, @aa:16	2				1	
MULXU	MULXU.B Rs, Rd	1					12
NEG	NEG.B Rd	1					
NOP	NOP	1					
NOT	NOT.B Rd	1					
OR	OR.B #xx:8, Rd	1					
	OR.B Rs, Rd	1					
ORC	ORC #xx:8, CCR	1					
ROTL	ROTL.B Rd	1					
ROTR	ROTR.B Rd	1					
ROTXL	ROTXL.B Rd	1					
ROTXR	ROTXR.B Rd	1					
RTE	RTE	2		2		α.	2
RTS	RTS	2		1			2
SHAL	SHAL.B Rd	1					
SHAR	SHAR.B Rd	1					
SHLL	SHLL.B Rd	1					
SHLR	SHLR.B Rd	1	1				
SLEEP	SLEEP	1					
STC	STC CCR, Rd	1					
SUB	SUB.B Rs, Rd	1	1				
	SUB.W Rs, Rd	1					
SUBS	SUBS.W #1, Rd	1	0		AI.	Α	0
	SUBS.W #2, Rd	1					
POP	POP Rd	1		1			2
PUSH	PUSH Rs	1		1			2
SUBX	SUBX.B #xx:8, Rd	1				1	
	SUBX.B Rs, Rd	1					
XOR	XOR.B #xx:8, Rd	1					
	XOR.B Rs, Rd	1					
XORC	XORC #xx:8, CCR	1					

Appendix B Internal I/O Registers

B.1 Addresses

Lower	Register				Bit N	lames				Module
Address	Name	Bit 7	Bit 6	Bit 5	Bit 4	Bit 3	Bit 2	Bit 1	Bit 0	Name
H'80										
H'81										
H'82										-
H'83										-
H'84										-
H'85										-
H'86										-
H'87										-
H'88										-
H'89										-
H'8A										-
H'8B										-
H'8C	ECPWCRH	ECPWCRH7	ECPWCRH6	ECPWCRH5	ECPWCRH4	ECPWCRH3	ECPWCRH2	ECPWCRH1	ECPWCRH0	Asynchronous
H'8D	ECPWCRL	ECPWCRL7	ECPWCRL6	ECPWCRL5	ECPWCRL4	ECPWCRL3	ECPWCRL2	ECPWCRL1	ECPWCRL0	event counter
H'8E		ECPWDRH7								-
H'8F	ECPWDRL	ECPWDRL7	ECPWDRL6	ECPWDRL5	ECPWDRL4	ECPWDRL3	ECPWDRL2	ECPWDRL1	ECPWDRL0	-
H'90	WEGR	WKEGS7	WKEGS6	WKEGS5	WKEGS4	WKEGS3	WKEGS2	WKEGS1	WKEGS0	System control
H'91	SPCR	_	_	SPC32	_	SCINV3	SCINV2	_	_	SCI
H'92	AEGSR	AHEGS1	AHEGS0	ALEGS1	ALEGS0	AIEGS1	AIEGS0	ECPWME	_	Asynchronous
H'93										event counter
H'94	ECCR	ACKH1	ACKH0	ACKL1	ACKL0	PWCK2	PWCK1	PWCK0	_	-
H'95	ECCSR	OVH	OVL	_	CH2	CUEH	CUEL	CRCH	CRCL	-
H'96	ECH	ECH7	ECH6	ECH5	ECH4	ECH3	ECH2	ECH1	ECH0	-
H'97	ECL	ECL7	ECL6	ECL5	ECL4	ECL3	ECL2	ECL1	ECL0	-
H'98										
H'99										-
H'9A										-
H'9B										-
H'9C										-
H'9D										-
H'9E										-
H'9F										-
H'A0										-
H'A1										-

Lower	Register				Bit	Names				Module
Address	Name	Bit 7	Bit 6	Bit 5	Bit 4	Bit 3	Bit 2	Bit 1	Bit 0	Name
H'A2										
H'A3										
H'A4										
H'A5										
H'A6										
H'A7										
H'A8	SMR	COM	CHR	PE	PM	STOP	MP	CKS1	CKS0	SCI
H'A9	BRR	BRR7	BRR6	BRR5	BRR4	BRR3	BRR2	BRR1	BRR0	
H'AA	SCR3	TIE	RIE	TE	RE	MPIE	TEIE	CKE1	CKE0	
H'AB	TDR	TDR7	TDR6	TDR5	TDR4	TDR3	TDR2	TDR1	TDR0	_
H'AC	SSR	TDRE	RDRF	OER	FER	PER	TEND	MPBR	MPBT	_
H"AD	RDR	RDR7	RDR6	RDR5	RDR4	RDR3	RDR2	RDR1	RDR0	_
H'AE										_
H'AF										_
H'B0	TMA	_	_	_	_	TMA3	TMA2	TMA1	TMA0	Timer A
H'B1	ТСА	TCA7	TCA6	TCA5	TCA4	TCA3	TCA2	TCA1	TCA0	_
H'B2										
H'B3	-0									
H'B4										
H'B5	1									
H'B6	TCRF	TOLH	CKSH2	CKSH1	CKSH0	TOLL	CKSL2	CKSL1	CKSL0	Timer F
H'B7	TCSRF	OVFH	CMFH	OVIEH	CCLRH	OVFL	CMFL	OVIEL	CCLRL	
H'B8	TCFH	TCFH7	TCFH6	TCFH5	TCFH4	TCFH3	TCFH2	TCFH1	TCFH0	
H'B9	TCFL	TCFL7	TCFL6	TCFL5	TCFL4	TCFL3	TCFL2	TCFL1	TCFL0	
H'BA	OCRFH	OCRFH7	OCRFH6	OCRFH5	OCRFH4	OCRFH3	OCRFH2	OCRFH1	OCRFH0	
H'BB	OCRFL	OCRFL7	OCRFL6	OCRFL5	OCRFL4	OCRFL3	OCRFL2	OCRFL1	OCRFL0	
H'BC	0	-14					-14	-14		
H'BD										
H'BE	0									
H'BF										
H'C0	LPCR	DTS1	DTS0	СМХ	_	SGS3	SGS2	SGS1	SGS0	LCD controller/
H'C1	LCR	_	PSW	ACT	DISP	CKS3	CKS2	CKS1	CKS0	driver
H'C2	LCR2	LCDAB	_	_	_	_	_	_	_	_
H'C3										
H'C4	ADRRH	ADR9	ADR8	ADR7	ADR6	ADR5	ADR4	ADR3	ADR2	A/D converter
H'C5	ADRRL	ADR1	ADR0	_	_	_	_	_	_	_
H'C6	AMR	CKS	_	_	_	СНЗ	CH2	CH1	CH0	_
H'C7	ADSR	ADSF	_	_	_	_	_	_	_	_
H'C8										I/O port
H'C9	PMR2		_	POF1		_				

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Lower	Register	erBit Names								
Address	-	Bit 7	Bit 6	Bit 5	Bit 4	Bit 3	Bit 2	Bit 1	Bit 0	_Module Name
H'CA	PMR3	AEVL	AEVH	_	_	_	TMOFH	TMOFL	_	I/O port
H'CB										_
H'CC	PMR5	WKP7	WKP6	WKP5	WKP4	WKP3	WKP2	WKP1	WKP0	_
H'CD	PWCR2	—	_	_	_	_	_	PWCR21	PWCR20	10 bit PWM2
H'CE	PWDRU2	_	_	_	_	_	_	PWDRU21	PWDRU20	_
H'CF	PWDRL2	PWDRL27	PWDRL26	PWDRL25	PWDRL24	PWDRL23	PWDRL22	PWDRL21	PWDRL20	_
H'D0	PWCR1	_	_	_	_	_	_	PWCR11	PWCR10	10 bit PWM1
H'D1	PWDRU1	_	_	_	_	_	_	PWDRU11	PWDRU10	_
H'D2	PWDRL1	PWDRL17	PWDRL16	PWDRL15	PWDRL14	PWDRL13	PWDRL12	PWDRL11	PWDRL10	_
H'D3										I/O port
H'D4										
H'D5										_
H'D6	PDR3	P37	P36	P35	P34	P33	P32	P31	_	-
H'D7	PDR4	_	_	_	_	P43	P42	P41	P40	_
H'D8	PDR5	P57	P56	P55	P54	P53	P52	P51	P50	_
H'D9	PDR6	P67	P66	P65	P64	P63	P62	P61	P60	_
H'DA	PDR7	P77	P76	P75	P74	P73	P72	P71	P70	_
H'DB	PDR8	_	_	_	_	_	_	_	P80	-
H'DC	PDR9	_	_	P95	P94	P93	P92	P91	P90	-
H'DD	PDRA	_	_	_	_	PA3	PA2	PA1	PA0	-
H'DE	PDRB	_				PB3	PB2	PB1	PB0	-
H'DF										_
H'E0										_
H'E1	PUCR3	PUCR37	PUCR36	PUCR35	PUCR34	PUCR33	PUCR32	PUCR31	_	_
H'E2	PUCR5	PUCR57	PUCR56	PUCR55	PUCR54	PUCR53	PUCR52	PUCR51	PUCR50	_
H'E3	PUCR6	PUCR67	PUCR66	PUCR65	PUCR64	PUCR63	PUCR62	PUCR61	PUCR60	_
H'E4										-
H'E5										-
H'E6	PCR3	PCR37	PCR36	PCR35	PCR34	PCR33	PCR32	PCR31	_	_
H'E7	PCR4	_	_	_	_	_	PCR42	PCR41	PCR40	_
H'E8	PCR5	PCR57	PCR56	PCR55	PCR54	PCR53	PCR52	PCR51	PCR50	_
H'E9	PCR6	PCR67	PCR66	PCR65	PCR64	PCR63	PCR62	PCR61	PCR60	_
H'EA	PCR7	PCR77	PCR76	PCR75	PCR74	PCR73	PCR72	PCR71	PCR70	-
H'EB	PCR8		_	_	_	_	_	_	PCR80	_
H'EC	PMR9	_	_	_	_	PIOFF	_	PWM2	PWM1	-
H'ED	PCRA		_	_	_	PCRA3	PCRA2	PCRA1	PCRA0	-
H'EE	PMRB	_	_	_	_	IRQ1	_	_	_	_
H'EF										_
H'F0	SYSCR1	SSBY	STS2	STS1	STS0	LSON	_	MA1	MA0	System contro
H'F1	SYSCR2		_	_	NESEL	DTON	MSON	SA1	SA0	-

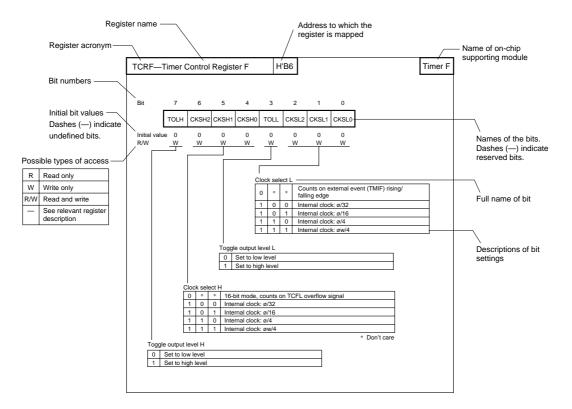
Lower	Register				Bit N	Names				Module Name
Address	Name	Bit 7	Bit 6	Bit 5	Bit 4	Bit 3	Bit 2	Bit 1	Bit 0	
H'F2	IEGR	_	_	_	_	_	_	IEG1	IEG0	System control
H'F3	IENR1	IENTA	_	IENWP	_	_	IENEC2	IEN1	IEN0	
H'F4	IENR2	IENDT	IENAD	_	_	IENTFH	IENTFL	_	IENEC	_
H'F5										_
H'F6	IRR1	IRRTA	_	_	_	_	IRREC2	IRRI1	IRRI0	
H'F7	IRR2	IRRDT	IRRAD	_	_	IRRTFH	IRRTFL	_	IRREC	
H'F8										
H'F9	IWPR	IWPF7	IWPF6	IWPF5	IWPF4	IWPF3	IWPF2	IWPF1	IWPF0	System control
H'FA	CKSTPR1	_	_	S32CKSTP	ADCKSTP	_	TFCKSTP	_	TACKSTP	_
H'FB	CKSTPR2	_	_	_	PW2CKSTP	AECKSTP	_	PW1CKSTP	LDCKSTP	_
H'FC										
H'FD										_
H'FE										_
H'FF										_
اممممط										

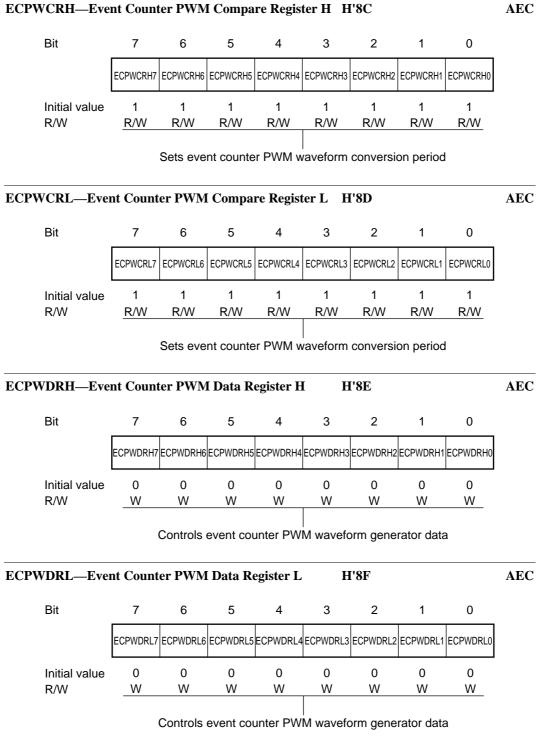
Legend

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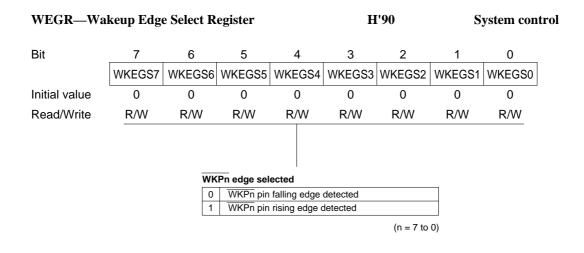
SCI: Serial Communication Interface

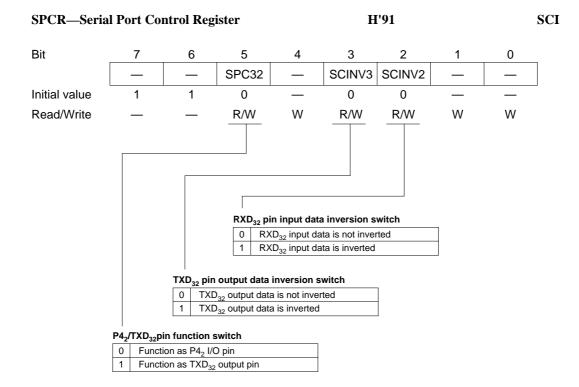
B.2 Functions



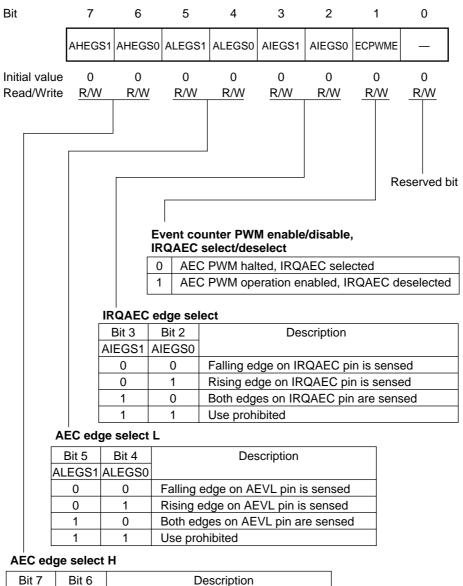


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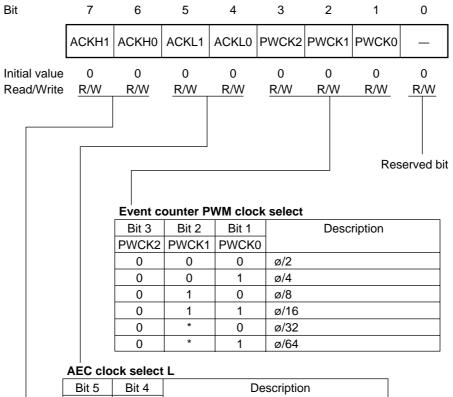






DIL /	ыю	Description
AHEGS1	AHEGS0	
0	0	Falling edge on AEVH pin is sensed
0	1	Rising edge on AEVH pin is sensed
1	0	Both edges on AEVH pin are sensed
1	1	Use prohibited

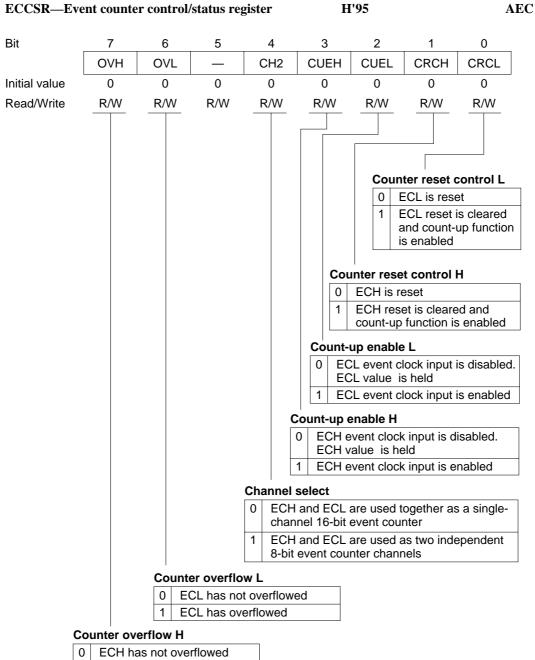
AEC



Bit 5	Bit 4	Description
ACKL1	ACKL0	
0	0	AEVL pin input
0	1	ø/2
1	0	ø/4
1	1	ø/8

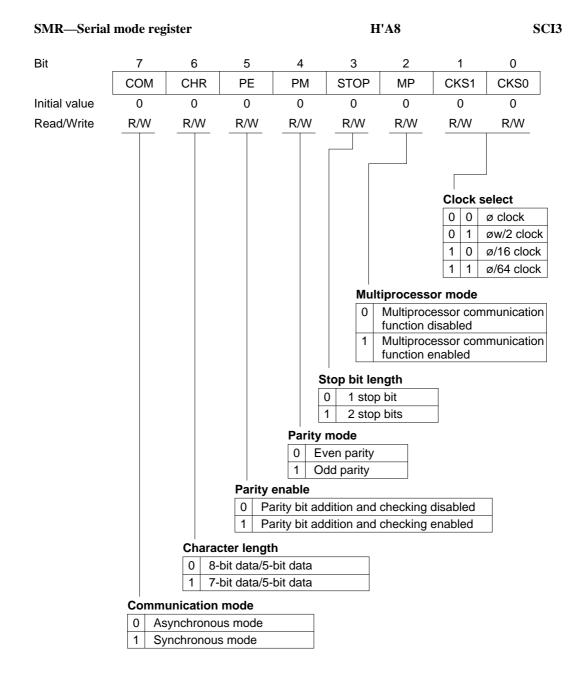
AEC clock select H

Bit 7	Bit 6	Description
ACKH1	ACKH0	
0	0	AEVH pin input
0	1	ø/2
1	0	ø/4
1	1	ø/8



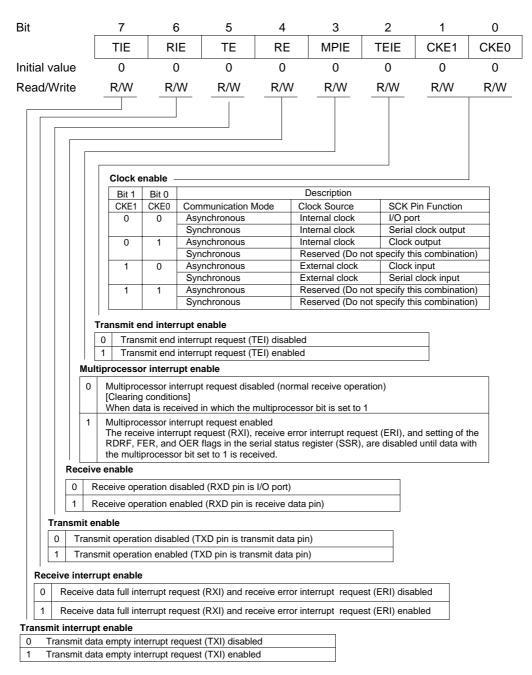
1 ECH has overflowed

ECH—Event	counter H	ł		H'96				AEC	
Bit	7	6	5	4	3	2	1	0	
	ECH7	ECH6	ECH5	ECH4	ECH3	ECH2	ECH1	ECH0	
Initial value	0	0	0	0	0	0	0	0	
Read/Write	R	R	R	R	R	R	R	R	
Read/White	<u></u>	<u>к</u>	<u>к</u>		к 	IX			
ECL—Event			<u>к</u>			' 97			AEC
			5	4			1		AEC
ECL—Event	counter L	4			Н	'97			AEC
ECL—Event	counter L	6	5	4	H 3	' 97 2	1	0	AEC



BRR—Bit rat	e register			H'A9				S	CI3
Bit	7	6	5	4	3	2	1	0	
	BRR7	BRR6	BRR5	BRR4	BRR3	BRR2	BRR1	BRR0	
Initial value	1	1	1	1	1	1	1	1	
Read/Write	R/W	R/W	R/W	R/W	R/W	R/W	R/W	R/W	

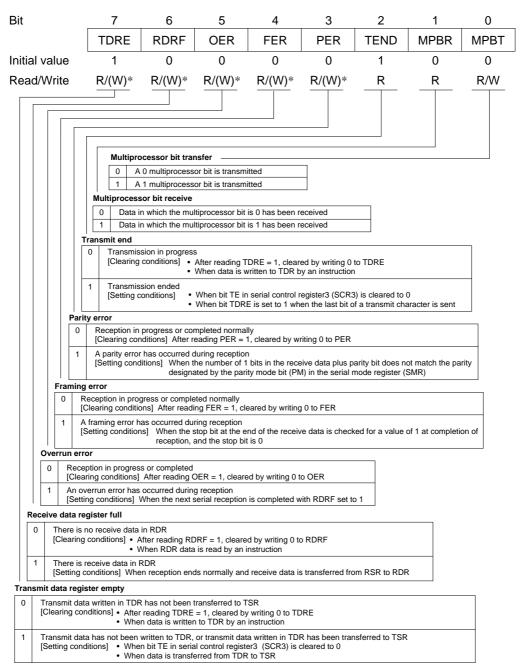
SCR3—Serial control register3



TDR—Transn	nit data r	egister		H'AB				SCI3	
Bit	7	6	5	4	3	2	1	0	
	TDR7	TDR6	TDR5	TDR4	TDR3	TDR2	TDR1	TDR0	
Initial value	1	1	1	1	1	1	1	1	-
Read/Write	R/W	R/W	R/W	R/W	R/W	R/W	R/W	R/W	

Data for transfer to TSR

SSR—Serial status register



Note: * Only a write of 0 for flag clearing is possible.

Bit	7	6	5	4	3	2	1	0	
2	RDR7	RDR6	RDR5	RDR4	RDR3	RDR2	RDR1	RDR0	
Initial value	0	0	0	0	0	0	0	0	
Read/Write	R	R	R	R	R	R	R	R	
TMA—Time	r mode reş	gister A			Н	'B0		Time	er A
TMA—Time	r mode reş	gister A			Н	'B0		Time	er A
TMA—Time Bit	r mode reş 7	gister A 6	5	4	Н 3	' B0 2	1	Time 0	er A
	r mode reg 7 	-	5	4			1 TMA1		er A

_

W

W

H'AD

SCI3

RDR—Receive data register

W

Read/Write

TNAAO	TNAAO	TNAA 4	TNAAO		r and Divider Ratio	Function
-	TMA2		-		ow Period	Function
0	0	0	0	PSS	ø/8192	
			1	PSS	ø/4096	timer
		1	0	PSS	ø/2048	
			1	PSS	ø/512	
	1	0	0	PSS	ø/256	
			1	PSS	ø/128	
		1	0	PSS	ø/32	
			1	PSS	ø/8	
1	0	0	0	PSW	1 s	Time
			1	PSW	0.5 s	base
		1	0	PSW	0.25 s	(when using
			1	PSW	0.03125 s	32.768 kHz)
	1	0	0	PSW and	d TCA are reset	
			1			
		1	0			
			1			

R/W

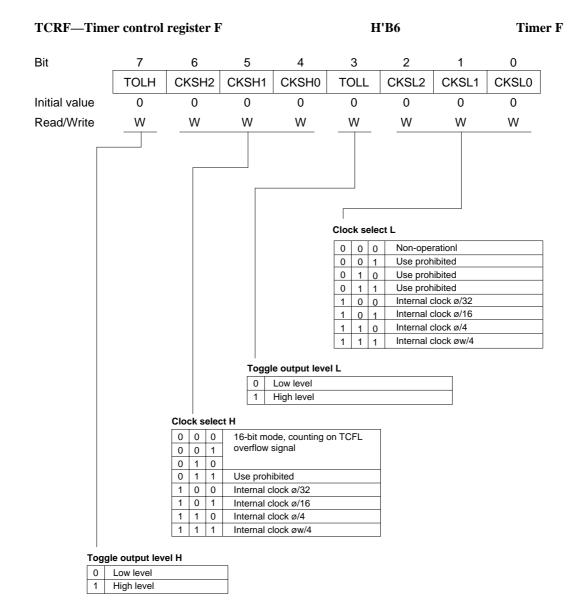
R/W

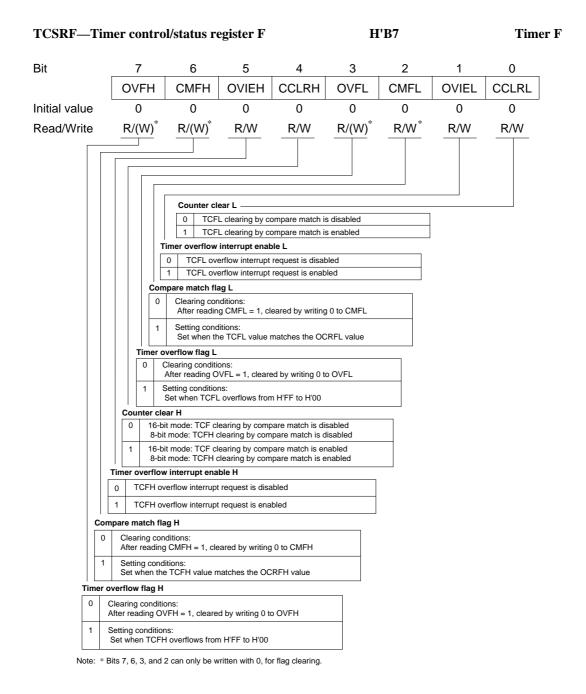
R/W

R/W

TCA—Timer	counter A	A			Н	Tim	er A		
Bit	7	6	5	4	3	2	1	0	7
	TCA7	TCA6	TCA5	TCA4	TCA3	TCA2	TCA1	TCA0	
Initial value	0	0	0	0	0	0	0	0	-
Read/Write	R	R	R	R	R	R	R	R	

Count value





TCFH—8-bit	timer cou	inter FH			Н	'B8		Time	er F
Bit	7	6	5	4	3	2	1	0	
	TCFH7	TCFH6	TCFH5	TCFH4	TCFH3	TCFH2	TCFH1	TCFH0	
Initial value	0	0	0	0	0	0	0	0	
Read/Write	R/W	R/W	R/W	R/W	R/W	R/W	R/W	R/W	
				Coun	t value				
TCFL—8-bit timer counter FLH'B9Timer F									er F
Bit	7	6	5	4	3	2	1	0	
	TCFL7	TCFL6	TCFL5	TCFL4	TCFL3	TCFL2	TCFL1	TCFL0	
Initial value	0	0	0	0	0	0	0	0	
Read/Write	R/W	R/W	R/W	R/W	R/W	R/W	R/W	R/W	
				Coun	t value				
OCRFH—Ou	tput com	pare regis	ter FH		Н	'BA		Time	er F
Bit	7	6	5	4	3	2	1	0	
	OCRFH7	OCRFH6	OCRFH5	OCRFH4	OCRFH3	OCRFH2	OCRFH1	OCRFH0	
Initial value	1	1	1	1	1	1	1	1	
Read/Write	R/W	R/W	R/W	R/W	R/W	R/W	R/W	R/W	
OCRFL—Ou	tput comp	pare regis	ter FL		Н	'BB		Time	er F
Bit	7	6	5	4	3	2	1	0	
	OCRFL7	OCRFL6	OCRFL5	OCRFL4	OCRFL3	OCRFL2	OCRFL1	OCRFL0	
Initial value	1	1	1	1	1	1	1	1	
Read/Write	R/W	R/W	R/W	R/W	R/W	R/W	R/W	R/W	

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LP	CR—	-LCD	port	control	register
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H'C0



Bit	7	6	5	4	3	2	1	0
	DTS1	DTS0	CMX	—	SGS3	SGS2	SGS1	SGS0
Initial value	0	0	0		0	0	0	0
Read/Write	R/W	R/W	R/W	W	R/W	R/W	R/W	R/W

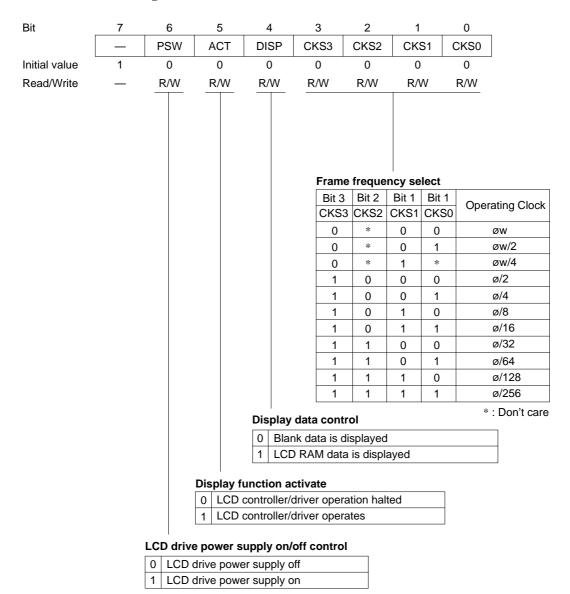
Bit 3	Bit 2	Bit 1	Bit 0		Function	n of Pins	s SEG ₂₅	to SEG	1		
SGS3	SGS2	SGS1	SGS0	SEG ₂₅	SEG24 to	SEG20 to	SEG16 to	SEG12 to	SEG ₈ to	SEG4 to	Notes
0000	0002	0001	0000		SEG ₂₁	SEG17	SEG ₁₃	SEG ₉	SEG₅	SEG1	
0	0	0	0	Port	Port	Port	Port	Port	Port	Port	(Initial value
0	0	0	1	Port	Port	Port	Port	Port	Port	SEG	
0	0	1	0	Port	Port	Port	Port	Port	SEG	SEG]
0	0	1	1	Port	Port	Port	Port	SEG	SEG	SEG	1
0	1	0	0	Port	Port	Port	SEG	SEG	SEG	SEG	1
0	1	0	1	Port	Port	SEG	SEG	SEG	SEG	SEG	
0	1	1	0	Port	SEG	SEG	SEG	SEG	SEG	SEG	
0	1	1	1	SEG	SEG	SEG	SEG	SEG	SEG	SEG	
1	0	0	0	SEG	SEG	SEG	SEG	SEG	SEG	SEG	
1	0	0	1	SEG	SEG	SEG	SEG	SEG	SEG	Port	
1	0	1	0	SEG	SEG	SEG	SEG	SEG	Port	Port	1
1	0	1	1	SEG	SEG	SEG	SEG	Port	Port	Port	
1	1	0	0	SEG	SEG	SEG	Port	Port	Port	Port	1
1	1	0	1	SEG	SEG	Port	Port	Port	Port	Port	1
1	1	1	0	SEG	Port	Port	Port	Port	Port	Port	1
1	1	1	1	Port	Port	Port	Port	Port	Port	Port	1

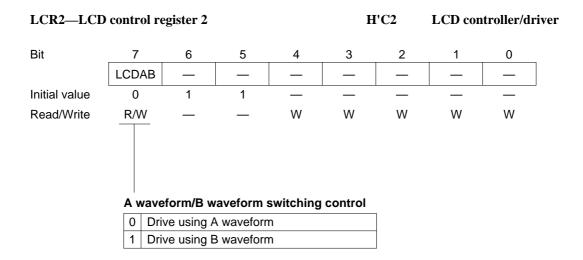
Duty select, common function select

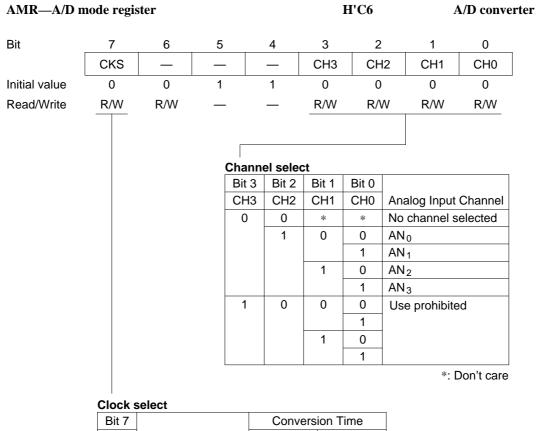
Bit 7	Bit 6	Bit 5	Duty Cycle	Common Drivers	Notes
DIS1	DTS0	CMX			
0	0	0	Static	COM ₁	
		1	Otatic	COM ₄ to COM ₁	COM ₄ to COM ₂ output the same waveform as COM ₁
0	1	0	1/2 duty	COM ₂ to COM ₁	
		1	1/2 Guty	COM ₄ to COM ₁	COM ₄ outputs the same waveform as COM ₃ and COM ₂ outputs the same waveform as COM ₁
1	0	0	1/3 duty	COM ₃ to COM ₁	
		1	1/5 duty	COM ₄ to COM ₁	COM ₄ outputs a non-selected waveform
1	1	0	1/4 duty	COM4 to COM1	_

LCR—LCD control register

H'C1 LCD controller/driver







BIt /		Convers	ion Time
CKS	Conversion Period	ø = 1 MHz	ø = 5 MHz
0	62/ø	62 µs	12.4 μs
1	31/ø	31 µs	*

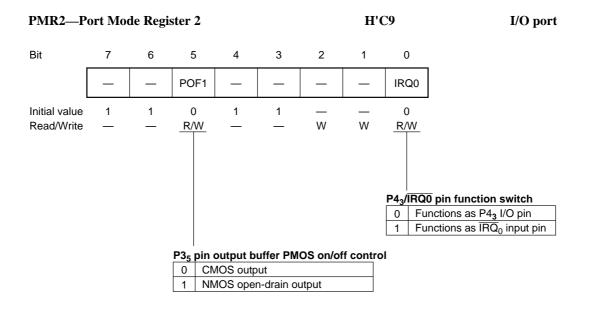
*: Operation is not guaranteed with a conversion time of less than 12.4 $\mu s.$

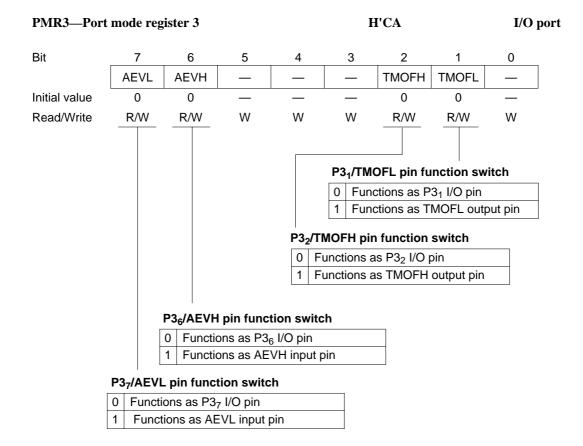
Select a setting that gives a conversion time of at least 12.4 μ s.

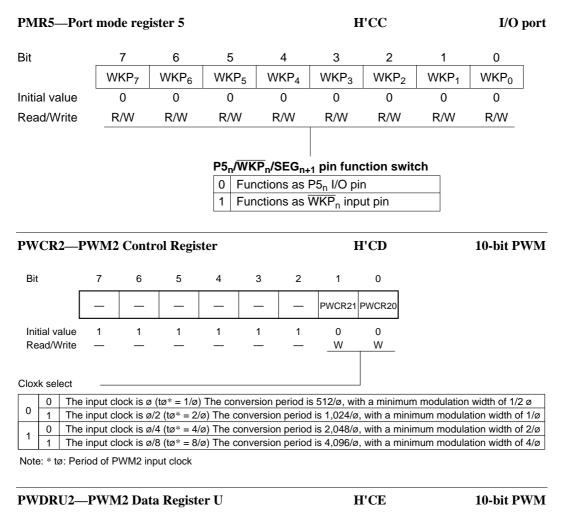
ADRRH—A/D result register H	H'C4	A/D converter
ADRRL—A/D result register L	H'C5	

ADRRH

Bit	7	6	5	4	3	2	1	0
	ADR9	ADR8	ADR7	ADR6	ADR5	ADR4	ADR3	ADR2
Initial value	Not fixed	Not fixed	Not fixed	Not fixed	Not fixed	Not fixed	Not fixed	Not fixed
Read/Write	R	R	R	R	R	R	R	R
	A/D conversion result							
ADRRL								
Bit	7	6	5	4	3	2	1	0
	ADR1	ADR0		_		_	_	
Initial value	Not fixed	Not fixed						
Read/Write	R	R		—	—	—	—	—
	A/D conv	ersion resu	ult					
ADSR—A/D	start regi	ster			Н	['C7	A	A/D convei
ADSR—A/D Bit) start regis	ster 6	5	4	H 3	2 ¹	A 1	A/D conver
	_		5	4				
	7		5 — 1	4				
Bit	7 ADSF	6	_	—	3	2	1	0
Bit Initial value	7 ADSF 0 R/W	6	_	—	3	2	1	0
Bit Initial value	7 ADSF 0 R/W A/D s	6 — 1 —	1	—	3 — 1 —	2	1	0
Bit Initial value	7 ADSF 0 R/W A/D s	6 1 	1	1 — Deletion of A	3 — 1 —	2	1	0
Bit Initial value	7 ADSF 0 R/W A/D s 0 R v	6 1 	1 	1 — Deletion of A	3 — 1 — A/D conver	2 1 rsion	1	0







Bit	7	6	5	4	3	2	1	0
	_	_		_			PWDRU21	PWDRU20
Initial value	1	1	1	1	1	1	0	0
Read/Write	—	_	—	_	_	_	W	W

Upper 2 bits of PWM2 waveform generation data

PWDRL2-	PWDRL2—PWM2 Data Register L						H'CF			
Bit	7	6	5	4	3	2	1	0		
	PWDRL27	PWDRL26	PWDRL25	PWDRL24	PWDRL23	PWDRL22	PWDRL21	PWDRL20		
Initial value	0	0	0	0	0	0	0	0		
Read/Write	W	W	W	W	W	W	W	W		
PWCR1—I	PWM1 c			PWM2 v	vaveform		on data		10-bit PW	VM
Bit	7		6	5	4	3	2	1	0	
	_		_	_	_	—	—	PWCR1 ₁	PWCR10	
Initial value	1		1	1	1	1	1	0	0	
Read/Write	_		_	_	—	—	—	W	W	
C	lock se	lect —								
Г	0 The i	nnut cloc	k is ø (tø	$* - 1/\alpha$]

0	The input clock is \emptyset (t \emptyset * = 1/ \emptyset)
	The conversion period is 512/ø, with a minimum modulation width of 1/2ø
	The input clock is $\emptyset/2$ (t $\emptyset^* = 2/\emptyset$)
	The conversion period is 1,024/ø, with a minimum modulation width of $1/ø$
1	The input clock is $\emptyset/4$ (t $\emptyset^* = 4/\emptyset$)
	The conversion period is 2,048/ø, with a minimum modulation width of 2/ø
	The input clock is Ø/8 (tø* = 8/Ø)
	The conversion period is 4,096/ø, with a minimum modulation width of 4/ø

Note:* tø: Period of PWM input clock

PWDRU1—	PWM1 dat	a register	·U		Н	10-bit PWN		
Bit	7	6	5	4	3	2	1	0
	_	—	_	_	_	—	PWDUR11	PWDRU10
Initial value	1	1	1	1	1	1	0	0
Read/Write	—	—	—		—	—	W	W
PWDRL1—	PWM1 dat	a register		2 bits of	data for ge H	enerating F	PWM1 wav	veform 10-bit PWN
Bit	7	6	5	4	3	2	1	0
		PWDRL1 ₆			PWDRL13	1		PWDRL10
Initial value	0	0	0	0	0	0	0	0
Read/Write	W	W	W	W	W	W	W	W
PDR3—Port	data regis	ter 3			Н	' D 6		I/O port
Bit	7	6	5	4	3	2	1	0
BR	P3 ₇	P3 ₆	P3 ₅	P3₄	P3 ₃	P3 ₂	P31	
Initial value	0	0	0	0	0	0	0	
Read/Write	R/W	R/W	R/W	R/W	R/W	R/W	R/W	_
			Stores	data of po	ort 3 pins			
PDR4—Port	data regis	ter 4			Н	['D7		I/O port
Bit	7	6	5	4	3	2	1	0
	_	_	_		P43	P42	P4 ₁	P4 ₀
Initial value	1	1	1	1	1	0	0	0
Read/Write	—	—			R	R/W	R/W	R/W
				Read	ds P4 ₃ pin		data of po	rt 4 pins

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PDR5—Port o	lata regis	ter 5			Н	I/O ports			
Bit	7	6	5	4	3	2	1	0	
	P57	P5 ₆	P5 ₅	P5 ₄	P5 ₃	P5 ₂	P5 ₁	P5 ₀	
Initial value	0	0	0	0	0	0	0	0	
Read/Write	R/W	R/W	R/W	R/W	R/W	R/W	R/W	R/W	
			Sto	ores data	of port 5 p	ins			
PDR6—Port	lata regis	ter 6			Н	'D9		I/O p	orts
Bit	7	6	5	4	3	2	1	0	
	P6 ₇	P6 ₆	P6 ₅	P6 ₄	P6 ₃	P6 ₂	P6 ₁	P6 ₀	
Initial value	0	0	0	0	0	0	0	0	
Read/Write	R/W	R/W	R/W	R/W	R/W	R/W	R/W	R/W	
			Ste	ores data					
PDR7_Port d	lata regis	ter 7			н	'DA		I/O n	orts
PDR7—Port o	lata regis	ter 7			Н	'DA		I/O p	orts
PDR7—Port o	lata regis 7	ter 7 6	5	4	н 3	' DA 2	1	І/О р 0	orts
			5 P7 ₅	4 P7 ₄			1 P7 ₁	_	orts
	7	6			3	2		0	orts
Bit	7 P7 ₇	6 P7 ₆	P7 ₅	P7 ₄	3 P7 ₃	2 P7 ₂	P7 ₁	0 P7 ₀	orts
Bit Initial value	7 P7 ₇ 0	6 P7 ₆ 0	P7 ₅ 0 R/W	P7 ₄ 0	3 P7 ₃ 0 R/W	2 P7 ₂ 0 R/W	P7 ₁ 0	0 P7 ₀ 0	orts
Bit Initial value	7 P7 ₇ 0 R/W	6 P7 ₆ 0 R/W	P7 ₅ 0 R/W	P7 ₄ 0 R/W	3 P7 ₃ 0 R/W of port 7 p	2 P7 ₂ 0 R/W	P7 ₁ 0	0 P7 ₀ 0	
Bit Initial value Read/Write	7 P7 ₇ 0 R/W	6 P7 ₆ 0 R/W	P7 ₅ 0 R/W	P7 ₄ 0 R/W	3 P7 ₃ 0 R/W of port 7 p	2 P72 0 R/W	P7 ₁ 0	0 P7 ₀ 0 R/W	
Bit Initial value Read/Write PDR8—Port o	7 P7 ₇ 0 R/W	6 P7 ₆ 0 R/W	P7 ₅ 0 R/W Sto	P7 ₄ 0 R/W pres data o	3 P7 ₃ 0 R/W of port 7 p	2 P72 0 R/W	P7 ₁ 0 R/W	0 P7 ₀ 0 R/W	
Bit Initial value Read/Write PDR8—Port o	7 P7 ₇ 0 R/W	6 P7 ₆ 0 R/W	P7 ₅ 0 R/W Sto	P7 ₄ 0 R/W pres data o	3 P7 ₃ 0 R/W of port 7 p	2 P72 0 R/W	P7 ₁ 0 R/W	0 P7 ₀ 0 R/W I/O p	
Bit Initial value Read/Write PDR8—Port of Bit	7 P7 ₇ 0 R/W	6 P7 ₆ 0 R/W	P7 ₅ 0 R/W Sto	P7 ₄ 0 R/W pres data o	3 P7 ₃ 0 R/W of port 7 p	2 P72 0 R/W	P7 ₁ 0 R/W	0 P7 ₀ 0 R/W I/O p 0 P8 ₀	
Bit Initial value Read/Write PDR8—Port of Bit Initial value	7 P7 ₇ 0 R/W	6 P7 ₆ 0 R/W	P7 ₅ 0 R/W Sto	P7 ₄ 0 R/W pres data o	3 P7 ₃ 0 R/W of port 7 p	2 P72 0 R/W ins 'DB 2 	P71 0 R/W	0 P7 ₀ 0 R/W I/O p 0 P8 ₀ 0	orts

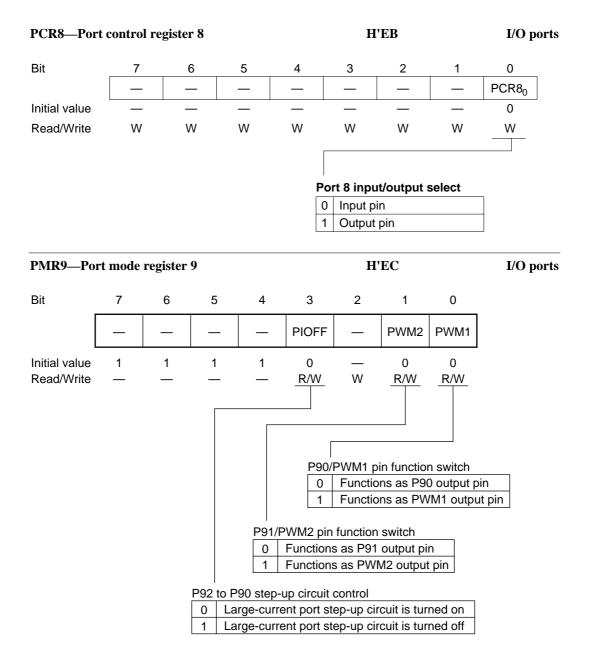
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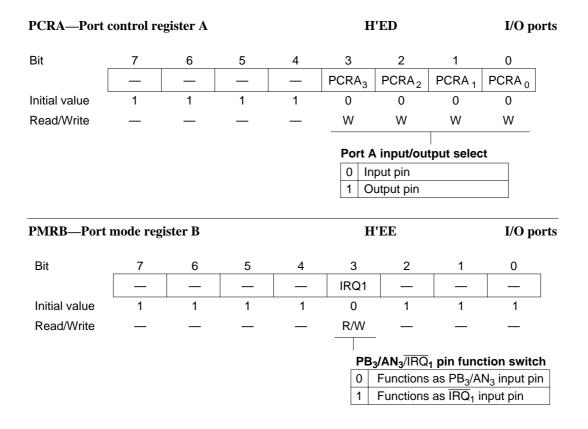
PDR9—Por			I/O ports							
Bit	7	6	5	4	3	2	1	0		
		_	P9 ₅	P9 ₄	P9 ₃	P9 ₂	P9 ₁	P9 ₀		
Initial value Read/Write	1	1	1 R/W	1 R/W	1 R/W	1 R/W	1 R/W	1 R/W		
				Sto	ores data	of port 9	pins			
PDRA—Poi	rt data r	egister A	\			Н	DD		I/O	po
Bit	7	(6	5	4	3	2	1	0	
	_	-	-	_	_	PA ₃	PA ₂	PA ₁	PA ₀	
Initial value	1		1	1	1	0	0	0	0	
Read/Write	_	-	_	_	—	R/W	R/W	R/W	R/W	-
								of port A p		
PDRB—Por	rt data r	egister F	3				res data o DE	of port A p	ins I/O	po
PDRB—Por	rt data r 7	-	B 6	5	4			of port A p		po
		-		5	4	H,	DE		I/O	po
		-		5	4	H' 3	2	1	I/O 0	p
Bit Initial value		-		5	4	H' 3 PB ₃ R	DE 2 PB ₂ R	1 PB ₁	I/O 0 PB ₀ R	p
Bit Initial value	7	-	6	_	4	H ¹ 3 PB ₃ R Rea	DE 2 PB ₂ R	1 PB ₁ R	I/O 0 PB ₀ R	
Bit Initial value Read/Write	7	up contr	6	_	4	H ¹ 3 PB ₃ R Rea	DE 2 PB ₂ R ds states	1 PB ₁ R	I/O 0 PB ₀ R bins	
Bit Initial value Read/Write PUCR3—Pe	7 	up contr	6 rol regist	 er 3 5	4	H ¹ 3 PB ₃ R Rea H ¹ 3	DE 2 PB ₂ R ds states E1	1 PB ₁ R of port B p	I/O 0 PB ₀ R bins	
Bit Initial value Read/Write PUCR3—Pe	7 	up contr (1) (1) (1) (1) (1) (1) (1) (1) (1) (1)	6 rol regist	 er 3 5	4	H ¹ 3 PB ₃ R Rea H ¹ 3	DE 2 PB ₂ R ds states E1 2	1 PB ₁ R of port B p	I/O 0 PB ₀ R bins	

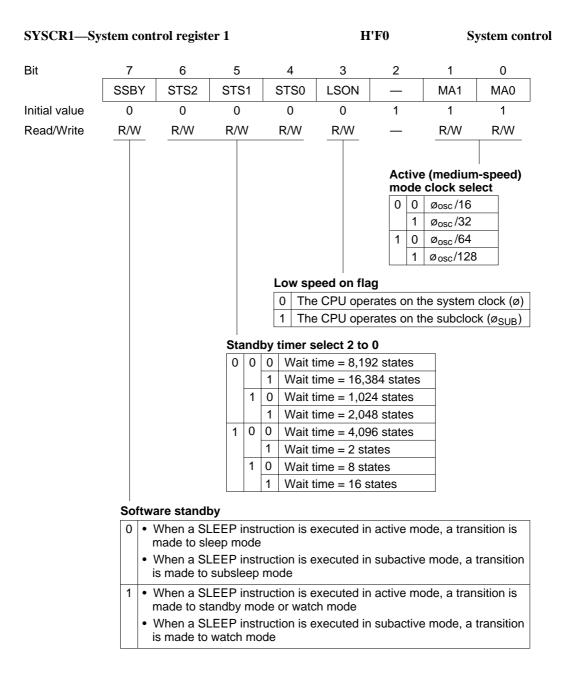
PUCR5—Por	t pull-up (control re	gister 5		Н	'E2		I/O p	orts
Bit	7	6	5	4	3	2	1	0	
	PUCR57	PUCR5 ₆	PUCR55	PUCR5 ₄	PUCR53	PUCR5 ₂	PUCR51	PUCR50	
Initial value	0	0	0	0	0	0	0	0	
Read/Write	R/W	R/W	R/W	R/W	R/W	R/W	R/W	R/W	
PUCR6—Por	t pull-up (control re	gister 6		Н	['E3		I/O p	orts
Bit	7	6	5	4	3	2	1	0	
	PUCR67	PUCR6 ₆	PUCR65	PUCR6 ₄	PUCR63	PUCR6 ₂	PUCR61	PUCR60	
Initial value	0	0	0	0	0	0	0	0	
Read/Write	R/W	R/W	R/W	R/W	R/W	R/W	R/W	R/W	
PCR3—Port o	control re	gister 3			Н	['E6		I/O p	orts
Bit	7	6	5	4	3	2	1	0	
	PCR37	PCR3 ₆	PCR3 ₅	PCR3 ₄	PCR3 ₃	PCR3 ₂	PCR3 ₁	_	
Initial value	0	0	0	0	0	0	0		
Read/Write	W	W	W	W	W	W	W	W	
				Р	∣ ort 3 inpu	t/output s	elect		
				(
					1 Output	pin			
PCR4—Port o	control re	gister 4			Н	['E7		I/O p	orts
Bit	7	6	5	4	3	2	1	0	
	_		—	_	—	PCR4 ₂	PCR4 ₁	PCR4 ₀	
Initial value	1	1	1	1	1	0	0	0	
Read/Write	—	—	_	_	_	W	W	W	
					Р	ort 4 inpu	it/output s	select	
					(0 Input p	in		
						1 Output	pin		

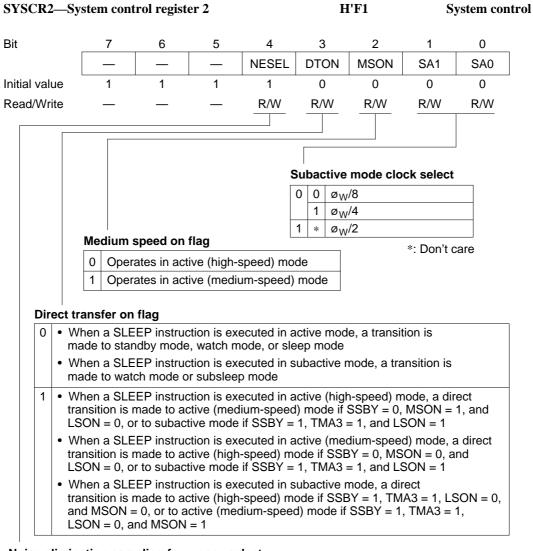
PCR5—Port	control re	gister 5		H'E8				I/O p	orts
Bit	7	6	5	4	3	2	1	0	_
	PCR57	PCR5 ₆	PCR55	PCR5 ₄	PCR5 ₃	PCR5 ₂	PCR5 ₁	PCR50]
Initial value	0	0	0	0	0	0	0	0	
Read/Write	W	W	W	W	W	W	W	W	
				P	ort 5 inpu		elect		
				1					
PCR6—Port	control re	gister 6			Н	'E9		I/O p	orts
Bit	7	6	5	4	3	2	1	0	
	PCR67	PCR6 ₆	PCR65	PCR6 ₄	PCR6 ₃	PCR6 ₂	PCR6 ₁	PCR6 ₀]
Initial value	0	0	0	0	0	0	0	0	1
Read/Write	W	W	W	W	W	W	W	W	
				P	ort 6 inpu	t/output s	elect		
) Input pi I Output				
PCR7—Port	control re	gister 7			Н	'EA		I/O p	orts
Bit	7	6	5	4	3	2	1	0	
	PCR77	PCR7 ₆	PCR75	PCR7 ₄	PCR73	PCR7 ₂	PCR7 ₁	PCR70]
Initial value	0	0	0	0	0	0	0	0	1
Read/Write	W	W	W	W	W	W	W	W	
				P	ort 7 inpu	t/output s	elect		
				(
				1	Output	pin			

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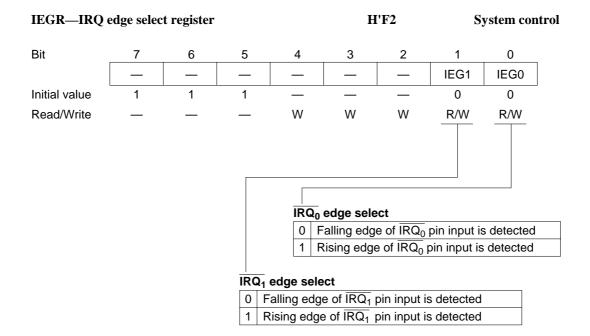


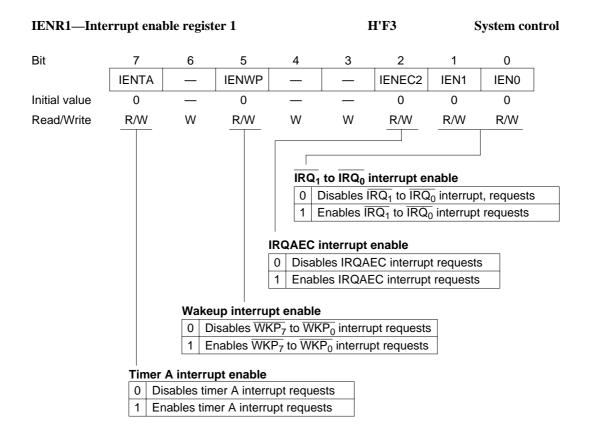


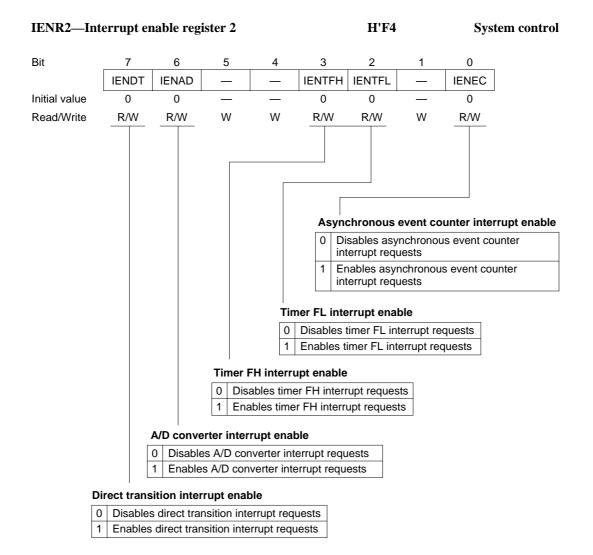


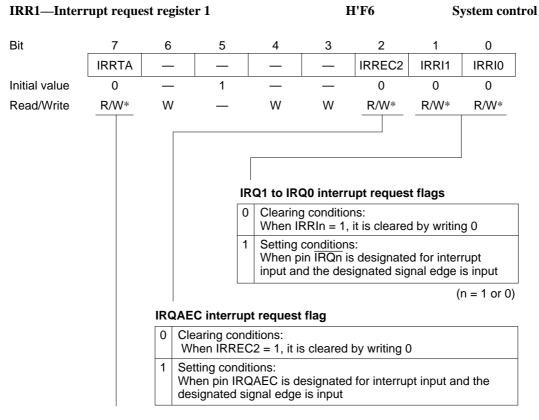
Noise elimination sampling frequency select

0	Sampling rate is ø _{OSC} /16 Sampling rate is ø _{OSC} /4
1	Sampling rate is Ø _{OSC} /4









Timer A interrupt request flag

0	Clearing conditions: When IRRTA = 1, it is cleared by writing 0
1	Setting conditions: When the timer A counter value overflows (from H'FF to H'00)

Note: * Bits 7 and 2 to 0 can only be written with 0, for flag clearing.

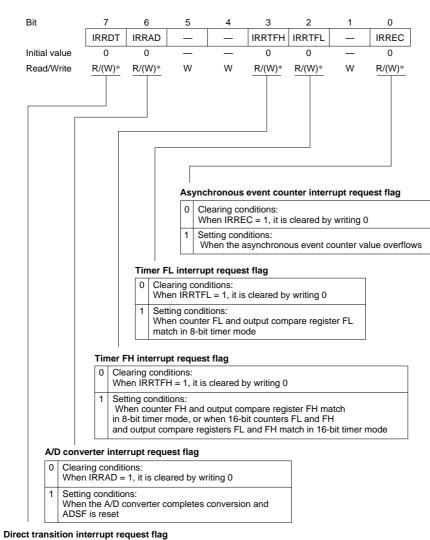
HITACHI

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IRR2—Interrupt request register 2

H'F7





0	Clearing conditions: When IRRDT = 1, it is cleared by writing 0
1	Setting conditions: When a SLEEP instruction is executed while DTON is set to 1, and a direct transition is made

Note: * Bits 7, 6, 3, 2, and 0 can only be written with 0, for flag clearing.

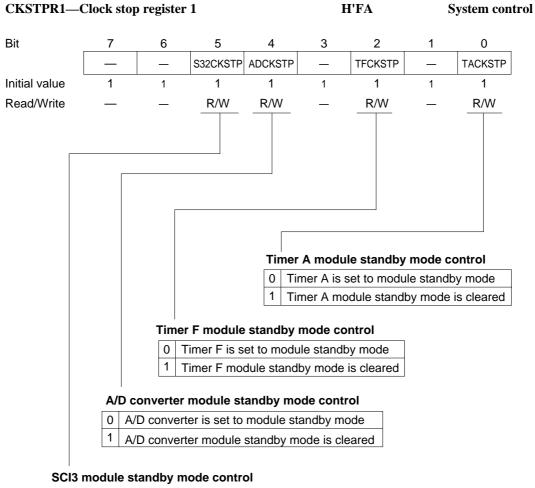
IWPR—Wakeup interrupt request register						'F9	System control		
Bit	7	6	5	4	3	2	1	0	
	IWPF7	IWPF6	IWPF5	IWPF4	IWPF3	IWPF2	IWPF1	IWPF0	
Initial value	0	0	0	0	0	0	0	0	
Read/Write	R/(W)*	R/(W)*							

Wakeup interrupt request register

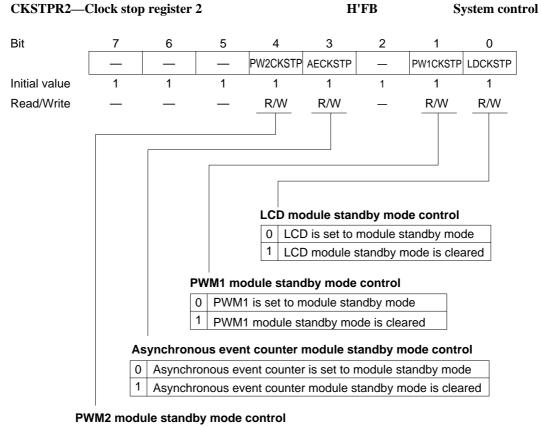
0	Clearing conditions: When IWPFn = 1, it is cleared by writing 0
1	Setting conditions: When pin WKPn is designated for wakeup input and a falling edge is input at that pin

(n = 7 to 0)

Note: * All bits can only be written with 0, for flag clearing.



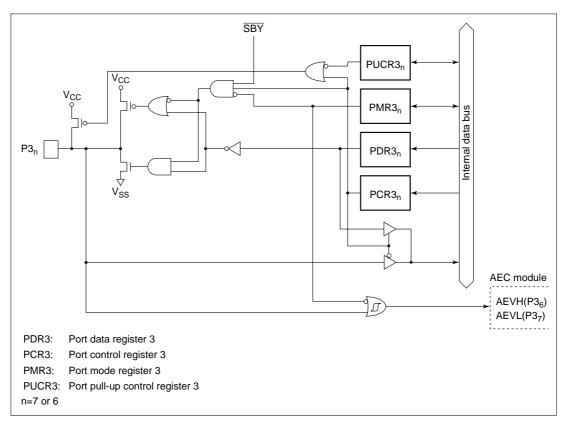
	-
0	SCI3 is set to module standby mode
1	SCI3 module standby mode is cleared



C)	PWM2 is set to module standby mode
1		PWM2 module standby mode is cleared

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Appendix C I/O Port Block Diagrams



C.1 Block Diagrams of Port 3

Figure C.1 (a) Port 3 Block Diagram (Pins P37 and P36)

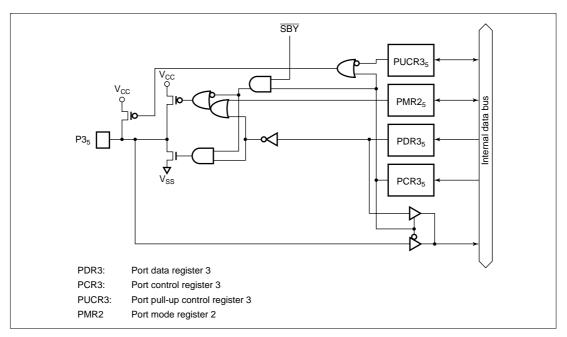


Figure C.1 (b) Port 3 Block Diagram (Pin P3₅)

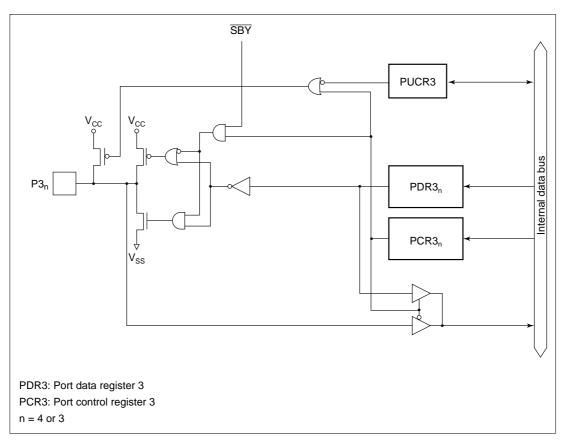


Figure C.1 (c) Port 3 Block Diagram (Pins P3₄ and P3₃)

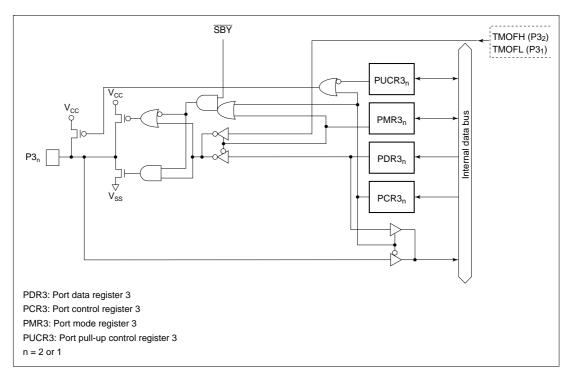


Figure C.1 (d) Port 3 Block Diagram (Pins $P3_2$ and $P3_1$)

C.2 Block Diagrams of Port 4

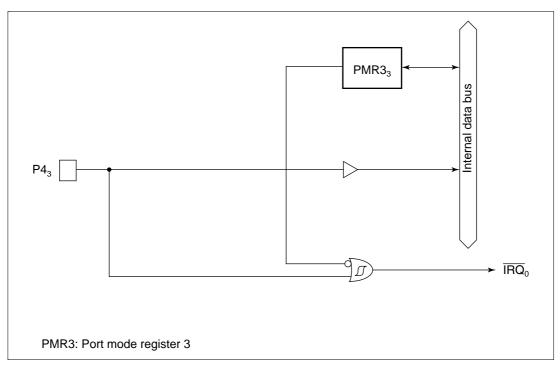


Figure C.2 (a) Port 4 Block Diagram (Pin P4₃)

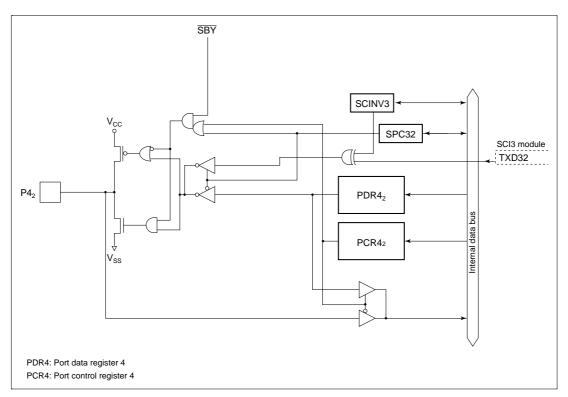


Figure C.2 (b) Port 4 Block Diagram (Pin P4₂)

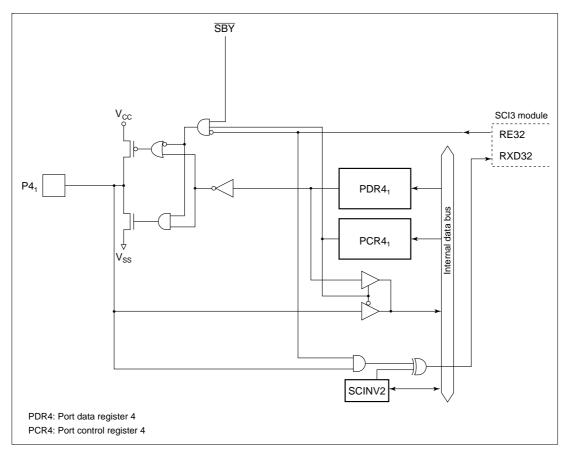


Figure C.2 (c) Port 4 Block Diagram (Pin P4₁)

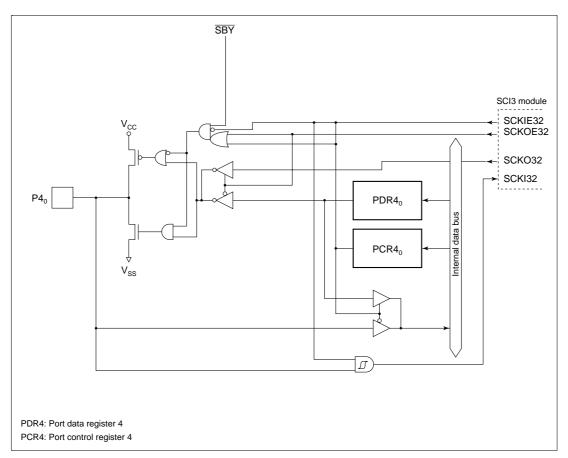


Figure C.2 (d) Port 4 Block Diagram (Pin P4₀)



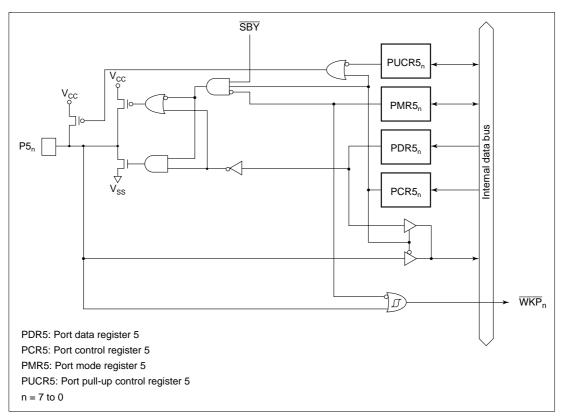


Figure C.3 Port 5 Block Diagram

C.4 Block Diagram of Port 6

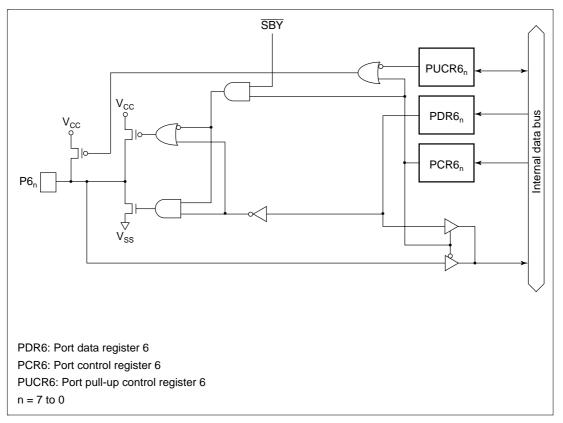


Figure C.4 Port 6 Block Diagram

C.5 Block Diagram of Port 7

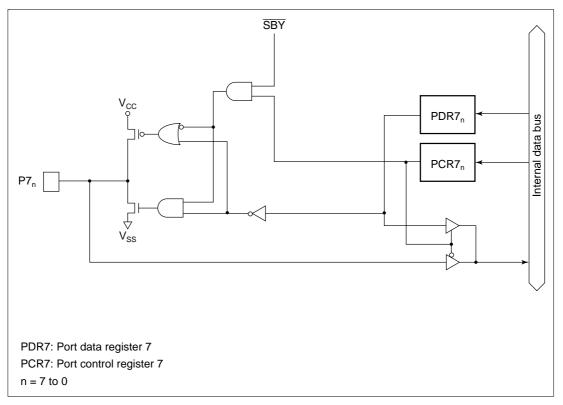


Figure C.5 Port 7 Block Diagram

C.6 Block Diagrams of Port 8

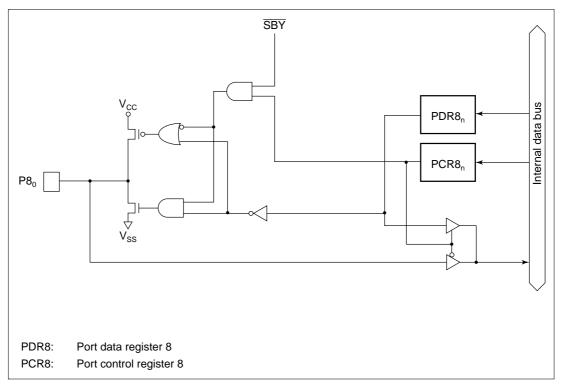


Figure C.6 Port 8 Block Diagram (Pin P8₀)

C.7 Block Diagrams of Port 9

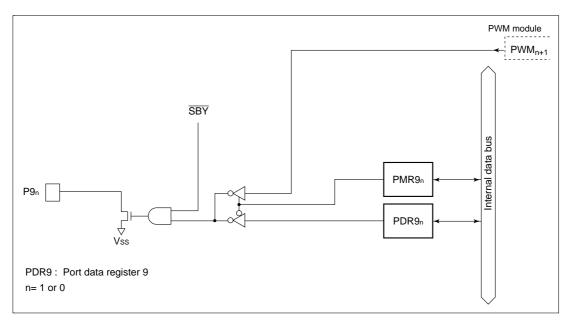
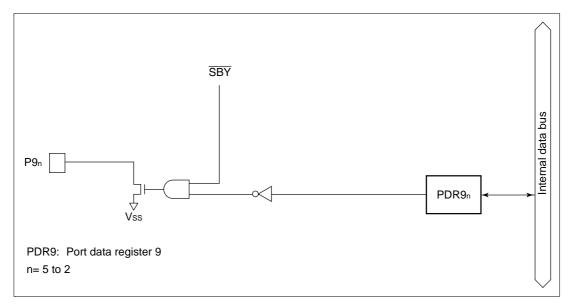
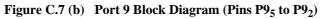


Figure C.7 (a) Port 9 Block Diagram (Pins P9₁ and P9₀)





C.8 Block Diagram of Port A

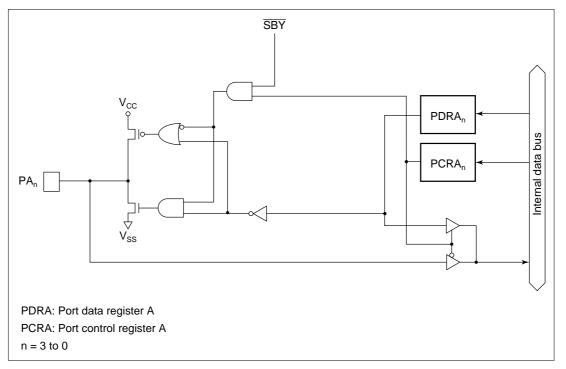


Figure C.8 Port A Block Diagram

C.9 Block Diagram of Port B

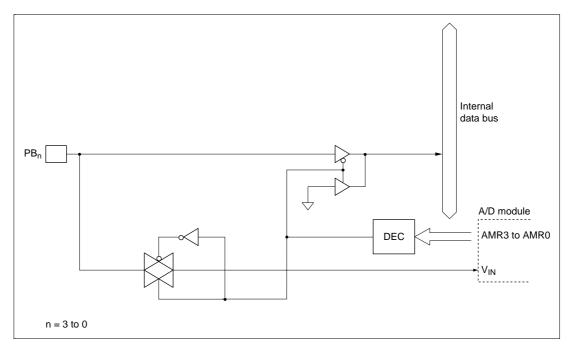


Figure C.9 Port B Block Diagram

Appendix D Port States in the Different Processing States

Port	Reset	Sleep	Subsleep	Standby	Watch	Subactive	Active
P3 ₇ to P3 ₁	High impedance	Retained	Retained	High impedance*	Retained	Functions	Functions
P4 ₃ to P4 ₀	High impedance	Retained	Retained	High impedance	Retained	Functions	Functions
P5 ₇ to P5 ₀	High impedance	Retained	Retained	High impedance*	Retained	Functions	Functions
P6 ₇ to P6 ₀	High impedance	Retained	Retained	High impedance	Retained	Functions	Functions
P7 ₇ to P7 ₀	High impedance	Retained	Retained	High impedance	Retained	Functions	Functions
P8 ₇ to P8 ₀	High impedance	Retained	Retained	High impedance	Retained	Functions	Functions
P9 ₅ to P9 ₀	High impedance	Retained	Retained	High impedance*	Retained	Functions	Functions
PA ₃ to PA ₀	High impedance	Retained	Retained	High impedance	Retained	Functions	Functions
PB ₃ to PB ₀	High impedance	High impedance	High impedance	High impedance	High impedance	High impedance	High impedance

Table D.1Port States Overview

Note: * High level output when MOS pull-up is in on state.

Appendix E List of Product Codes

	Product Type		Product Code	Mark Code	Package(Hitachi Package Code)
H8/3802	H8/3802	Mask ROM versions	HD6433802H	HD6433802 (***) H	64-pin QFP (FP-64A)
series			HD64433802FP	HD6433802 (***) FP	64-pin LQFP (FP-64E)
			HD6433802P	HD6433802 (***) P	64-pin DILP (DP-64S)
		ZTAT versions	HD6473802H	HD6473802H	64-pin QFP (FP-64A)
			HD6473802FP	HD6473802FP	64-pin LQFP (FP-64E)
			HD6473802P	HD6473802P	64-pin DILP (DP-64S)
	H8/3801	Mask ROM versions	HD6433801H	HD6433801 (***) H	64-pin QFP (FP-64A)
			HD6433801FP	HD6433801 (***) FP	64-pin LQFP (FP-64E)
			HD6433801P	HD6433801 (***) P	64-pin DILP (DP-64S)
	H8/3800	Mask ROM versions	HD6433800H	HD6433800 (***) H	64-pin QFP (FP-64A)
			HD6433800FP	HD6433800 (***) FP	64-pin LQFP (FP-64E)
			HD6433800P	HD6433800 (***) P	64-pin DILP (DP-64S)

Table E.1 H8/3802 Series Product Code Lineup

Note: For mask ROM versions, (***) is the ROM code.

Appendix F Package Dimensions

Dimensional drawings of H8/3802 Series packages FP-64A, FP-64E, and DP-64S are shown in figures F.1, F.2, and F.3 below.

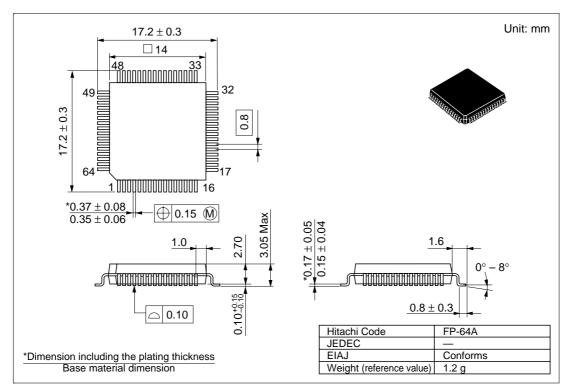


Figure F.1 FP-64A Package Dimensions

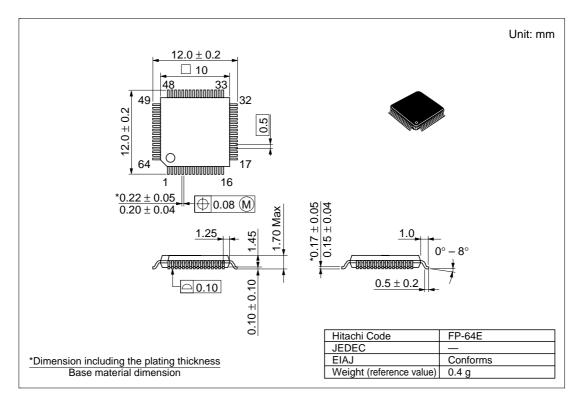


Figure F.2 FP-64E Package Dimensions

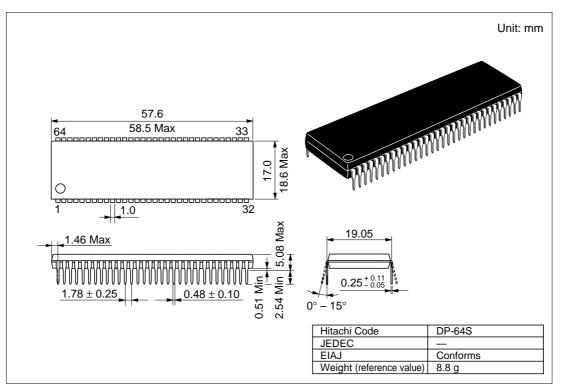


Figure F.3 DP-64S Package Dimensions

H8/3802 Series Hardware Manual

Publication Date:	1st Edition, November 1999			
	2nd Edition, January 2001			
Published by:	Electronic Devices Sales & Marketing Group			
	Semiconductor & Integrated Circuits			
	Hitachi, Ltd.			
Edited by:	Technical Documentation Group			
	Hitachi Kodaira Semiconductor Co., Ltd.			
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